# Appendix: Upper bounds for Model-Free Row-Sparse Principal Component Analysis 

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## 1. Problem Setting

The row-sparse principal component analysis problem is defined as follows: Given a sample covariance matrix $\boldsymbol{A} \in \mathbb{R}^{d \times d}$, a sparsity parameter $k(\leq d)$, the task is to find the top- $r k$-sparsity principal components $\boldsymbol{V}=\left(\boldsymbol{v}_{1}, \ldots, \boldsymbol{v}_{r}\right) \in \mathbb{R}^{d \times r}$,

$$
\begin{equation*}
\underset{\boldsymbol{V}=\boldsymbol{I}_{r},\|\boldsymbol{V}\|_{0} \leq k}{\arg \max } \operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{A} \boldsymbol{V}\right) . \tag{SPCA}
\end{equation*}
$$

where the row-sparsity constraint $\|\boldsymbol{V}\|_{0} \leq k$ denotes that there are at most $k$ non-zero rows in matrix $\boldsymbol{V}$, i.e., the principal components share the global support of cardinality at most $k$.

### 1.1. Notations

Let the bold upper case letters, for example, $\boldsymbol{A}, \boldsymbol{B}$ be matrices, and denote its $(i, j)$-th component as $[\boldsymbol{A}]_{i j},[\boldsymbol{B}]_{i j}$. Let $\operatorname{supp}(\boldsymbol{A})$ be the support of non-zero rows of matrix $\boldsymbol{A}$. Let the bold lower case letters, for example, $\boldsymbol{a}, \boldsymbol{b}$ be vectors, and denote its $i$-th component as $[\boldsymbol{a}]_{i},[\boldsymbol{b}]_{i}$. Let the regular case letters, for example, $I, J$ be the set of indices. For an integer $k$, let $[k]:=\{1, \ldots, k\}$. Given any matrix $\boldsymbol{A} \in \mathbb{R}^{n \times m}$ and $I \subseteq[n], J \subseteq[m]$, let $[\boldsymbol{A}]_{I, J}$ (or $\boldsymbol{A}_{I, J}$ in short) be the sub-matrix of $A$ with rows in $I$ and columns in $J$. In order to simplify notation, let $[\boldsymbol{A}]_{I}$ be the sub-matrix of $A$ when considering all rows in $I$, and let $[\boldsymbol{A}]_{j}$ be the $j$ th row of matrix $\boldsymbol{A}$. Let the regular lower case letters, for example, $\alpha, \beta$ be the reals. Let $\oplus$ be the sign of direct plus, i.e., given two square matrices $\boldsymbol{A} \in \mathbb{R}^{n_{1} \times n_{1}}, \boldsymbol{B} \in \mathbb{R}^{n_{2} \times n_{2}}, \boldsymbol{A} \oplus \boldsymbol{B}:=\operatorname{diag}(\boldsymbol{A}, \boldsymbol{B}) \in$ $\mathbb{R}^{\left(n_{1}+n_{2}\right) \times\left(n_{1}+n_{2}\right)}$.

## 2. Proof of Dual (Upper) Bounds

Let $\mathcal{F}$ be the feasible region of SPCA as

$$
\mathcal{F}:=\left\{\begin{array}{ll}
\boldsymbol{V} \in \mathbb{R}^{d \times r}: \begin{array}{l}
\boldsymbol{V}^{\top} \boldsymbol{V}=\boldsymbol{I}_{r} \\
\|\boldsymbol{V}\|_{0} \leq k
\end{array} & (1) \\
& (2)
\end{array}\right\}
$$

where the constraint (1) is the so-called Stiefel manifold (denoted as $\operatorname{St}(d, r)$ ) and the constraint (2) is the row-sparsity constraint. For the constraint (1), recall that the convex-hull of Stiefel manifold can be explicitly represented (?) as,

$$
\operatorname{Conv}(\operatorname{St}(d, r)):=\left\{\boldsymbol{V}:\left(\begin{array}{cc}
\boldsymbol{I}_{d} & -\boldsymbol{V} \\
-\boldsymbol{V}^{\top} & \boldsymbol{I}_{r}
\end{array}\right) \succeq \mathbf{0}_{d+r}\right\} \Leftrightarrow\left\{\boldsymbol{V}: \boldsymbol{I}_{r}-\boldsymbol{V}^{\top} \boldsymbol{V} \succeq \mathbf{0}_{r}\right\} \Leftrightarrow\left\{\boldsymbol{V}:\|\boldsymbol{V}\|_{\mathrm{op}} \leq 1\right\}
$$

For the constraint (2),
Proposition 2.1. If $\boldsymbol{V} \in \mathcal{F}$, then $\left\|[\boldsymbol{V}]_{[d], i}\right\|_{1} \leq \sqrt{k}$ holds for all $i \in[r]$.
Proof. Since the operator norm of $\boldsymbol{V}$ is upper bounded by 1 , we have that the $\ell_{2}$-norm of each column of $\boldsymbol{V}$ is at most 1 . Moreover, each column is $k$-sparse, then we have $\left\|[\boldsymbol{V}]_{[d], i}\right\|_{1} \leq \sqrt{k}$ holds for all $i \in[r]$.

[^0]The above proposition can be viewed as the $\ell_{1}$-relaxation of the sparsity constraint for each column in $\boldsymbol{V}$. Moreover, the row-sparsity property can be futher captured by
Proposition 2.2. If $\boldsymbol{V} \in \mathcal{F}$, then $\sum_{j=1}^{d}\left\|[\boldsymbol{V}]_{j,[r]}\right\|_{2} \leq \sqrt{r k}$.
Proof. For any $\boldsymbol{V} \in \mathcal{F}$, based on the row-sparsity condition $\|\boldsymbol{V}\|_{0} \leq k$, there are at most $k$ non-zero values of $\left\|[\boldsymbol{V}]_{j,[r]}\right\|_{2}$ among $j \in[d]$. Since $\|\boldsymbol{V}\|_{\text {op }} \leq 1$ with rank $r$, then

$$
r \geq\|\boldsymbol{V}\|_{F}^{2}=\sum_{j=1}^{d}\left\|[\boldsymbol{V}]_{j,[r]}\right\|_{2}^{2}
$$

which implies that $\sum_{j=1}^{d}\left\|[\boldsymbol{V}]_{j,[r]}\right\|_{2} \leq \sqrt{r k}$.
From Proposition 2.1 and 2.2, we obtain the following result.
Corollary 2.1. [SDP-relaxation] Let $\mathcal{F}$ be the feasible region of SPCA. We have $\operatorname{conv}(\mathcal{F})$ is contained in the following convex set

$$
\mathcal{C}:=\left\{\begin{array}{ll} 
& \boldsymbol{I}_{r}-\boldsymbol{V}^{\top} \boldsymbol{V} \succeq \mathbf{0}_{r}, \\
\boldsymbol{V}: & \sum_{j=1}^{d}\left\|\boldsymbol{v}_{i}\right\|_{1} \leq \sqrt{k}, \forall i \in[r] \\
& \sum_{j=1}^{d}\left\|[\boldsymbol{V}]_{j,[r]}\right\|_{2} \leq \sqrt{r k}, \forall j \in[d]
\end{array}\right\}
$$

Since SDP-relaxations are usually difficult to solve, to be more scalable in practice, instead of using semi-definite constraint, we replace it with second-order cone constraints. In particular, we will replace the constraints defining the convex hull of the Stiefel manifold by a simple second-order-cone representable relaxation to obtain the following result.
Corollary 2.2. [SOCP-relaxation] Let $\mathcal{F}$ be the feasible region of SPCA. We have $\operatorname{conv}(\mathcal{F})$ is contained in the following convex set

$$
\mathcal{C}^{\prime}:=\left\{\begin{array}{ll} 
& \left\|[\boldsymbol{V}]_{[d], i}\right\|_{2}^{2} \leq 1, \forall i \in[r] \\
\boldsymbol{V}: & \left\|[\boldsymbol{V}]_{[d], i_{1}} \pm[\boldsymbol{V}]_{[d], i_{2}}\right\|_{2}^{2} \leq 2, \forall i_{1} \neq i_{2} \in[r] \\
& \sum_{j=1}^{d}\left\|[\boldsymbol{V}]_{j,[r]}\right\|_{2} \leq \sqrt{r k} \\
& \left\|[\boldsymbol{V}]_{j,[r]}\right\|_{2} \in[0,1], \forall j \in[d]
\end{array}\right\}
$$

### 2.1. Proof of Theorem 1

Let $\mathrm{opt}^{\mathcal{F}}, \mathrm{opt}^{\mathcal{C}}, \mathrm{opt}^{\mathcal{C}^{\prime}}$ be the optimal values of the following:

$$
\begin{align*}
& \mathrm{opt}^{\mathcal{F}}:=\max _{\boldsymbol{V} \in \mathcal{F}} \operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{A} \boldsymbol{V}\right), \\
& \mathrm{opt}^{\mathcal{C}}:=\max _{\boldsymbol{V} \in \mathcal{C}} \operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{A} \boldsymbol{V}\right),  \tag{Relax}\\
& \mathrm{opt}^{\mathcal{C}^{\prime}}:=\max _{\boldsymbol{V} \in \mathcal{C}^{\prime}} \operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{A} \boldsymbol{V}\right)
\end{align*}
$$

(SOCP-Relax)
Our first main result is that:
Theorem 1. $\mathrm{opt}^{\mathcal{F}} \leq \mathrm{opt}^{\mathcal{C}^{\prime}} \leq(1+\sqrt{r})^{2} \mathrm{opt}^{\mathcal{F}}$.
Proof. Consider any $\boldsymbol{V} \in \mathcal{C}^{\prime}$, sort the row of $\boldsymbol{V}$ by its $\ell_{2}$-norm in decreasing order $\left\{j_{1}, j_{2}, \ldots, j_{d}\right\}$, i.e.,

$$
\left\|[\boldsymbol{V}]_{j_{1},[r]}\right\|_{2} \geq \ldots \geq\left\|[\boldsymbol{V}]_{j_{d},[r]}\right\|_{2}
$$

Decompose the matrix $\boldsymbol{V}$ based on its top- $k$ largest rows, second top- $k$ largest rows, and so on, i.e., let $m=\lceil d / k\rceil$, $\boldsymbol{V}=\boldsymbol{V}^{1}+\cdots \boldsymbol{V}^{m}$ with

$$
\operatorname{supp}\left(\boldsymbol{V}^{1}\right)=\left\{j_{1}, \ldots, j_{k}\right\}=: J^{1}, \ldots, \boldsymbol{V}^{m}=\left\{j_{d-(m-1) k}, \ldots, j_{d}\right\}=: J^{m}
$$

For each $p=1, \ldots, m$, we have $\left\|\boldsymbol{V}^{p} /\right\| \boldsymbol{V}^{p}\left\|_{\mathrm{op}}\right\|_{0} \leq k,\left\|\boldsymbol{V}^{p} /\right\| \boldsymbol{V}^{p}\left\|_{\mathrm{op}}\right\|_{\mathrm{op}}=1$, thus $\boldsymbol{V}^{p} /\left\|\boldsymbol{V}^{p}\right\|_{\mathrm{op}} \in \operatorname{conv}(\mathcal{F})$. Since $\operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{A} \boldsymbol{V}\right)$ is convex, then $\max _{\boldsymbol{V} \in \mathcal{F}} \operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{A} \boldsymbol{V}\right)=\max _{\boldsymbol{V} \in \operatorname{conv}(\mathcal{F})} \operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{A} \boldsymbol{V}\right)$. To verify the approximation ratio,

$$
\begin{aligned}
& \boldsymbol{V}=\boldsymbol{V}^{1}+\cdots \boldsymbol{V}^{m}=\left\|\boldsymbol{V}^{1}\right\|_{\mathrm{op}} \frac{\boldsymbol{V}^{1}}{\left\|\boldsymbol{V}^{1}\right\|_{\mathrm{op}}}+\cdots+\left\|\boldsymbol{V}^{m}\right\|_{\mathrm{op}} \frac{\boldsymbol{V}^{m}}{\left\|\boldsymbol{V}^{m}\right\|_{\mathrm{op}}} \\
\Leftrightarrow & \frac{\boldsymbol{V}}{\sum_{p=1}^{m}\left\|\boldsymbol{V}^{p}\right\|_{\mathrm{op}}}=\frac{\left\|\boldsymbol{V}^{1}\right\|_{\mathrm{op}}}{\sum_{p=1}^{m}\left\|\boldsymbol{V}^{p}\right\|_{\mathrm{op}}} \frac{\boldsymbol{V}^{1}}{\left\|\boldsymbol{V}^{1}\right\|_{\mathrm{op}}}+\cdots+\frac{\left\|\boldsymbol{V}^{m}\right\|_{\mathrm{op}}}{\sum_{p=1}^{m}\left\|\boldsymbol{V}^{p}\right\|_{\mathrm{op}}} \frac{\boldsymbol{V}^{m}}{\left\|\boldsymbol{V}^{m}\right\|_{\mathrm{op}}} \in \operatorname{conv}(\mathcal{F}) .
\end{aligned}
$$

Notice that

$$
\left\|\boldsymbol{V}^{p}\right\|_{\mathrm{op}} \leq \max \left\{1, \sqrt{\sum_{j \in J^{p}}\left\|[\boldsymbol{V}]_{j,[r]}\right\|_{2}^{2}}\right\}
$$

then based on the decomposition of $\ell_{2}$ norms of rows,

$$
\begin{aligned}
\sum_{p=1}^{m}\left\|\boldsymbol{V}^{p}\right\|_{\mathrm{op}} & =\left\|\boldsymbol{V}^{1}\right\|_{\mathrm{op}}+\sum_{p=2}^{m}\left\|\boldsymbol{V}^{p}\right\|_{\mathrm{op}} \\
& \leq 1+\sum_{p=2}^{m} \sqrt{\left(\frac{\sum_{j \in J^{p-1}}\left\|[\boldsymbol{V}]_{j,[r]}\right\|_{2}}{k}\right)^{2} \cdot k} \\
& \leq 1+\frac{1}{\sqrt{k}} \cdot \sum_{p=2}^{m} \sum_{j \in J^{p-1}}\left\|[\boldsymbol{V}]_{j,[r]}\right\|_{2} \\
& \leq 1+\frac{1}{\sqrt{k}} \sum_{j=1}^{d}\left\|[\boldsymbol{V}]_{j,[r]}\right\|_{2} \\
& \leq 1+\sqrt{r}
\end{aligned}
$$

where the final inequality holds since the constraint $\sum_{j=1}^{d}\left\|[\boldsymbol{V}]_{j,[r]}\right\|_{2} \leq \sqrt{r k}$ in $\mathcal{C}^{\prime}$. Therefore, we have

$$
\boldsymbol{V} \in\left(\sum_{p=1}^{m}\left\|\boldsymbol{V}^{p}\right\|_{\mathrm{op}}\right) \cdot \operatorname{conv}(\mathcal{F}) \subseteq(1+\sqrt{r}) \cdot \operatorname{conv}(\mathcal{F})
$$

i.e., $\mathcal{C}^{\prime} \subseteq(1+\sqrt{r}) \cdot \operatorname{conv}(\mathcal{F})$. Hence opt ${ }^{\mathcal{F}} \leq \mathrm{opt}^{\mathcal{C}^{\prime}} \leq(1+\sqrt{r})^{2} \mathrm{opt}^{\mathcal{F}}$ holds.

A corollary can be derived from the Theorem 1 based on the containment $\mathcal{C} \subseteq \mathcal{C}^{\prime}$ as follows:
Corollary 2.3. $\mathrm{opt}^{\mathcal{F}} \leq \mathrm{opt}^{\mathcal{C}} \leq(1+\sqrt{r})^{2} \mathrm{opt}^{\mathcal{F}}$.
Remark: For $r=1$ case, Theorem 1 and Corollary 2.3 provide constant multiplicative approximation ratios. Thus inapproximability results from (??) implies that solving Relax or SOCP-Relax to optimality is NP-hard.

### 2.2. Proof of Proposition 2.3 \& 2.4

To overcome the non-convex part of Relax or SOCP-Relax, the objective function of Relax or SOCP-Relax is further relaxed via piecewise-linear functions using special-ordered sets type-2 constraints. Recall that the piecewise-linear approximation (PLA) set is defined as
in which SOS-II denotes the set of special-ordered sets type-2 constraints as follows:

$$
\text { SOS-II }:=\left\{\begin{array}{lll} 
& \sum_{\ell=-N}^{N} \eta_{j i}^{\ell}=1 & \\
& \sum_{\ell=-N}^{N-1} y^{\ell}=1 & \\
\left(\eta_{j i}^{\ell}\right)_{\ell=-N}^{N}: & \eta_{j i}^{\ell}+\eta_{j i}^{\ell+1} \leq y^{\ell} & \ell=-N, \ldots, N-1 \\
& \eta_{j i}^{\ell} \geq 0 & \ell=-N, \ldots, N \\
& y^{\ell} \in\{0,1\} & \ell=-N, \ldots, N-1
\end{array}\right\}
$$

since SOS-II contains the integer variables, we name the convex relaxation of Relax or SOCP-Relax be semi-definite convex integer program (SDCIP)

$$
\begin{aligned}
\mathrm{ub}^{\mathcal{C}}:=\max & \sum_{j=1}^{d} \lambda_{j} \sum_{i=1}^{r} \xi_{j i} \\
\text { s.t. } & \boldsymbol{V} \in \mathcal{C} \\
& (g, \xi, \eta) \in \mathrm{PLA}
\end{aligned}
$$

(SDCIP)
or second-order-cone convex integer program (SOCIP)

$$
\begin{aligned}
\mathrm{ub}^{\mathcal{C}^{\prime}}:=\max & \sum_{j=1}^{d} \lambda_{j} \sum_{i=1}^{r} \xi_{j i} \\
\text { s.t. } & \boldsymbol{V} \in \mathcal{C}^{\prime} \\
& (g, \xi, \eta) \in \text { PLA }
\end{aligned}
$$

(SOCIP)

Thus we have the Proposition 2.3 and Proposition 2.4,
Proposition 2.3. The optimal value $u b^{\mathcal{C}^{\prime}}$ of SOCIP is an upper bound of SPCA.
Proof. Based on Theorem 2.3, we have the optimal value opt ${ }^{\mathcal{F}}$ of SPCA

$$
\text { opt }^{\mathcal{F}}:=\max _{\boldsymbol{V}^{\top} \boldsymbol{V}=\boldsymbol{I}_{r},\|\boldsymbol{V}\|_{0} \leq k} \operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{A} \boldsymbol{V}\right)
$$

is upper bounded by the optimal value opt ${ }^{\mathcal{C}^{\prime}}$ of Relax

$$
\mathrm{opt}^{\mathcal{C}}:=\max _{\boldsymbol{V} \in \mathcal{C}^{\prime}} \operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{A} \boldsymbol{V}\right)
$$

To show that $\mathrm{ub}^{\mathcal{C}^{\prime}}$ is an upper bound, it is sufficient to show that $\mathrm{ub}^{\mathcal{C}^{\prime}} \geq \mathrm{opt}^{\mathcal{C}^{\prime}}$. Consider the auxiliary variable $\xi_{j i}$ for all $(j, i) \in[d] \times[r]$ with $g_{j i}=\boldsymbol{a}_{j}^{\top} \boldsymbol{v}_{i}$. Based on the property of SOS-II constraint, for a fixed $(j, i)$, there are at most two non-zero continuous variables in $\eta_{j i}^{\ell}$, say $\eta_{j i}^{\ell^{*}}, \eta_{j i}^{\ell^{*}+1}$, such that $\eta_{j i}^{\ell^{*}}+\eta_{j i}^{\ell^{*}+1}=1$. Combining with the constraints in set PLA, we have

$$
g_{j i}=\gamma_{j i}^{\ell^{*}} \eta_{j i}^{\ell^{*}}+\gamma_{j i}^{\ell^{*}+1} \eta_{j i}^{\ell^{*}+1}, \quad \quad \xi_{j i}=\left(\gamma_{j i}^{\ell^{*}}\right)^{2} \eta_{j i}^{\ell^{*}}+\left(\gamma_{j i}^{\ell^{*}+1}\right)^{2} \eta_{j i}^{\ell^{*}+1}
$$

and hence,

$$
\xi_{j i}-g_{j i}^{2}=\left(\gamma_{j i}^{\ell^{*}+1}-\gamma_{j i}^{\ell^{*}}\right)^{2} \eta_{j i}^{\ell_{i}^{*}} \eta_{j i}^{\ell^{*}+1} \geq 0
$$

That is to say, the objective function $\sum_{j=1}^{d} \lambda_{j} \sum_{i=1}^{r} \xi_{j i}$ in SOCIP is greater than or equal to the function $\sum_{j=1}^{d} \lambda_{j} \sum_{i=1}^{r} g_{j i}^{2}$, where $\sum_{\mathcal{C}^{\prime}=1}^{d} \lambda_{j} \sum_{i=1}^{r} g_{j i}^{2}$ is equivalent to the objective function $\operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{A} \boldsymbol{V}\right)$ in Relax by the definition of $g_{j i}$. Therefore, $\mathrm{ub}^{\mathcal{C}^{\prime}} \geq \mathrm{opt}^{\mathcal{C}^{\prime}}$ holds.
Proposition 2.4. The optimal value $\mathrm{ub}^{\mathcal{C}^{\prime}}$ of SOCIP can be upper bounded by $(1+\sqrt{r})^{2} \mathrm{opt}^{\mathcal{F}}+\sum_{j=1}^{d} \frac{r \lambda_{j} \theta_{j}^{2}}{4 N^{2}}$, which is an affine function of opt ${ }^{\mathcal{F}}$.

Proof. For SOCIP, the objective function $\operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{A} \boldsymbol{V}\right)$ equals to $\sum_{j=1}^{d} \lambda_{j} \sum_{i=1}^{r}\left(\boldsymbol{a}_{j}^{\top} \boldsymbol{v}_{i}\right)^{2}=\sum_{j=1}^{d} \lambda_{j} \sum_{i=1}^{r} g_{j i}^{2}$. By Theorem 1, we have $\sum_{j=1}^{d} \lambda_{j} \sum_{i=1}^{r} g_{j i}^{2} \leq 4 r$ opt $^{\mathcal{F}}$. Since $g_{j i} \in\left[-\theta_{j}, \theta_{j}\right]$, without any prior information, we split the
interval $\left[-\theta_{j}, \theta_{j}\right]$ evenly via splitting points $\left(\gamma_{j i}^{\ell}\right)_{\ell=-N}^{N}$ such that $\gamma_{j i}^{\ell}=\frac{\ell}{N} \cdot \theta_{j}$. Based on the proof of Proposition 2.1, we have

$$
\xi_{j i}-g_{j i}^{2}=\left(\gamma_{j i}^{\ell^{*}+1}-\gamma_{j i}^{\ell^{*}}\right)^{2} \eta_{j i}^{\ell^{*}} \eta_{j i}^{\ell^{*}+1}=\frac{\theta_{j}^{2}}{N^{2}} \eta_{j i}^{\ell^{*}} \eta_{j i}^{\ell^{*}+1} \leq \frac{\theta_{j}^{2}}{4 N^{2}}
$$

Therefore, the objective function in SOCIP $\sum_{j=1}^{d} \lambda_{j} \sum_{i=1}^{r} \xi_{j i} \leq \sum_{j=1}^{d} \lambda_{j} \sum_{i=1}^{r} g_{j i}^{2}+\sum_{j=1}^{d} \frac{r \lambda_{j} \theta_{j}^{2}}{4 N^{2}} \leq(1+\sqrt{r})^{2}$ opt $^{\mathcal{F}}+$ $\sum_{j=1}^{d} \frac{r \lambda_{j} \theta_{j}^{2}}{4 N^{2}}$.

## 3. Proof of Primal (Lower) Bounds

Recall that Given a sample covariance matrix $\boldsymbol{A}$, let $\boldsymbol{A}^{1 / 2}$ be its positive semi-definite square root such that $\boldsymbol{A}=$ $\boldsymbol{A}^{1 / 2} \boldsymbol{A}^{1 / 2}$, the SPCA can be represented in the following fashion:

$$
\begin{aligned}
\min _{\boldsymbol{V} \in \mathbb{R}^{d \times r}} & \left\|\boldsymbol{A}^{1 / 2}-\boldsymbol{V} \boldsymbol{V}^{\top} \boldsymbol{A}^{1 / 2}\right\|_{F}^{2} \\
\text { s.t. } & \boldsymbol{V}^{\top} \boldsymbol{V}=\boldsymbol{I}_{r} \\
& \|\boldsymbol{V}\|_{0} \leq k
\end{aligned}
$$

(SPCA-lasso)
where $\|\cdot\|_{F}$ denotes the Frobenius norm. Furthermore, SPCA-lasso can be reformulated into the following two-stage (inner \& outer) optimization problem:

$$
\underbrace{\min _{S \subseteq[d],|S| \leq k}}_{\text {outer optimization }} \underbrace{\min _{[\boldsymbol{V}]_{S}^{\top}[\boldsymbol{V}]_{S}=\boldsymbol{I}_{r}}}_{\text {inner optimization }}\left\|\left[\boldsymbol{A}^{1 / 2}\right]_{S}-[\boldsymbol{V}]_{S}[\boldsymbol{V}]_{S}^{\top}\left[\boldsymbol{A}^{1 / 2}\right]_{S}\right\|_{F}^{2}+\left\|\left[\boldsymbol{A}^{1 / 2}\right]_{S^{C}}\right\|_{F}^{2}
$$

Given support $S$, there is a closed form solution of the inner optimization by eigenvalue decomposition. Thus the main challenging of solving SPCA-lasso is to find a support set $S$ within $\binom{d}{k}$ possible support set, which is known to be NP-hard. The local search algorithm 1 is therefore proposed to find a relative good primal solution (via updating the support set $S$ in each epoch).

```
Algorithm 1 Local Search Method
Input: Covariance matrix \(\boldsymbol{A}\), sparsity parameter \(k\), number of eigenvectors \(r\), number of maximum iterations \(T\).
Output: A feasible solution \(\boldsymbol{V}\) for SPCA.
    Initialize with \(S_{0} \subseteq[d]\) with \(\left|S_{0}\right|=k\).
    while epoch \(t=1, \ldots, T\) do
        For each \(j \in S_{t-1}\), set the reduced value \(\Delta_{j}\), set removing candidate \(j^{\text {out }}:=\arg \min _{j \in S_{t-1}} \Delta_{j}\).
        For each \(j^{\prime} \in S_{t-1}^{C}\), set the will-reduced value \(\Delta_{j^{\prime}}\), set adding candidate \(j^{\text {in }}:=\arg \min _{j^{\prime} \in S_{t-1}} \Delta_{j^{\prime}}\)
        if \(\Delta_{j \text { in }}>\Delta_{j \text { out }}\) then
            Set \(S_{t}:=S_{t-1}-\left\{j^{\text {out }}\right\}+\left\{j^{\text {in }}\right\}\).
            By eigenvalue decomposition, \(\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t}}\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t}}^{\top}=\boldsymbol{U}_{S_{t}} \boldsymbol{\Lambda}_{S_{t}} \boldsymbol{U}_{S_{t}}^{\top}\), set \([\boldsymbol{V}]_{S_{t}}=\left[\boldsymbol{U}_{S_{t}}\right]_{[k],[r]}\).
        else
            Break while loop.
        end if
    end while
    Return \(V\).
```

Remark 3.1. Note that, for general instance (i.e., model-free case), we initialize the support set $S_{0}$ by picking $S_{0} \subseteq[d]$ uniformly at random, and repeat the Algorithm 1 for serval times for best solution. But in Section 3.2, we show that under some statistical assumptions, a specific initialization method will find the support set with respect to the optimal solution with high probability.

### 3.1. Primal Method in Model-free Case

Here we arrive the following results: In model-free case,

Theorem 2. Algorithm 1 is a monotone decreasing algorithm in the objective value of SPCA.

Proof. Based on the updating rule in Algorithm 1, when $\Delta_{j \text { in }}>\Delta_{j \text { out }}$ and set $S_{t}=S_{t-1}-\left\{j^{\text {out }}\right\}+\left\{j^{\text {in }}\right\}$, the objective value of SPCA-lasso satisfies

$$
\begin{aligned}
\left\|\boldsymbol{A}^{1 / 2}-\boldsymbol{V} \boldsymbol{V}^{\top} \boldsymbol{A}^{1 / 2}\right\|_{F}^{2} & =\left\|\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t-1}}-[\boldsymbol{V}]_{S_{t-1}}[\boldsymbol{V}]_{S_{t-1}}^{\top}\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t-1}}\right\|_{F}^{2}+\left\|\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t-1}^{C}}\right\|_{F}^{2} \\
& >\left\|\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t}}-[\boldsymbol{V}]_{S_{t-1}}[\boldsymbol{V}]_{S_{t-1}}^{\top}\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t}}\right\|_{F}^{2}+\left\|\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t}^{C}}\right\|_{F}^{2} \\
& \geq\left\|\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t}}-[\boldsymbol{V}]_{S_{t}}[\boldsymbol{V}]_{S_{t}}^{\top}\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t}}\right\|_{F}^{2}+\left\|\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t}^{C}}\right\|_{F}^{2}
\end{aligned}
$$

Thus Algorithm 1 is a monotone decreasing algorithm.
Theorem 3. Algorithm 1 terminates in at most $\binom{d}{k}$ epochs.
Proof. We claim that if a subset $S \subseteq[d]$ exists in $t$ th epoch, then such set $S$ will not exist in future epochs. Otherwise, suppose $S_{t_{1}}=S_{t_{2}}$ and $t_{1}<t_{2}$, then based on the monotonicity of Algorithm 1, we have

$$
\begin{aligned}
& \left\|\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t_{1}}}-[\boldsymbol{V}]_{S_{t_{1}}}[\boldsymbol{V}]_{S_{t_{1}}}^{\top}\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t_{1}}}\right\|_{F}^{2}+\left\|\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t_{1}}}\right\|_{F}^{2} \\
> & \left\|\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t_{2}}}-[\boldsymbol{V}]_{S_{t_{2}}}[\boldsymbol{V}]_{S_{t_{2}}}^{\top}\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t_{2}}}\right\|_{F}^{2}+\left\|\left[\boldsymbol{A}^{1 / 2}\right]_{S_{t_{2}}}\right\|_{F}^{2}
\end{aligned}
$$

which contradicts to $S_{t_{1}}=S_{t_{2}}$, and the claim holds. Since there are at most $\binom{d}{k}$ possible subsets, then the Algorithm 1 terminates in finite epochs.

### 3.2. Primal Method with Statistical Model

Although this is not the main contribution of our paper, to complete the framework, we demonstrate that the primal feasible solution obtained from our primal heuristic ensures some statistical properties when the following assumptions hold.
Assumption 1. Assume that $\left\{\boldsymbol{x}_{m}\right\}_{m=1}^{M} \in \mathbb{R}^{d}$ is a sequence of i.i.d. random samples generated from Gaussian distribution with zero mean and true covariance matrix $\boldsymbol{\Sigma}$.

To demonstrate the main difference between usual sparse PCA and row-sparse PCA defined in SPCA, consider the following block spiked covariance matrix model (described in Assumption 2) and support-inconsistency condition (Assumption 3).
Assumption 2. Let $\Theta$ be the collection of all true covariance matrix $\boldsymbol{\Sigma}$ with block diagonal structure $\boldsymbol{\Sigma}=\boldsymbol{\Sigma}_{1} \oplus \boldsymbol{\Sigma}_{2} \oplus \mathbf{0}_{\text {rest }}$ such that:

- $\boldsymbol{\Sigma}_{1} \in \mathbb{R}^{k \times k}$ is a rank $r^{*}(<k)$ matrix with top- $r^{*}$ largest eigen-pairs $\left(\lambda_{1}, \boldsymbol{v}_{1}\right), \ldots,\left(\lambda_{r^{*}}, \boldsymbol{v}_{r^{*}}\right), \boldsymbol{\Sigma}_{2} \in \mathbb{R}^{k \times k}$ is a fullrank matrix with the next top- $k$ largest eigen-pairs $\left(\lambda_{r^{*}+1}, \boldsymbol{v}_{r^{*}+1}\right), \ldots,\left(\lambda_{r^{*}+k}, \boldsymbol{v}_{r^{*}+k}\right)$, in which $\lambda_{1}>\cdots>\lambda_{r^{*}}>$ $\lambda_{r^{*}+1}>\cdots>\lambda_{r^{*}+k}>0$.
- The diagonal entries of $\boldsymbol{\Sigma}$ satisfy $\max _{j \in[k]}\left[\boldsymbol{\Sigma}_{1}\right]_{j, j}<\min _{j \in[k]}\left[\boldsymbol{\Sigma}_{2}\right]_{j, j}$, i.e., the top- $k$ diagonal entries are all in submatrix $\boldsymbol{\Sigma}_{2}$.

Let $\Sigma$ be a true covariance matrix in set $\Theta$.
Assumption 3. Assume that $\boldsymbol{\Sigma}$ satisfies: the sum of eigenvalues in first block $\boldsymbol{\Sigma}_{1}$ is upper bounded by $\sum_{j=r^{*}+1}^{2 r^{*}+1} \lambda_{j}$, i.e.,

$$
\sum_{j=1}^{r^{*}} \lambda_{j}<\sum_{j=r^{*}+1}^{2 r^{*}+1} \lambda_{j}
$$

Remark 3.2. Based on Assumption 2 and 3, easy to observe that:

- given $k$ and $r \leq r^{*}$, the global optimal solution of SPCA with truth covariance matrix $\boldsymbol{\Sigma}$ has support set $\{1, \ldots, k\}$,
- in contrast, given $k$ and $r>r^{*}$, the global optimal solution of SPCA with truth covariance matrix $\boldsymbol{\Sigma}$ has support set $\{k+1, \ldots, 2 k\}$.

Thus a significant shortage of existing support recovery methods is that these methods fail to recover the optimal support set of SPCA.

To overcome this shortage, we proposed an initialization method that recovers the optimal support set with two procedures:

- Recover the union of support sets: Based on the existing results in ??, the Covariance Thresholding algorithm (Algorithm 1 in ?) with input sample covariance matrix $\boldsymbol{A}=\frac{1}{M} \sum_{m=1}^{M} \boldsymbol{x}_{m} \boldsymbol{x}_{m}^{\top}$, sparse parameter $2 k$, thresholding parameter $\tau, \rho$ (defined in Theorem 1, Theorem 3 of ?) is able to recover the union of support sets

$$
S:=\bigcup_{i=1}^{r^{*}+k} \operatorname{supp}\left(\boldsymbol{v}_{i}\right)
$$

using the soft-thresholding matrix $\hat{\boldsymbol{A}}$ (mentioned in Algorithm 1 in ?) with high probability, if the following Assumption 4 is satisfied.
Assumption 4. Let $\boldsymbol{v}_{1}, \ldots, \boldsymbol{v}_{d}$ be the set of eigenvectors of $\boldsymbol{\Sigma}$ corresponding to eigenvalues $\lambda_{1}, \ldots, \lambda_{d}$ respectively as we defined in Assumption 2. Let $S_{1}, \ldots, S_{d}$ be the set of supports of $\boldsymbol{v}_{1}, \ldots, \boldsymbol{v}_{d}$. There exists constants $c_{1}, c_{2}>0$ such that the following holds. The non-zero components satisfy $\left|\left[\boldsymbol{v}_{j}\right]_{i}\right| \geq c_{1} / \sqrt{k}$ for all $i \in S_{j}$ and $j=1, \ldots, r^{*}+k$. Furthermore, for any $j, j^{\prime},\left|\left[\boldsymbol{v}_{j}\right]_{i} /\left[\boldsymbol{v}_{j^{\prime}}\right]_{i}\right| \leq c_{2}$ for all $i \in S_{j} \cap S_{j^{\prime}}$. Here we assume that $c_{2} \geq 1$.

In the block spiked covariance matrix assumption, the union of support sets is the set with respect to the blocks $\Sigma_{1}, \boldsymbol{\Sigma}_{2}$.

- Search for optimal support for SPCA: In the second procedure, we search for the global optimal support set of SPCA problem based the structure of blocked spiked covariance matrix.

Here is the pseudocode of Initialization method 2.

```
Algorithm 2 Initialization - Statistical Model
Input: Sample covariance matrix \(\boldsymbol{A}\), sparsity parameter \(k\), number of eigenvectors \(r\).
Output: An initial support set \(S_{0}\).
Do Covariance Thresholding Algorithm with input \((\boldsymbol{A}, 2 k, \tau, \rho)\) described in ?, let \(S\) be the union of support sets, and \(\hat{\boldsymbol{A}}\) be the soft-thresholding matrix of \(\boldsymbol{A}\).
Set \(\overline{\boldsymbol{A}}=[\hat{\boldsymbol{A}}]_{S, S}\). Without loss of generality, let \(\{k+1, \ldots, 2 k\}\) be the set of indices corresponding to the top- \(k\) diagonal entries in \(\bar{A}\), and let \(\{1, \ldots, k\}\) be the rest.
Return support set \(S_{0}=\{1, \ldots, k\}\) if \(r \leq r^{*}\), and \(S_{0}=\{k+1, \ldots, 2 k\}\) if \(r>r^{*}\).
```

Notice that Assumptions 1, 2, 3, 4 can hold simultaneously, one possible example is the artificial instance mentioned in Section 3 of the main content.

Let $\beta:=\min _{q \neq q^{\prime}}\left\{\lambda_{1}-\lambda_{2}, \lambda_{2}-\lambda_{3}, \ldots, \lambda_{r^{*}+k-1}-\lambda_{r^{*}+k}, \lambda_{r^{*}+k}\right\}$ be the minimum eigenvalue gap.
Theorem 4. Assume the assumptions 1, 2, 3, 4 hold, and further we have $M \geq \frac{C^{2}}{\epsilon^{2}}(2 k)^{2} \max \{\beta, 1\} \log \frac{d}{(2 k)^{2}}, M>$ $(2 k)^{2}, \frac{d}{e}>(2 k)^{2}$ for $C$ a constant. If

$$
\max _{j \in[k]}\left[\boldsymbol{\Sigma}_{1}\right]_{j, j}+\epsilon<\min _{j \in[k]}\left[\boldsymbol{\Sigma}_{2}\right]_{j, j}
$$

then the solution $\boldsymbol{V} \in \mathbb{R}^{d \times r}$ obtained from Algorithm 1 satisfies

$$
\operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{\Sigma} \boldsymbol{V}\right) \geq \max _{\boldsymbol{V} \in \mathcal{F}} \operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{\Sigma} \boldsymbol{V}\right)-2 r \epsilon
$$

with high probability $1-o(1)$.

Proof. We show this in the following steps:

Estimating the sample covariance matrix $A$ : Recall the Theorem 1 and Covariance Thresholding Algorithm in ?, there exists a constant $C>0$ such that: for any $\epsilon>0$, when the number of samples $M \geq \frac{C^{2}}{\epsilon^{2}}(2 k)^{2} \max \{\beta, 1\} \log \frac{d}{(2 k)^{2}}$, the soft-thresholding matrix $\hat{\boldsymbol{A}}$ (mentioned in Covariance Thresholding Algorithm) satisfies $\|\hat{\boldsymbol{A}}-\boldsymbol{\Sigma}\|_{\text {op }} \leq \epsilon$ with probability $1-o(1)$. This implies that for any diagonal entries $\hat{\boldsymbol{A}}_{j j}$ with $j=1, \ldots, d$, we have $\left|\hat{\boldsymbol{A}}_{j j}-\hat{\boldsymbol{\Sigma}}_{j j}\right|<\epsilon$ holds with probability $1-o(1)$.

Recovering the union of supports: Since the Assumptions 1, 2, 3, 4 hold, and the sample size $M$ satisfies the conditions given in 4, Theorem 3 in (?) guarantees that: there exists a constant $C_{0}=C_{0}\left(c_{1}, c_{2}, \lambda_{1}, \beta\right)$ such that if $M \geq$ $C_{0}(2 k)^{2}\left(k+r^{*}\right) \log \frac{d}{(2 k)^{2}}$, then the covariance thresholding algorithm recovers the union of supports $\bigcup_{j=1}^{r^{*}+k} S_{j}$ with probability $1-o(1)$.

Recovering the exact support of SPCA: Note that the set $S$ obtained from the Covariance Thresholding Algorithm only recovers the union of supports $\bigcup_{j=1}^{r^{*}+k} S_{j}$ with high probability. Condition on $S$ is successfully recovered, suppose $\max _{j \in[k]}\left[\boldsymbol{\Sigma}_{1}\right]_{j, j}+2 \epsilon<\min _{j \in[k]}\left[\boldsymbol{\Sigma}_{2}\right]_{j, j}$, we have the block indexed by the top- $k$ diagonal entries corresponding to the block $\boldsymbol{\Sigma}_{2}$ since $\left|\hat{\boldsymbol{A}}_{j j}-\hat{\boldsymbol{\Sigma}}_{j j}\right|<\epsilon$ holds with probability $1-o(1)$. Hence the initialization method 2 recover the true support by $S_{0}$.
Quality of primal feasible solution: Let $\boldsymbol{V}$ be the primal solution obtained from Algorithm 1. Let $\boldsymbol{A}_{S_{0}, S_{0}}=$ $\boldsymbol{U}_{S_{0}} \boldsymbol{\Lambda}_{S_{0}} \boldsymbol{U}_{S_{0}}^{\top} \in \mathbb{R}^{k \times k}$ be the eigenvalue decomposition of $\boldsymbol{A}_{S_{0}, S_{0}}$ with diagonal entries in $\boldsymbol{\Lambda}_{S_{0}}$ sorted in decreasing order. We have the solution obtained by Algorithm 1 satisfies:

$$
\begin{array}{rlr}
\operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{\Sigma} \boldsymbol{V}\right) & \geq \operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{A} \boldsymbol{V}\right)-\left|\operatorname{Tr}\left(\boldsymbol{V}^{\top}(\boldsymbol{\Sigma}-\boldsymbol{A}) \boldsymbol{V}\right)\right| & \text { by }\|\hat{\boldsymbol{A}}-\boldsymbol{\Sigma}\|_{\mathrm{op}} \leq \epsilon \\
& \geq \operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{A} \boldsymbol{V}\right)-r \epsilon & \\
& \geq \operatorname{Tr}\left(\left[\boldsymbol{U}_{S_{0}}\right]_{[k],[r]}^{\top} \boldsymbol{A}_{S_{0}, S_{0}}\left[\boldsymbol{U}_{S_{0}}\right]_{[k],[r]}\right)-r \epsilon & \text { by monotone local search 1 with initial }\left[\boldsymbol{U}_{S_{0}}\right]_{[k],[r]} \\
& =\sum_{i=1}^{r} \lambda_{i}\left(\boldsymbol{A}_{S_{0}, S_{0}}\right)-r \epsilon & \text { let } \lambda_{i}\left(\boldsymbol{A}_{S_{0}, S_{0}}\right) \text { be } i \text {-th largest eigenvalue } \\
& \geq \sum_{i=1}^{r} \lambda_{i}\left(\boldsymbol{\Sigma}_{S_{0}, S_{0}}\right)-2 r \epsilon & \text { by }\|\hat{\boldsymbol{A}}-\boldsymbol{\Sigma}\|_{\text {op }} \leq \epsilon \\
& =\max _{\boldsymbol{V} \in \mathcal{F}} \operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{\Sigma} \boldsymbol{V}\right)-2 r \epsilon &
\end{array}
$$

with probability $1-o(1)$.
Therefore, we have the quality of primal feasible solution $V$ is close enough to the global optimal value.

### 3.3. Proof of Proposition 1.1

However, the above Theorem 4 requires the Assumptions 1, 2, 3, 4. In a more general case, suppose the support set cannot be recovered correctly. We could still verify the quality of a given primal feasible solution by solving SOCIP-impl. Here we formally state the Proposition 1.1:
Proposition 1.1. Let samples $\boldsymbol{x}_{1}, \ldots, \boldsymbol{x}_{M}$ be i.i.d. generated from sub-Gaussian distribution with zero-mean, and $\boldsymbol{\Sigma}$ as the true underlying covariance matrix (second moment). Let $\boldsymbol{A}:=\frac{1}{M} \boldsymbol{X} \boldsymbol{X}^{\top}$ be the sample covariance matrix defined as before. Let $\boldsymbol{V}_{\text {app }}$ be a $(1-\Delta)$-approximation primal feasible solution corresponding to $\boldsymbol{A}$ where the value of $\Delta$ is obtained from its upper (dual) bound via solving SOCIP-impl. If the number of samples $M$ is sufficiently large (i.e., $M \geq[(C \sqrt{d}+t) / \epsilon]^{2}$ with $C$ a constant depends on the sub-Gaussian norm of $\boldsymbol{x}$ for any $\left.t>0\right)$, then:

$$
\operatorname{Tr}\left(\boldsymbol{V}_{\text {app }}^{\top} \boldsymbol{\Sigma} \boldsymbol{V}_{\text {app }}\right) \geq(1-\Delta) \cdot \max _{\boldsymbol{V} \in \mathcal{F}} \operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{\Sigma} \boldsymbol{V}\right)-(2-\Delta) r \epsilon
$$

holds with high probability, where $\epsilon>0$ is any given constant.

Proof. Note the Remark 5.40 in ? ensures that: as $M \geq[(C \sqrt{d}+t) / \epsilon]^{2}$ with $C$ and $t$ be constants defined as above, we have $\|\boldsymbol{A}-\boldsymbol{\Sigma}\|_{\text {op }} \leq \epsilon$ holds with probability at least $1-2 \exp \left(-c t^{2}\right)$ with $c>0$ another constant only depends on the sub-Gaussian norm of $\boldsymbol{x}$.

Now we have the objective value corresponding to $\boldsymbol{V}_{\text {app }}$ satisfies

$$
\operatorname{Tr}\left(\boldsymbol{V}_{\text {app }}^{\top} \boldsymbol{\Sigma} \boldsymbol{V}_{\text {app }}\right)
$$

$$
\geq \operatorname{Tr}\left(\boldsymbol{V}_{\text {app }}^{\top} \boldsymbol{A} \boldsymbol{V}_{\mathrm{app}}\right)-r \epsilon \quad \text { by }\|\boldsymbol{A}-\boldsymbol{\Sigma}\|_{\mathrm{op}} \leq \epsilon \text { with high probability }
$$

$$
\geq(1-\Delta) \cdot \max _{\boldsymbol{V} \in \mathcal{F}} \operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{A} \boldsymbol{V}\right)-r \epsilon \quad \text { by definition of }(1-\Delta) \text {-approximation primal feasible }
$$

$$
\geq(1-\Delta) \cdot \max _{\boldsymbol{V} \in \mathcal{F}} \operatorname{Tr}\left(\boldsymbol{V}^{\top} \boldsymbol{\Sigma} \boldsymbol{V}\right)-(2-\Delta) r \epsilon \quad \text { by }\|\boldsymbol{A}-\boldsymbol{\Sigma}\|_{\mathrm{op}} \leq \epsilon \text { with high probability. }
$$

Therefore, the proposition 1.1 holds with high probability.
Remark 3.3. Notice that the Proposition 1.1 does not require Assumptions 1, 2, 3, 4. In contrast, since the gap between the primal bounds and the dual bounds of $\boldsymbol{A}$ is computed via convex integer program SOCIP-impl, the Proposition 1.1 only request a generalized version of Assumption 1 to ensure the a "good" estimating for the true covariance $\boldsymbol{\Sigma}$.

## 4. Polynomial Running Time

Recall that the second-order-cone $\mathcal{C}^{\prime}$ and revised piecewise linear upper approximation set PLA ${ }^{\prime}$ are defined as follows:

$$
\begin{aligned}
\mathcal{C}^{\prime}:=\left\{\begin{array}{c}
\left\|[\boldsymbol{V}]_{[d], i}\right\|_{2}^{2} \leq 1, \forall i \in[r] \\
\left.\boldsymbol{V}: \begin{array}{l}
\left\|[\boldsymbol{V}]_{[d], i_{1}} \pm[\boldsymbol{V}]_{[d], i_{2}}\right\|_{2}^{2} \leq 2, \forall i_{1} \neq i_{2} \in[r] \\
\sum_{j=1}^{d}\left\|[\boldsymbol{V}]_{j,[r]}\right\|_{2} \leq \sqrt{r k} \\
\left\|[\boldsymbol{V}]_{j,[r]}\right\|_{2} \in[0,1], \forall j \in[d]
\end{array}\right\} . \\
\text { PLA }^{\prime}:=\left\{\begin{array}{l}
g_{j i}=\boldsymbol{a}_{j}^{\top} \boldsymbol{v}_{i},(j, i) \in[d] \times[r] \\
g_{j i}=\sum_{\ell=-N}^{N} \gamma_{j i}^{\ell} \eta_{j i}^{\ell},(j, i) \in J^{+} \times[r] \\
\xi_{j i}=\sum_{\ell=-N}^{N}\left(\gamma_{j i}^{\ell}\right)^{2} \eta_{j i}^{\ell} \\
\left(\eta_{j i}^{\ell}\right)_{\ell=-N}^{N} \in \operatorname{SOS}-\mathrm{II}
\end{array}\right\} .
\end{array} .\right.
\end{aligned}
$$

The implemented version of second-order-cone integer program is

$$
\begin{align*}
\max & \sum_{j \in J^{+}}\left(\lambda_{j}-\phi\right) \sum_{i=1}^{r} \xi_{j i}-s+r \phi \\
\text { s.t } & \boldsymbol{V} \in \mathcal{C}^{\prime},(g, \xi, \eta) \in \mathrm{PLA}^{\prime} \\
& \sum_{j \in J^{-}}\left(\phi-\lambda_{j}\right) \sum_{i=1}^{r} g_{j i}^{2} \leq s  \tag{2}\\
& \sum_{i=1}^{r} g_{j i}^{2} \leq \theta_{j}^{2}, \forall j \in[d]  \tag{3}\\
& \sum_{j \in J^{+}}\left(\lambda_{j}-\phi\right) \sum_{i=1}^{r} \xi_{j i}-s+r \phi \leq[\boldsymbol{A}]_{j_{1}, j_{1}}+\cdots+[\boldsymbol{A}]_{j_{k}, j_{k}}+\sum_{j \in J^{+}} \frac{r\left(\lambda_{j}-\phi\right) \theta_{j}^{2}}{4 N^{2}} \\
& \sum_{j=1}^{d} \lambda_{j} \xi_{j i} \geq \boldsymbol{v}_{i+1}^{\top} \boldsymbol{A} \boldsymbol{v}_{i+1}, \forall i \in[r-1] \\
& \sum_{j=1}^{d}\left[\boldsymbol{v}_{i}\right]_{j} \geq 0, \forall i \in[r]
\end{align*}
$$

(SOCIP-impl)
where constraint (1) is the convex relaxation and piecewise linear approximation for SPCA combined with thresholding technique that we discussed in Section 2.4 of main content; constraints (2), (3), (4) are the cutting planes; constraints (5), (6) are the symmetric breaking constraints. Here are have

Theorem 5. Given the number of splitting points $N$, the size of large eigenvalue set $\left|J^{+}\right|$and its corresponding threshold parameter $\phi$, then the SOCIP-impl can be solved within polynomial time.

Proof. The only non-convex part in SOCIP-impl is the SOS-II constraints for variables $\eta$. Since the SOS-II constraints imply that: for each $(j, i) \in J^{+} \times[r]$, at most two continuous variables of $\left(\eta_{j i}^{\ell}\right)_{\ell=-N}^{N}$ are non-zero, then there are $2 N$ possible situations for non-zero $\left(\eta_{j i}^{\ell}\right)_{\ell=-N}^{N}$, and thus $(2 N)^{r\left|J^{+}\right|}$possible situations for all $(j, i) \in J^{+} \times[r]$. Given a fixed possible situation of SOS-II variables $\left(\eta_{j i}^{\ell}\right)_{\ell=-N}^{N}$ for all $(j, i) \in J^{+} \times[r]$, note that the SOCIP-impl can be transferred into
a continuous convex programming (CP) which is able to optimized exactly within polynomial time $T$ (corresponding to the size of input CP, and the additive gap). Therefore the SOCIP-impl can be solved within polynomial time $T \cdot(2 N)^{r \mid J^{+}}$| by checking each possible situations.

Remark 4.1. Note that the Theorem 5 still holds even when the SOCIP-impl only contains constraint (1). The rest of constraints (2)-(6) are used to improve the running time in practice. Thus people can decide which constraint from (2)-(6) is needed when optimizing their own instances.

## 5. Main Numerical Results

In this section, we present the original numerical results of SOCIP-impl and Baselines on two types of instances described in Section 3 of main content. Three methods for upper bounds are listed as follows:

$$
\begin{array}{lll}
\text { Baseline-1 }:=[\boldsymbol{A}]_{j_{1}, j_{1}}+\cdots+[\boldsymbol{A}]_{j_{k}, j_{k}}, & \text { where }[\boldsymbol{A}]_{j_{1}, j_{1}} \geq[\boldsymbol{A}]_{j_{2}, j_{3}} \geq \cdots[\boldsymbol{A}]_{j_{d}, j_{d}} \\
\text { Baseline-2 }:= & \max _{\boldsymbol{P}} \operatorname{Tr}(\boldsymbol{A P}), & \text { s.t. } \boldsymbol{I}_{d} \succeq \boldsymbol{P} \succeq \mathbf{0}, \operatorname{Tr}(\boldsymbol{P})=r, \mathbf{1}^{\top}|\boldsymbol{P}| \mathbf{1} \leq r k \\
\text { SOCIP-impl }:= & \sum_{j \in J^{+}}\left(\lambda_{j}-\phi\right) \sum_{i=1}^{r} \xi_{j i}-s+r \phi & \text { s.t. constraint (1), (2), (3), (4), (5) }
\end{array}
$$

The lower bounds (LB) are computed by feasible solutions of SPCA which obtained from local search algorithm 1 with randomized initialization. The column "Gap" denotes the gap between upper bounds of SOCIP-impl and lower bounds (LB) defined as Gap $:=\frac{\text { SOCIP-LB }}{\text { LB }}$. All original numerical results are reported in the following tables.

Notice that because of the limitation of hardware and software, the Baseline-2 (SDP relaxation method) is hard to scalable since the quadratic increasing of the number of variables in the lifted space. Thus the Baseline-2 (SDP relaxation method) only works for the case Eisen-1 and Eisen-2, and in rest of the tables we remove the column of "Baseline-2".

| Para $(d, k, r)$ | LB | Baseline-1 | Baseline-2 | SOCIP | Gap |
| :--- | :---: | :---: | :---: | :---: | :---: |
| $(79,5,2)$ | 15.277 | 16.295 | 20.081 | 16.351 | $7.028 \%$ |
| $(79,10,2)$ | 19.705 | 21.376 | 21.530 | 21.184 | $7.505 \%$ |
| $(79,15,2)$ | 20.590 | 24.034 | 22.036 | 21.678 | $5.282 \%$ |
| $(79,20,2)$ | 21.020 | 25.881 | 22.197 | 21.833 | $3.866 \%$ |
| $(79,25,2)$ | 21.274 | 27.279 | 22.203 | 21.981 | $3.321 \%$ |
| $(79,30,2)$ | 21.481 | 28.535 | 22.203 | 22.069 | $2.736 \%$ |
| $(79,5,3)$ | 16.150 | 16.295 | 21.998 | 16.438 | $1.783 \%$ |
| $(79,10,3)$ | 20.569 | 21.376 | 23.806 | 21.54 | $4.721 \%$ |
| $(79,15,3)$ | 21.553 | 24.034 | 24.493 | 24.139 | $11.997 \%$ |
| $(79,20,3)$ | 21.683 | 25.881 | 24.725 | 24.599 | $13.447 \%$ |
| $(79,25,3)$ | 23.205 | 27.279 | 24.738 | 24.427 | $5.268 \%$ |
| $(79,30,3)$ | 23.229 | 28.535 | 24.738 | 24.527 | $5.588 \%$ |

Table 1. Compare SOCIP-impl with baseline for Eisen-1

| Para $(d, k, r)$ | LB | Baseline-1 | SOCIP | Gap |
| :--- | :---: | :---: | :---: | :---: |
| $(500,5,1)$ | 1646.454 | 1720.878 | 1723.119 | $4.656 \%$ |
| $(500,10,1)$ | 2641.229 | 2970.226 | 2970.658 | $12.476 \%$ |
| $(500,20,1)$ | 4255.287 | 5015.718 | 5007.157 | $17.669 \%$ |
| $(500,40,1)$ | 6924.530 | 8280.635 | 8242.987 | $19.040 \%$ |
| $(500,80,1)$ | 10741.925 | 13292.953 | 13183.996 | $22.734 \%$ |
| $(500,120,1)$ | 13660.302 | 17349.272 | 17165.391 | $25.659 \%$ |
| $(500,160,1)$ | 15666.335 | 20599.533 | 19154.019 | $22.262 \%$ |
| $(500,5,2)$ | 1709.958 | 1720.878 | 1733.222 | $1.361 \%$ |
| $(500,10,2)$ | 2794.140 | 2970.226 | 2989.336 | $6.986 \%$ |
| $(500,20,2)$ | 4510.085 | 5015.718 | 5045.065 | $11.862 \%$ |
| $(500,40,2)$ | 7245.277 | 8280.635 | 8326.816 | $14.928 \%$ |
| $(500,80,2)$ | 11226.442 | 13292.953 | 13359.151 | $18.997 \%$ |
| $(500,120,2)$ | 14163.219 | 17349.272 | 17429.896 | $23.065 \%$ |
| $(500,160,2)$ | 16457.275 | 20599.533 | 19070.61 | $15.880 \%$ |

Table 3. Compare SOCIP-impl with baseline for CovColon

| Para $(d, k, r)$ | LB | Baseline-1 | Baseline-2 | SOCIP | Gap |
| :--- | :---: | :---: | :---: | :---: | :---: |
| $(118,5,2)$ | 8.144 | 8.574 | 14.008 | 8.61 | $5.724 \%$ |
| $(118,10,2)$ | 13.686 | 15.051 | 22.211 | 15.094 | $10.291 \%$ |
| $(118,15,2)$ | 18.328 | 20.641 | 26.984 | 20.696 | $12.917 \%$ |
| $(118,20,2)$ | 22.155 | 25.845 | 29.322 | 25.473 | $14.975 \%$ |
| $(118,25,2)$ | 25.040 | 30.018 | 30.786 | 27.807 | $11.052 \%$ |
| $(118,30,2)$ | 27.311 | 33.461 | 31.814 | 29.604 | $8.397 \%$ |
| $(118,5,3)$ | 8.434 | 8.574 | 19.328 | 8.635 | $2.385 \%$ |
| $(118,10,3)$ | 14.457 | 15.051 | 28.708 | 15.148 | $4.777 \%$ |
| $(118,15,3)$ | 19.296 | 20.641 | 32.086 | 20.762 | $7.596 \%$ |
| $(118,20,3)$ | 23.583 | 25.845 | 34.152 | 25.977 | $10.152 \%$ |
| $(118,25,3)$ | 26.734 | 30.018 | 35.545 | 30.164 | $12.831 \%$ |
| $(118,30,3)$ | 28.741 | 33.461 | 36.495 | 33.409 | $16.242 \%$ |

Table 2. Compare SOCIP-impl with baseline for Eisen-2

| Para $(d, k, r)$ | LB | Baseline-1 | SOCIP | Gap |
| :--- | :---: | :---: | :---: | :---: |
| $(500,5,1)$ | 4300.497 | 5177.405 | 5184.741 | $20.561 \%$ |
| $(500,10,1)$ | 6008.317 | 8901.180 | 8902.889 | $48.176 \%$ |
| $(500,20,1)$ | 9082.158 | 15160.617 | 14641.435 | $61.211 \%$ |
| $(500,40,1)$ | 13107.045 | 24293.415 | 19557.092 | $49.211 \%$ |
| $(500,80,1)$ | 17544.277 | 38358.967 | 24458.286 | $39.409 \%$ |
| $(500,120,1)$ | 20797.933 | 48691.305 | 27058.292 | $30.101 \%$ |
| $(500,160,1)$ | 23310.903 | 57395.584 | 28527.242 | $22.377 \%$ |
| $(500,5,2)$ | 4990.132 | 5177.405 | 5220.062 | $4.608 \%$ |
| $(500,10,2)$ | 8125.266 | 8901.180 | 8951.928 | $10.174 \%$ |
| $(500,20,2)$ | 11868.012 | 15160.617 | 15226.865 | $28.302 \%$ |
| $(500,40,2)$ | 16138.886 | 24293.415 | 24378.169 | $51.052 \%$ |
| $(500,80,2)$ | 21396.692 | 38358.968 | 30178.957 | $41.045 \%$ |
| $(500,120,2)$ | 25579.788 | 48691.305 | 33116.474 | $29.463 \%$ |
| $(500,160,2)$ | 28950.488 | 57395.584 | 35369.37 | $22.172 \%$ |

Table 4. Compare SOCIP-impl with baseline for Lymp

| Para $(d, k, r)$ | LB | Baseline-1 | SOCIP | Gap |
| :--- | :---: | :---: | :---: | :---: |
| $(1000,5,1)$ | 1665.714 | 3144.0 | 1994.758 | $19.754 \%$ |
| $(1000,10,1)$ | 1766.499 | 3936.0 | 2036.647 | $15.293 \%$ |
| $(1000,20,1)$ | 1834.473 | 5058.585 | 2061.048 | $12.351 \%$ |
| $(1000,40,1)$ | 1918.158 | 6355.687 | 2083.412 | $8.615 \%$ |
| $(1000,80,1)$ | 1993.947 | 7876.078 | 2094.868 | $5.061 \%$ |
| $(1000,5,2)$ | 2208.062 | 3144.0 | 2669.481 | $20.897 \%$ |
| $(1000,10,2)$ | 2358.341 | 3936.0 | 2842.534 | $20.531 \%$ |
| $(1000,20,2)$ | 2583.484 | 5058.585 | 2963.481 | $14.709 \%$ |
| $(1000,40,2)$ | 2762.794 | 6355.687 | 3093.255 | $11.961 \%$ |
| $(1000,80,2)$ | 2681.272 | 7876.078 | 3128.75 | $16.689 \%$ |
| $(1000,5,3)$ | 2650.855 | 3144.0 | 6745.96 | $154.482 \%$ |
| $(1000,10,3)$ | 2839.634 | 3936.0 | 7373.92 | $159.679 \%$ |
| $(1000,20,3)$ | 2996.054 | 5058.585 | 7616.909 | $154.231 \%$ |
| $(1000,40,3)$ | 3361.852 | 6355.687 | 7659.558 | $127.837 \%$ |
| $(1000,80,3)$ | 3657.417 | 7876.078 | 7362.175 | $101.294 \%$ |


| Para $(d, k, r)$ | LB | Baseline-1 | SOCIP | Gap |
| :--- | :---: | :---: | :---: | :---: |
| $(1500,5,1)$ | 1694.560 | 3327.0 | 2198.699 | $29.750 \%$ |
| $(1500,10,1)$ | 1834.871 | 4344.0 | 2333.223 | $27.160 \%$ |
| $(1500,20,1)$ | 1965.926 | 5649.585 | 2449.516 | $24.599 \%$ |
| $(1500,40,1)$ | 2317.149 | 7316.316 | 2593.001 | $11.905 \%$ |
| $(1500,80,1)$ | 2537.644 | 9129.687 | 2630.608 | $3.663 \%$ |
| $(1500,5,2)$ | 2241.333 | 3327.0 | 4679.819 | $108.796 \%$ |
| $(1500,10,2)$ | 2552.943 | 4344.0 | 5023.951 | $96.791 \%$ |
| $(1500,20,2)$ | 3073.350 | 5649.585 | 5348.092 | $74.015 \%$ |
| $(1500,40,2)$ | 3178.420 | 7316.316 | 5461.002 | $71.815 \%$ |
| $(1500,80,2)$ | 3596.531 | 9129.687 | 5912.314 | $64.389 \%$ |
| $(1500,5,3)$ | 2675.918 | 3327.0 | 7548.85 | $182.103 \%$ |
| $(1500,10,3)$ | 3025.420 | 4344.0 | 8432.259 | $178.714 \%$ |
| $(1500,20,3)$ | 3587.461 | 5649.585 | 9568.021 | $166.707 \%$ |
| $(1500,40,3)$ | 3722.302 | 7316.316 | 11705.781 | $214.477 \%$ |
| $(1500,80,3)$ | 4639.675 | 9129.687 | 12523.76 | $169.928 \%$ |

Table 5. Compare SOCIP-impl with baseline for Reddit of size 1000Table 6. Compare SOCIP-impl with baseline for Reddit of size 1500

| Para $(d, k, r)$ | LB | Baseline-1 | SOCIP | Gap |
| :--- | :---: | :---: | :---: | :---: |
| $(500,5,1)$ | 55.085 | 252.808 | 56.22 | $2.060 \%$ |
| $(500,10,1)$ | 56.318 | 360.936 | 56.474 | $0.277 \%$ |
| $(500,20,1)$ | 56.342 | 371.594 | 56.474 | $0.235 \%$ |
| $(500,40,1)$ | 56.370 | 392.540 | 56.474 | $0.184 \%$ |
| $(500,80,1)$ | 56.403 | 433.838 | 56.474 | $0.126 \%$ |
| $(500,5,2)$ | 107.422 | 252.810 | 109.825 | $2.237 \%$ |
| $(500,10,2)$ | 109.971 | 360.936 | 110.307 | $0.306 \%$ |
| $(500,20,2)$ | 110.001 | 371.594 | 110.308 | $0.279 \%$ |
| $(500,40,2)$ | 110.046 | 392.540 | 110.309 | $0.239 \%$ |
| $(500,80,2)$ | 110.109 | 433.834 | 113.142 | $2.754 \%$ |

Table 7. Compare SOCIP-impl with baseline for block spiked covariance of size 500

| Para $(d, k, r)$ | LB | Baseline-1 | SOCIP | Gap |
| :--- | :---: | :---: | :---: | :---: |
| $(2000,5,1)$ | 4300.497 | 5177.405 | 5184.741 | $20.561 \%$ |
| $(2000,10,1)$ | 6008.317 | 8901.180 | 8902.889 | $48.176 \%$ |
| $(2000,20,1)$ | 9082.158 | 15160.617 | 14641.435 | $61.211 \%$ |
| $(2000,40,1)$ | 13107.045 | 24293.415 | 19557.092 | $49.211 \%$ |
| $(2000,80,1)$ | 17544.277 | 38358.967 | 24458.286 | $39.409 \%$ |
| $(2000,120,1)$ | 20797.933 | 48691.305 | 27058.292 | $30.101 \%$ |
| $(2000,160,1)$ | 23310.903 | 57395.584 | 28527.242 | $22.377 \%$ |
| $(2000,5,2)$ | 4990.132 | 5177.405 | 5220.062 | $4.608 \%$ |
| $(2000,10,2)$ | 8125.266 | 8901.180 | 8951.928 | $10.174 \%$ |
| $(2000,20,2)$ | 11868.012 | 15160.617 | 15226.865 | $28.302 \%$ |
| $(2000,40,2)$ | 16138.886 | 24293.415 | 24378.169 | $51.052 \%$ |
| $(2000,80,2)$ | 21396.692 | 38358.968 | 30178.957 | $41.045 \%$ |
| $(2000,120,2)$ | 25579.788 | 48691.305 | 33116.474 | $29.463 \%$ |
| $(2000,160,2)$ | 28950.488 | 57395.584 | 35369.37 | $22.172 \%$ |

Table 8. Compare SOCIP-impl with baseline for Reddit of size 2000


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