Free Energy Wells and Overlap Gap Property in Sparse PCA  
(Extended Abstract)

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Abstract

We study a variant of the sparse PCA (principal component analysis) problem in the “hard” regime, where the inference task is possible yet no polynomial-time algorithm is known to exist. Prior work, based on the low-degree likelihood ratio, has conjectured a precise expression for the best possible (sub-exponential) runtime throughout the hard regime. Following instead a statistical physics inspired point of view, we show bounds on the depth of free energy wells for various Gibbs measures naturally associated to the problem. These free energy wells imply hitting time lower bounds that corroborate the low-degree conjecture: we show that a class of natural MCMC (Markov chain Monte Carlo) methods (with worst-case initialization) cannot solve sparse PCA with less than the conjectured runtime. These lower bounds apply to a wide range of values for two tuning parameters: temperature and sparsity misparametrization. Finally, we prove that the Overlap Gap Property (OGP), a structural property that implies failure of certain local search algorithms, holds in a significant part of the hard regime.\footnote{This paper is an extended abstract. The full version appears as arXiv preprint arXiv:2006.10689v1.}

1. The Model

We consider the following variant of sparse PCA in the spiked Wigner model (also called principal submatrix recovery). Let $W$ be a GOE($n$) matrix, i.e., $n \times n$ symmetric with off-diagonal entries $\mathcal{N}(0, 1/n)$ and diagonal entries $\mathcal{N}(0, 2/n)$, all independent aside from the symmetry $W_{ij} = W_{ji}$. Let $x$ be an unknown $k$-sparse vector in $\{0, 1\}^n$, i.e., exactly $k$ entries are equal to 1. We are interested in recovering $x$ from the observation

$$Y = \frac{\lambda}{k} xx^\top + W$$

where $\lambda > 0$ is the signal-to-noise ratio. We study the problem in the limit $n \to \infty$, where the parameters $\lambda = \lambda_n$ and $k = k_n$ may depend on $n$. We are primarily interested in the exact recovery problem: we study algorithms which given $Y$, output $x$ with high probability, i.e., probability tending to 1 as $n \to \infty$. Our regime of interest will be $1 \ll k \ll n$. Throughout, we use the notation $\ll$ to hide factors of $n^{o(1)}$ (although in most cases, $\ll$ will only hide logarithmic factors).

2. Our Contributions

Prior work suggests the existence of a “hard regime” \( \sqrt{k/n} \ll \lambda \ll \min\{1, k/\sqrt{n}\} \) where exact recovery is information-theoretically possible but no polynomial-time algorithm is known (see e.g. Baik et al. (2005); Féral and Péché (2007); Amini and Wainwright (2008); Benaych-Georges and Nadakuditi (2009); Johnstone and Lu (2009); Banks et al. (2018); Ding et al. (2019)). More specifically, the work of Ding et al. (2019) suggests the following conjecture regarding a precise expression for the best possible (sub-exponential) runtime throughout the hard regime.

**Conjecture 1** Consider the sparse PCA problem as defined in Section 1. For any \( \lambda \) in the “hard” regime \( \sqrt{k/n} \ll \lambda \ll \min\{1, k/\sqrt{n}\} \), any algorithm requires runtime \( \exp \left( \tilde{\Omega} \left( \frac{k^2}{\lambda n} \right) \right) \) to achieve exact recovery.

This prediction is made by Ding et al. (2019) (for a variant of our model where \( x_i \in \{0, -1, 1\} \)) using the low-degree likelihood ratio (Hopkins and Steurer, 2017; Hopkins et al., 2017; Hopkins, 2018), which amounts to studying the power of algorithms based on low-degree polynomials. There are known algorithms which achieve the matching runtime \( \exp \left( \tilde{O} \left( \frac{k^2}{\lambda n} \right) \right) \) (Ding et al., 2019; Holtzman et al., 2019). For instance, the following simple algorithm of Ding et al. (2019) proceeds in two steps. The first step is to let \( k' \approx \frac{k^2}{\lambda n} \) and solve, by exhaustive search, the optimization problem

\[
\arg\max_{v \in S_{k'}} v^\top Y v
\]

(1)

where \( S_{k'} \) is the space of \( k' \)-sparse vectors

\[
S_{k'} = \{ v \in \{0, 1\}^n : \|v\|_0 = k' \},
\]

(2)

and the final step uses the optimizer \( v^* \) to exactly recover \( x \) via a simple boosting procedure.

In this work we give evidence in support of Conjecture 1 by showing

- the existence of free energy wells (Ben Arous et al., 2018; Gamarnik and Zadik, 2017b, 2019; Gamarnik et al., 2019) in the Gibbs measure (at various temperatures) associated with the optimization problem (1),

- and the existence of the overlap gap property (Gamarnik and Zadik, 2017a,b; Gamarnik et al., 2019; Gamarnik and Zadik, 2019; Zadik, 2019) in the space of feasible solutions of the optimization problem (1),

for various choices of the tuning parameter \( k' \). As explained in the full version of the present work, a free energy well of depth \( D \) at inverse temperature \( \beta \) implies that a certain class of MCMC methods with parameter \( \beta \) requires time at least \( \exp(\Omega(D)) \) to solve (1). Our main result can be stated informally as follows.

**Theorem** (Main result, informal) Suppose \( \lambda \) is in the “hard” regime \( \sqrt{k/n} \ll \lambda \ll \min\{1, k/\sqrt{n}\} \) and that additionally, \( \lambda \ll (k/n)^{1/4} \). For any “informative” \( k' \) and any \( \beta \geq 0 \) (possibly depending on \( n \)), there exists a free energy well of depth \( \tilde{\Omega} \left( \frac{k^2}{\lambda^2 n} \right) \) and the overlap gap property holds, with high probability.
Here "informative" $k'$ refers to the condition, $\frac{k'^2}{\lambda^2 n} \lesssim k' \lesssim \lambda^2 n$, which captures the $k'$ values for which solving the optimization problem (1) is actually useful in the sense that a near-optimal solution can be used to exactly recover $x$ via a simple boosting procedure. Our main result shows that if the condition $\lambda \ll (k/n)^{1/4}$ is satisfied then MCMC cannot improve the runtime of Ding et al. (2019); Holtzman et al. (2019) for any choice of inverse temperature $\beta$ and any (informative) choice of misparametrization $k'$. The main weakness of the result is the condition $\lambda \ll (k/n)^{1/4}$, which is an artifact of the proof. However, in the relatively sparse regime $k \ll n^{1/3}$, the condition $\lambda \ll (k/n)^{1/4}$ holds throughout the entire "hard" regime. Thus we obtain a complete refutation of MCMC methods (across all $\beta$ and $k'$) throughout a large range of sparsity values (namely $k \ll n^{1/3}$).

Our results are actually somewhat stronger than what we have stated here: even when the condition $\lambda \ll (k/n)^{1/4}$ does not hold, the result still holds for some $k'$ values; in particular, it always holds for all informative $k' \leq k$. One consequence of this is that it is not possible to speed up the algorithm of Ding et al. (2019) by taking their choice of $k' \approx \frac{k^2}{\lambda^2 n}$ (the smallest "informative" $k'$, which in particular is less than $k$) and solving (1) via MCMC instead of exhaustive search.

Remark 2 In order for a computational hardness result to be most compelling, the class of algorithms ruled out should capture the best known algorithms. This is indeed the case here in the sense that there exists a choice of parameters, namely $k' \approx \frac{k^2}{\lambda^2 n}$ and $\beta = 0$, for which MCMC (followed by boosting) mimics the algorithm of Ding et al. (2019) and achieves the runtime $\exp(\tilde{O}(\frac{k^2}{\lambda^2 n}))$. For this choice of parameters, MCMC is simply a random walk (ignoring the data $Y$) on the space of $k'$-sparse vectors, which will visit all states within time $\exp(\tilde{O}(k')) = \exp(\tilde{O}(\frac{k^2}{\lambda^2 n}))$ with high probability. A consequence is that for this choice of $k'$ and $\beta$, any free energy well has depth $\tilde{O}(\frac{k^2}{\lambda^2 n})$ and so our lower bound is tight. It is not clear whether MCMC with more natural parameters (e.g. $k' = k$) matches the above runtime; it might in fact be strictly worse. This highlights the importance of allowing $k' \neq k$ in our main result.

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