

A Notation

\mathbb{N}	The set of nonnegative integers.
\mathbb{R}	The set of real numbers.
\mathbf{E}	The expectation operator.
\Pr	$\Pr(x y)$ denotes the probability of the stochastic variable x given y .
arg max	$\pi^* = \arg \max_{\pi \in \Pi} f_\pi$ denotes an element $\pi^* \in \Pi$ that maximizes the function f_π .
arg max min	$\pi^* = \arg \max_{\pi \in \Pi} \min_{o \in O} f_{\pi,o}$ denotes an element $(\pi^*, o^*) \in \Pi \times O$ that takes the maxmin over $f_{\pi,o}$.
\geq	For $\lambda = (\lambda_1, \dots, \lambda_J)$, $\lambda \geq 0$ denotes that $\lambda_i \geq 0$ for $i = 1, \dots, J$.
$1_X(x)$	$1_X(x) = 1$ if $x \in \{X\}$ and 0 otherwise
$\mathbf{1}_n$	$\mathbf{1}_n = (1, 1, \dots, 1) \in \mathbb{R}^n$.
$N(t, s, a, b)$	$N(t, s, a, b) = \sum_{k=1}^t 1_{(s,a,b)}(s_k, a_k, b_k)$.
e	$e : (s, a, o) \mapsto 1$.
$ S $	Denotes the number of elements in S .

B Examples

The following example from [Altman \(1999\)](#) describes in more detail a model where we have a Markov decision process with constraints and where the agent doesn't have model knowledge.

Example 1 (Altman, 1999). *Consider a discrete time single-server queue with a buffer of finite size L . For a given time slot, we assume that at most one customer may join the system. The state of the system at a given time slot is the number of customers in the queue. There is a delay cost $c(s)$ given a state $s \in \{S_1, \dots, S_n\}$ which one would like to keep as low as possible. The probability of a service to be completed is a^1 , where $1/a_1$ is the Quality of Service (QoS). The probability of queue arrival at time t is a^2 . The actions are given by a^1 and a^2 . Let $c^1(a^1)$ be the cost to complete the service (c^1 is increasing in a^1). c^1 should be bounded by some value v^1 . There is a cost corresponding to the throughput, $c^2(a^2)$, (c^2 is decreasing in a^2). c^2 should be bounded by some value v^2 . We assume that the number of actions is finite and actions sets are given by $a^1 \in \{A_1^1, \dots, A_{l_1}^1\}$ and $a^2 \in \{A_1^2, \dots, A_{l_2}^2\}$ where $0 < A_1^1 \leq \dots \leq A_{l_1}^1 \leq 1$ and $0 \leq A_1^2 \leq \dots \leq A_{l_2}^2 \leq 1$. The transition probability $P(s_{k+1}, s_k, a_k^1, a_k^2)$ from state s_k to s_{k+1} given actions a_k^1 and a_k^2 is given by*

$$P(s_+, s, a^1, a^2) = \begin{cases} (1 - a^2)a^1 & \text{if } L \geq s \geq 1, \\ & s_+ = s - 1 \\ a^2a^1 + (1 - a^2)(1 - a^1) & \text{if } L \geq s \geq 1, \\ & s_+ = s \\ a^2(1 - a^1) & \text{if } L \geq s \geq 0, \\ & s_+ = s + 1 \\ 1 - a^2(1 - a^1) & \text{if } L \geq s \geq 0, \\ & s_+ = s = 0 \end{cases}$$

For $\gamma \in (0, 1)$, the constrained Markov decision process problem is given by

$$\begin{aligned} \min_{\pi^1, \pi^2} \quad & \mathbf{E} \left(\sum_{k=0}^{\infty} \gamma^k c(s_k) \right) \\ \text{s. t.} \quad & \mathbf{E} \left(\sum_{k=0}^{\infty} \gamma^k c^1(\pi^1(s_k)) \right) \leq v^1 \\ & \mathbf{E} \left(\sum_{k=0}^{\infty} \gamma^k c^2(\pi^2(s_k)) \right) \leq v^2, \end{aligned} \tag{28}$$

which is equivalent to

$$\begin{aligned}
 & \max_{\pi^1, \pi^2} \mathbf{E} \left(\sum_{k=0}^{\infty} \gamma^k r(s_k, \pi(s_k)) \right) \\
 & \text{s. t. } \mathbf{E} \left(\sum_{k=0}^{\infty} \gamma^k r^1(s_k, \pi(s_k)) \right) \geq 0 \\
 & \mathbf{E} \left(\sum_{k=0}^{\infty} \gamma^k r^2(s_k, \pi(s_k)) \right) \geq 0,
 \end{aligned} \tag{29}$$

where $a_k = (a_k^1, a_k^2)$, $\pi(s_k) = (\pi^1(s_k), \pi^2(s_k))$, $r(s_k, a_k) = -c(s_k)$, $r^1(s_k, a_k) = -c^1(a_k^1) + v^1 \cdot (1 - \gamma)$, $r^2(s_k, a_k) = -c^2(a_k^2) + v^2 \cdot (1 - \gamma)$

Example 2 (Search Engine). *In a search engine, there is a number of documents that are related to a certain query. There are two values that are related to every document, the first being a (advertisement) value u_i of document i for the search engine and the second being a value v_i for the user (could be a measure of how strongly related the document is to the user query). The task of the search engine is to display the documents in a row some order, where each row has an attention value, A_j for row j . We assume that u_i and v_i are known to the search engine for all i , whereas the attention values $\{A_j\}$ are not known. The strategy π of the search engine is to display document i in position j , $\pi(i) = j$, with probability p_{ij} . Thus, the expected average reward for the search engine is*

$$R^e = \lim_{N \rightarrow \infty} \frac{1}{N} \sum_{i=1}^N \mathbf{E}(u_i A_{\pi(i)})$$

and for the user

$$R^u = \lim_{N \rightarrow \infty} \frac{1}{N} \sum_{i=1}^N \mathbf{E}(v_i A_{\pi(i)}).$$

The search engine has multiple objectives here where it wants to maximize the rewards for the user and itself. One solution is to define a measure for the quality of service for the user, $R^u \geq \underline{R}^u$ and at the same time satisfy a certain lower bound \underline{R}^e of its own reward, that is

$$\begin{aligned}
 & \text{find } \pi \\
 & \text{s. t. } \lim_{N \rightarrow \infty} \frac{1}{N} \sum_{i=1}^N \mathbf{E}(u_i A_{\pi(i)}) \geq \underline{R}^e \\
 & \lim_{N \rightarrow \infty} \frac{1}{N} \sum_{i=1}^N \mathbf{E}(v_i A_{\pi(i)}) \geq \underline{R}^u
 \end{aligned}$$

C Proof of Theorem 1

The proof will rely on the following result.

Proposition 2. *The random process $\{\Delta_k\}$ taking values in \mathbb{R} and defined as*

$$\Delta_{k+1}(x) = (1 - \alpha_k(x))\Delta_k(x) + \alpha_k(x)F_k(x)$$

converges to zero with probability 1 under the following assumptions:

- i. For all x , $0 \leq \alpha_k(x) \leq 1$, $\sum_k \alpha_k(x) = \infty$, and $\sum_k \alpha_k^2(x) < \infty$
- ii. $\|\mathbf{E}(F_k(x) | \mathcal{F}_k)\|_{\infty} \leq \gamma \|\Delta_k\|_{\infty}$, with $\gamma < 1$
- iii. $\mathbf{E}(F_k - \mathbf{E}(F_k(x) | \mathcal{F}_k))^2 \leq C(1 + \|\Delta_k\|_{\infty}^2)$, for some constant $C > 0$

where \mathcal{F}_k is the sigma algebra $\sigma(\Delta_t, F_{t-1}, \alpha_{t-1}, t \leq k)$.

Proof. Consult [\(Jaakkola et al., 1994\)](#). □

Now let

$$\Delta_k(s, a, o) = Q_k(s, a, o) - Q^*(s, a, o)$$

Subtracting Q^* from the right and left hand sides of the second equality in (12) implies that

$$\begin{aligned} \Delta_{k+1}(s, a, o) &= (1 - \alpha(s, a, o))\Delta_k(s, a, o) + \\ &+ \alpha(s, a, o)(R(s, a, o) + \gamma\mathbf{E}(Q_k(s_+, \pi_k(s_+), o)) - Q^*(s, a, o)). \end{aligned}$$

We will show that Δ_k satisfies the conditions of Proposition 2. Introduce the sigma algebra $\mathcal{F}_k = \sigma(\Delta_t, F_{t-1}, \alpha_{t-1}, t \leq k)$.

Define

$$F_k(s, a, o) = 1_{(s, a, o)}(s_k, a_k, o_k) \times (R(s, a, o) + \gamma\mathbf{E}(Q_k(s_+, \pi_k(s_+), o)) - Q^*(s, a, o))$$

If $(s, a, o) \neq (s_k, a_k, o_k)$, then $F_k(s, a, o) = 0$. Else,

$$\begin{aligned} \mathbf{E}(F_k(s, a, o) \mid \mathcal{F}_k) &= \sum_{s_+} P(s, a, s_+) 1_{(s, a, o)}(s_k, a_k, o_k) \times (R(s, a, o) + \\ &\quad \gamma\mathbf{E}(Q_k(s_+, \pi_k(s_+), o)) - Q^*(s, a, o)) \\ &= \sum_{s_+} P(s, a, s_+) (R(s, a, o_k) + \\ &\quad \gamma\mathbf{E}(Q_k(s_+, \pi_k(s_+), o_k)) - Q^*(s, a, o_k)) \\ &= \sum_{s_+} P(s, a, s_+) (R(s, a, o_k) + \gamma\mathbf{E}(Q_k(s_+, \pi_k(s_+), o_k)) - \\ &\quad R(s, a, o_k) - \gamma\mathbf{E}(Q^*(s_+, \pi^*(s_+), o_k))) \\ &= \gamma \sum_{s_+} P(s, a, s_+) (\mathbf{E}(Q_k(s_+, \pi_k(s_+), o_k)) - \\ &\quad \mathbf{E}(Q^*(s_+, \pi^*(s_+), o_k))) \end{aligned}$$

If $\mathbf{E}(Q_k(s_+, \pi_k(s_+), o_k)) \geq \mathbf{E}(Q^*(s_+, \pi^*(s_+), o_k))$, then

$$\begin{aligned} &\left| \mathbf{E}[Q_k(s_+, \pi_k(s_+), o_k)] - \mathbf{E}[Q^*(s_+, \pi^*(s_+), o_k)] \right| \\ &= \mathbf{E}(Q_k(s_+, \pi_k(s_+), o_k)) - \mathbf{E}(Q^*(s_+, \pi^*(s_+), o_k)) \\ &= R(s, a, o_k) + \mathbf{E}(Q_k(s_+, \pi_k(s_+), o_k)) - R(s, a, o_k) - \mathbf{E}(Q^*(s_+, \pi^*(s_+), o_k)) \\ &\leq R(s, a, o_k) + \mathbf{E}(Q_k(s_+, \pi_k(s_+), o_k)) - R(s, a, o^*) - \mathbf{E}(Q^*(s_+, \pi^*(s_+), o^*)) \\ &\leq R(s, a, o^*) + \mathbf{E}(Q_k(s_+, \pi_k(s_+), o^*)) - R(s, a, o^*) - \mathbf{E}(Q^*(s_+, \pi^*(s_+), o^*)) \\ &\leq R(s, a, o^*) + \mathbf{E}(Q_k(s_+, \pi_k(s_+), o^*)) - R(s, a, o^*) - \mathbf{E}(Q^*(s_+, \pi_k(s_+), o^*)) \\ &= |\mathbf{E}(Q_k(s_+, \pi_k(s_+), o^*) - Q^*(s_+, \pi_k(s_+), o^*))| \\ &\leq \max_{s, a, o} |Q_k(s, a, o) - Q^*(s, a, o)| \\ &= \|Q_k - Q^*\|_\infty. \end{aligned}$$

Else, if $\mathbf{E}(Q_k(s_+, \pi_k(s_+), o_k)) \leq \mathbf{E}(Q^*(s_+, \pi^*(s_+), o_k))$, then

$$\begin{aligned} &\left| \mathbf{E}[Q_k(s_+, \pi_k(s_+), o_k)] - \mathbf{E}[Q^*(s_+, \pi^*(s_+), o_k)] \right| \\ &= \mathbf{E}(Q^*(s_+, \pi^*(s_+), o_k)) - \mathbf{E}(Q_k(s_+, \pi_k(s_+), o_k)) \\ &\leq \mathbf{E}(Q^*(s_+, \pi^*(s_+), o_k)) - \mathbf{E}(Q_k(s_+, \pi^*(s_+), o_k)) \\ &= |\mathbf{E}(Q_k(s_+, \pi^*(s_+), o_k) - Q^*(s_+, \pi^*(s_+), o_k))| \\ &\leq \max_{s, a, o} |Q_k(s, a, o) - Q^*(s, a, o)| \\ &= \|Q_k - Q^*\|_\infty. \end{aligned}$$

Thus,

$$\begin{aligned}
 & \|\mathbf{E}(F_k(s, a, o))\|_\infty = \\
 & = \gamma \max_{s, a, o} \left| \sum_{s_+} P(s, a, s_+) (\mathbf{E}(Q_k(s_+, \pi_k(s_+), \phi_k(s_+))) - \mathbf{E}(Q^*(s_+, \pi^*(s_+), \phi_k(s_+)))) \right| \\
 & \leq \gamma \max_{s, a, o} \sum_{s_+} P(s, a, s_+) |(\mathbf{E}(Q_k(s_+, \pi_k(s_+), \phi_k(s_+))) - \mathbf{E}(Q^*(s_+, \pi^*(s_+), \phi_k(s_+))))| \\
 & \leq \gamma \max_{s, a, o} \sum_{s_+} P(s, a, s_+) \|Q_k - Q^*\|_\infty \\
 & = \gamma \|Q_k - Q^*\|_\infty \\
 & = \gamma \|\Delta_k\|_\infty
 \end{aligned} \tag{30}$$

where the first inequality follows from the triangle inequality and the fact that $P(s, a, s_+) \geq 0$. Also, we have that

$$\begin{aligned}
 & \mathbf{E}(F_k - \mathbf{E}(F_k) | \mathcal{F}_k)^2 = \\
 & = \gamma^2 \mathbf{E} \left(Q_k(s_+, \pi_k(s_+), \phi_k(s_+)) - Q^*(s_+, \pi^*(s_+), \phi_k(s_+)) - \right. \\
 & \quad \left. - \sum_{s_+} P(s, a, s_+) (Q_k(s_+, \pi_k(s_+), \phi_k(s_+)) - Q^*(s_+, \pi^*(s_+), \phi_k(s_+))) \right)^2 \\
 & = \gamma^2 \mathbf{E} \left(\Delta_k(s_+, \pi_k(s_+), \phi_k(s_+)) - \right. \\
 & \quad \left. - \sum_{s_+} P(s, a, s_+) (\Delta_k(s_+, \pi_k(s_+), \phi_k(s_+))) \right)^2 \\
 & \leq C(1 + \|\Delta_k\|_\infty^2).
 \end{aligned}$$

Thus, $\Delta_k = Q_k - Q^*$ satisfies the conditions of Proposition 2 and hence converges to zero with probability 1, i. e. Q_k converges to Q^* with probability 1.

D Proof of Theorem 2

Lemma 1. *Let the operator \mathbf{T} be given by*

$$(\mathbf{T}Q)(s, a, o) = \sum_{s_+} P(s, a, s_+) \max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q(s_+, \pi(s_+), o))). \tag{31}$$

Then,

$$\|\mathbf{T}Q_1 - \mathbf{T}Q_2\|_\infty \leq \|Q_1 - Q_2\|_\infty.$$

Proof.

$$\begin{aligned}
 & \|\mathbf{T}Q_1 - \mathbf{T}Q_2\|_\infty = \\
 & = \max_{s, a, o} \left| \sum_{s_+} P(s, a, s_+) \left(\max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_1(s_+, \pi(s_+), o))) - \right. \right. \\
 & \quad \left. \left. - \max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_2(s_+, \pi(s_+), o))) \right) \right| \\
 & \leq \max_{s, a, o} \sum_{s_+} P(s, a, s_+) \left| \max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_1(s_+, \pi(s_+), o))) - \right. \\
 & \quad \left. - \max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_2(s_+, \pi(s_+), o))) \right|
 \end{aligned} \tag{32}$$

where the last inequality follows from the triangle inequality and the fact that $P(s, a, s_+) \geq 0$. Without loss of generality, assume that

$$\begin{aligned} & \max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_1(s_+, \pi(s_+), o))) \\ & \geq \max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_2(s_+, \pi(s_+), o))). \end{aligned}$$

Introduce

$$(\pi_i, o_i) = \arg \max_{\pi \in \Pi} \min_{o \in O} R(s, a, o) + \mathbf{E}(Q_i(s_+, \pi(s_+), o)).$$

Then,

$$\left| \max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_1(s_+, \pi(s_+), o))) - \right. \quad (33)$$

$$\left. - \max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_2(s_+, \pi(s_+), o))) \right|$$

$$= \max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_1(s_+, \pi(s_+), o))) - \quad (34)$$

$$- \max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_2(s_+, \pi(s_+), o)))$$

$$= R(s, a, o_1) + \mathbf{E}(Q_1(s_+, \pi_1(s_+), o_1)) - (R(s, a, o_2) + \mathbf{E}(Q_2(s_+, \pi_2(s_+), o_2)))$$

$$\leq R(s, a, o_2) + \mathbf{E}(Q_1(s_+, \pi_1(s_+), o_2)) - (R(s, a, o_2) + \mathbf{E}(Q_2(s_+, \pi_2(s_+), o_2)))$$

$$\leq R(s, a, o_2) + \mathbf{E}(Q_1(s_+, \pi_1(s_+), o_2)) - (R(s, a, o_2) + \mathbf{E}(Q_2(s_+, \pi_1(s_+), o_2)))$$

$$= |\mathbf{E}(Q_1(s_+, \pi_1(s_+), o_2)) - Q_2(s_+, \pi_1(s_+), o_2)|$$

$$\leq \max_{s_+, a, o} |Q_1(s_+, a, o) - Q_2(s_+, a, o)| \quad (35)$$

$$= \|Q_1 - Q_2\|_\infty. \quad (36)$$

Combining (32)–(36) implies that

$$\|\mathbf{T}Q_1 - \mathbf{T}Q_2\|_\infty \leq$$

$$\leq \max_{s, a, o} \sum_{s_+} P(s, a, s_+) \|Q_1 - Q_2\|_\infty \quad (37)$$

$$= \|Q_1 - Q_2\|_\infty$$

and the proof is complete. \square

Lemma 2. *The operator \mathbf{T} given by (31) is a span semi-norm, that is*

$$\|\mathbf{T}Q_1 - \mathbf{T}Q_2\|_s \leq \|Q_1 - Q_2\|_s \quad (38)$$

where

$$\|Q\|_s \triangleq \max_{s, a, o} Q(s, a, o) - \min_{s, a, o} Q(s, a, o).$$

Proof. We start off by noting the trivial inequalities

$$\max_{s', a', o'} (Q_1(s', a', o') - Q_2(s', a', o'))$$

$$\geq Q_1(s_+, a_+, o) - Q_2(s_+, a_+, o) \quad (39)$$

$$\geq \min_{s', a', o'} (Q_1(s', a', o') - Q_2(s', a', o')).$$

Also, let

$$o_i = \arg \min_{o \in O} R(s, a, o) + Q_i(s_+, \pi(s_+), o)$$

and

$$a_i = \arg \max_{a \in A} Q_i(s, a, o_j), \quad i \neq j.$$

The definition of the span semi-norm implies that

$$\begin{aligned}
 & \|\mathbf{T}Q_1 - \mathbf{T}Q_2\|_s = \\
 & = \left\| \sum_{s_+} P(s, a, s_+) \left(\max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_1(s_+, \pi(s_+), o))) - \right. \right. \\
 & \quad \left. \left. - \max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_2(s_+, \pi(s_+), o))) \right) \right\|_s \\
 & = \max_{s, a, o} \sum_{s_+} P(s, a, s_+) \left(\max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_1(s_+, \pi(s_+), o))) - \right. \\
 & \quad \left. - \max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_2(s_+, \pi(s_+), o))) \right) \\
 & \quad - \min_{s, a, o} \sum_{s_+} P(s, a, s_+) \left(\max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_1(s_+, \pi(s_+), o))) - \right. \\
 & \quad \left. - \max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_2(s_+, \pi(s_+), o))) \right) \\
 & \leq \max_{s, a, o} \sum_{s_+} P(s, a, s_+) \left(\max_{\pi \in \Pi} (R(s, a, o_2) + \mathbf{E}(Q_1(s_+, \pi(s_+), o_2))) - \right. \\
 & \quad \left. - \max_{\pi \in \Pi} (R(s, a, o_2) + \mathbf{E}(Q_2(s_+, \pi(s_+), o_2))) \right) \\
 & \quad - \min_{s, a, o} \sum_{s_+} P(s, a, s_+) \left(\max_{\pi \in \Pi} (R(s, a, o_1) + \mathbf{E}(Q_1(s_+, \pi(s_+), o_1))) - \right. \\
 & \quad \left. - \max_{\pi \in \Pi} (R(s, a, o_1) + \mathbf{E}(Q_2(s_+, \pi(s_+), o_1))) \right) \\
 & \leq \max_{s, a, o} \sum_{s_+} P(s, a, s_+) (Q_1(s_+, a_1, o_2) - Q_2(s_+, a_1, o_2)) \\
 & \quad - \min_{s, a, o} \sum_{s_+} P(s, a, s_+) (Q_1(s_+, a_2, o_1) - Q_2(s_+, a_2, o_1)) \\
 & \leq \max_{s, a, o} \sum_{s_+} P(s, a, s_+) \times \max_{s', a', o'} (Q_1(s', a', o') - Q_2(s', a', o')) \\
 & \quad - \min_{s, a, o} \sum_{s_+} P(s, a, s_+) \times \min_{s', a', o'} (Q_1(s', a', o') - Q_2(s', a', o')) \\
 & = \max_{s', a', o'} (Q_1(s', a', o') - Q_2(s', a', o')) - \min_{s', a', o'} (Q_1(s', a', o') - Q_2(s', a', o')) \\
 & = \|Q_1 - Q_2\|_s.
 \end{aligned} \tag{40}$$

□

For convenience, let $e : (s, a, o) \mapsto 1$ be a constant tensor with all elements equal to 1.

Lemma 3. Let $f \in \Phi$ be given, where the set Φ is defined as in Definition 2 and let

$$\mathbf{T}'(Q) = \mathbf{T}(Q) - f(Q) \cdot e$$

The ordinary differential equation (ODE)

$$\dot{Q}(t) = \mathbf{T}'(Q(t)) - Q(t) \tag{41}$$

has a unique globally asymptotically stable equilibrium Q^* , with $f(Q^*) = v^*$, where Q^* and v^* satisfy (17).

Proof. Introduce the operator

$$\widehat{\mathbf{T}}(Q) = \mathbf{T}(Q) - v \cdot e.$$

According to lemma [1](#) we have that

$$\|\mathbf{T}Q_1 - \mathbf{T}Q_2\|_\infty \leq \|Q_1 - Q_2\|_\infty$$

and hence, \mathbf{T} is Lipschitz. It's easy to verify that

$$\widehat{\mathbf{T}}(Q_1) - \widehat{\mathbf{T}}(Q_2) = \mathbf{T}(Q_1) - \mathbf{T}(Q_2)$$

and therefore

$$\begin{aligned} \|\widehat{\mathbf{T}}(Q_1) - \widehat{\mathbf{T}}(Q_2)\|_\infty &\leq \|Q_1 - Q_2\|_\infty, \\ \|\widehat{\mathbf{T}}(Q_1) - \widehat{\mathbf{T}}(Q_2)\|_s &\leq \|Q_1 - Q_2\|_s. \end{aligned}$$

Now consider the ODE:s

$$\dot{Q}(t) = \widehat{\mathbf{T}}(Q(t)) - Q(t) \tag{42}$$

and

$$\dot{Q}(t) = \mathbf{T}'(Q(t)) - Q(t) = \widehat{\mathbf{T}}(Q(t)) + (v - f(Q)) \cdot e. \tag{43}$$

Note that since \mathbf{T} and f are Lipschitz, the ODE:s [\(42\)](#) and [\(43\)](#) are well posed.

Since \mathbf{T} is Lipschitz and span semi-norm, the rest of the proof becomes identical to Theorem 3.4 along with Lemma 3.1, 3.2, and 3.3 in [\(Abounadi et al. 2001b\)](#) and hence omitted here. \square

Proposition 3 (Borkar & Meyn, 2000: Theorem 2.5). *Consider the asynchronous algorithm given by*

$$Q_{k+1} = Q_k + \alpha_k(h(Q_k) + M_{k+1})$$

where $\alpha_k(s, a, o) = 1_{(s,a,o)}(s_k, a_k, o_k) \times \beta_{N(k,s,a,o)}$. Suppose that

1. M_k is a martingale sequence with respect to the sigma algebra $\mathcal{F}_k = \sigma(Q_t, M_t, t \leq k)$, that is

$$\mathbf{E}(M_{k+1} \mid \mathcal{F}_k) = 0$$

and that there exists a constant $C_1 > 0$ such that

$$\mathbf{E}(\|M_{k+1}\|^2 \mid \mathcal{F}_k) \leq C_1(1 + \|Q_k\|^2).$$

2. Assumptions [4](#) and [5](#) hold.

3. The limit

$$h_\infty(X) = \lim_{z \rightarrow \infty} \frac{h(zX)}{z}$$

exists.

4. $\dot{Q}(t) = h(Q(t))$ has a unique globally asymptotically stable equilibrium Q^* .

Then, $Q_k \rightarrow Q^*$ with probability 1 as $k \rightarrow \infty$ for any initial value $Q(0)$.

Proof of Theorem [2](#). Introduce the operator

$$(\mathbf{T}Q)(s, a, o) = \sum_{s_+} P(s, a, s_+) \max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q(s_+, \pi(s_+), o))).$$

For convenience, let

$$\alpha_k(s, a, o) = 1_{(s,a,o)}(s_k, a_k, o_k) \cdot \beta_{N(k,s,a,o)},$$

$$M_{k+1}(s, a, o) = \max_{\pi \in \Pi} \min_{o \in O} (R(s, a, o) + \mathbf{E}(Q_k(s_{k+1}, \pi(s_{k+1}), o))) - (\mathbf{T}Q_k)(s, a, o),$$

and

$$h(Q) = \mathbf{T}Q - f(Q) \cdot e - Q.$$

Then,

$$Q_{k+1} = Q_k + \alpha_k(h(Q_k) + M_{k+1}).$$

We will now show that conditions 1 - 4 in Proposition [3](#) hold, and therefore $Q_k \rightarrow Q^*$ with probability 1, where Q^* is the solution to [\(17\)](#).

1. Let \mathcal{F}_k be the sigma algebra $\sigma(Q_t, M_t, t \leq k)$. Clearly,

$$\mathbf{E}(M_{k+1} \mid \mathcal{F}_k) = 0$$

and

$$\mathbf{E}(\|M_{k+1}\|^2 \mid \mathcal{F}_k) \leq C_1(1 + \|Q_k\|^2)$$

for some constant $C_1 > 0$.

2. We have supposed that assumptions [4](#) and [5](#) hold.

3. Let $h(X) = \mathbf{T}(X) - X - f(X) \cdot \mathbf{e}$ and introduce

$$(\bar{\mathbf{T}}Q)(s, a, o) = \max_{a_+ \in A} \sum_{s_+} P(s, a, s_+) Q(s_+, a_+, o). \quad (44)$$

Then, the limit

$$\begin{aligned} h_\infty(X) &= \lim_{z \rightarrow \infty} h(zX)/z \\ &= \bar{\mathbf{T}}(X) - X - f(X) \cdot \mathbf{e} \end{aligned}$$

exists.

4. By noting that

$$h(x) = \mathbf{T}(X) - X - f(X) \cdot \mathbf{e} = \mathbf{T}'(X) - X$$

we can apply Lemma [3](#) and conclude that $\dot{Q}(t) = h(Q(t))$ has a unique globally asymptotically stable equilibrium Q^* .

Thus, according to Proposition [3](#) the iterators Q_k in [18](#) converge to Q^* , where $h(Q^*) = 0$ and hence the unique solution to [17](#). Thus, the policy $\pi^* \in \Pi$ given by

$$\pi^*(s) = \arg \max_{\pi} \min_{o \in O} Q^*(s, \pi(s), o)$$

maximizes [13](#), and the proof is complete. \square

E Proof of Theorem [3](#)

Let

$$\mathcal{L}(\pi, j) = \mathbf{E} \left(\sum_{k=0}^{\infty} \gamma^k r^j(s_k, \pi(s_k)) \right).$$

Consider the zero-sum game

$$\max_{\pi \in \Pi} \min_{j \in [J]} \mathcal{L}(\pi, j).$$

Suppose that π is a policy such that

$$\mathbf{E} \left(\sum_{k=0}^{\infty} \gamma^k r^j(s_k, \pi(s_k)) \right) < 0$$

for some j . Then,

$$\mathcal{L}(\pi, j) < 0$$

which implies

$$\min_{j \in [J]} \mathcal{L}(\pi, j) < 0.$$

Thus, if

$$\max_{\pi \in \Pi} \min_{j \in [J]} \mathcal{L}(\pi, j) \geq 0$$

then, there must exist a policy π that satisfies

$$\mathbf{E} \left(\sum_{k=0}^{\infty} \gamma^k r^j(s_k, \pi(s_k)) \right) \geq 0 \quad (45)$$

for all j , and we get

$$\min_{j \in [J]} \mathcal{L}(\pi, j) \geq 0.$$

On the other hand, suppose that

$$\max_{\pi \in \Pi} \min_{j \in [J]} \mathcal{L}(\pi, j) < 0.$$

Then, there doesn't exist a policy π such that

$$\mathbf{E} \left(\sum_{k=0}^{\infty} \gamma^k r^j(s_k, \pi(s_k)) \right) \geq 0$$

for all j , because it would imply that

$$\max_{\pi \in \Pi} \min_{j \in [J]} \mathcal{L}(\pi, j) \geq 0$$

which is a contradiction, and the proof is complete.

F Proof of Theorem 5

Let

$$\mathcal{L}(\pi, j) = \lim_{T \rightarrow \infty} \mathbf{E} \left(\frac{1}{T} \sum_{k=0}^{T-1} r^j(s_k, \pi(s_k)) \right)$$

where the expectation is taken over s_k and π . The rest of the proof is similar to the proof of Theorem 3

G Proof of Theorem 6

According to Theorem 5 (23) is equivalent to the zero-sum Markov-Bandit game (24), which is equivalent to the zero-sum Markov-Bandit game given by the tuple (S, A, O, P, R) with the objective

$$\max_{\pi \in \Pi} \min_{o \in O} \lim_{T \rightarrow \infty} \mathbf{E} \left(\frac{1}{T} \sum_{k=0}^{T-1} R(s_k, \pi(s_k), o) \right). \quad (46)$$

Assumption 3 implies that $|R(s, a, o)| \leq 2c$ for all $(s, a, o) \in S \times A \times O$. Now let Q^* be the solution to the maximin optimality equation (17). According to Theorem 2 Q_k in the recursion given by (18) converges to Q^* with probability 1 under Assumptions 2, 3, 4, and 5. By definition, the optimal policy π^* maximizes the expected average reward of the zero-sum Markov-Bandit game (46). Hence,

$$\pi^*(s) = \arg \max_{\pi \in \Pi} \min_{o \in O} \mathbf{E} (Q^*(s, \pi(s), o))$$

and the proof is complete.

H Simulations

In this section we will consider two additional examples for discounted rewards.

H.1 Static Process Example 1

In this subsection, we consider an example with 1 state (denoted as 1), 2 actions (denoted as 1, 2), and two constraints. Let the reward function for the two constraints, $r^j(s, a)$ be given as

$$r^1(1, 1) = 1 \quad r^1(1, 2) = -1 \quad r^2(1, 1) = -1 \quad r^2(1, 2) = 1 \quad (47)$$

The aim of this example is to find a feasible policy that satisfies the discounted constraints. We let $\gamma = \frac{1}{2}$ in this example. Since there is only a single state, we will ignore the first variable of state in the following. We note that the only stationary policy that satisfies the constraints in this example is $\pi(1) = \pi(2) = 0.5$ due to the symmetry of the two constraints. We will now illustrate that the proposed algorithm will achieve a feasible policy that satisfies the constraints.

First, we define the reward function $R(a, o)$ for Markov zero-sum Bandit Game, $a, o \in \{1, 2\}$ as

$$R(1, 1) = 1 \quad R(2, 1) = -1 \quad R(1, 2) = -1 \quad R(2, 2) = 1 \quad (48)$$

We let the initial value for the Q-function be 0 and assume that the action for $k = 0$ is 1. For the learning rate, we adopt $\alpha_k = \frac{1}{k+1}$. We also label the policy in time-step i as π_i . According to Theorem 4, we can use the update rule in Eq. (12) to obtain the feasible policy. For $k = 0$, we have

$$(\pi_1, o_0) = \arg \max_{\pi \in \Pi} \min_{o \in \mathcal{O}} Q_0(\pi_0(s), o) \quad (49)$$

Since $Q_0 = 0$ for all $(a, o) \in \mathcal{A} \times \mathcal{O}$ and then the objective is not dependent on π , any arbitrarily policy can be used. Let us choose π as a half-half policy such that $\pi_1(1) = \pi_1(2) = 0.5$ and assume $a_1 = 2$. Similarly, o_0 can be arbitrary and we assume $o_0 = 1$. We also let $a_0 = 1$. Using $a_0 = 1, o_0 = 1, \pi_1(1) = \pi_1(2) = 0.5$, the Q-table update is given as

$$\begin{aligned} Q_1(1, 1) &= (1 - \alpha_0(1, 1))Q_0(1, 1) + \alpha_0(1, 1)(R(1, 1) + \gamma \mathbf{E}(Q_0(\pi_0, 1))) \\ &= R(1, 1) = 1 \end{aligned} \quad (50)$$

At the end of $k = 0$, we get $Q_1(1, 1) = 1$ and $Q_1(1, 2) = Q_1(2, 1) = Q_1(2, 2) = 0$.

For $k = 1$, we have

$$(\pi_2, o_1) = \arg \max_{\pi \in \Pi} \min_{o \in \mathcal{O}} Q_1(s, \pi(s), o) \quad (51)$$

Since $Q_1(2, 1) = Q_1(2, 2) = 0$, the maxmin problem will again have result 0 whatever the policy π_2 is. Thus, we still assume that $\pi_2(1) = \pi_2(2) = 0.5$ and next action $a_2 = 1$. However, it follows that $o_1 = 2$ because $Q_1(1, 1) = 1$. Since $a_1 = 2, o_1 = 2, \pi_2(1) = \pi_2(2) = 0.5$, the Q-table update is

$$\begin{aligned} Q_2(2, 2) &= (1 - \alpha_1(2, 2))Q_1(2, 2) + \alpha_1(2, 2)(R(2, 2) + \gamma \mathbf{E}(Q_1(\pi_1, 2))) \\ &= 0.5 * 0 + 0.5 * (1 + 0.5 * 0) = 0.5 \end{aligned} \quad (52)$$

At the end of $k = 1$, we get $Q_2(1, 1) = 1, Q_2(2, 2) = 0.5$ and $Q_2(1, 2) = Q_2(2, 1) = 0$.

For $k = 2$, we have

$$(\pi_3, o_2) = \arg \max_{\pi \in \Pi} \min_{o \in \mathcal{O}} Q_2(s, \pi(s), o) \quad (53)$$

To solve this problem, it is equivalent to solve the following problem

$$\begin{aligned} \arg \max_z \quad & z \\ \text{s.t.} \quad & z \leq Q_2(s, \pi(s), o) \quad \text{for } o = 1, 2 \end{aligned} \quad (54)$$

Assume $\pi_3(1) = p, \pi_3(2) = 1 - p$, this is equivalent to solve the equation that $p * Q_2(1, 1) + (1 - p) * Q_2(1, 2) = p * Q_2(1, 2) + (1 - p) * Q_2(2, 2)$, which gives the result $\pi_3(1) = \frac{1}{3}$ and $\pi_3(2) = \frac{2}{3}$ and we assume the next action

$a_3 = 2$. Due to the equality in the above equation, o_2 can again be arbitrary and we assume $o_2 = 2$. Since $a_2 = 1$, the Q-table update is

$$\begin{aligned} Q_3(1, 2) &= (1 - \alpha_2(1, 2))Q_2(1, 2) + \alpha_2(1, 2)(R(1, 2) + \gamma \mathbf{E}(Q_2(\pi_2, 2))) \\ &= \frac{2}{3} * 0 + \frac{1}{3} * [-1 + 0.5 * (\frac{1}{3} * Q_2(1, 2) + \frac{2}{3} * Q_2(2, 2))] = -\frac{5}{18} \end{aligned} \quad (55)$$

At the end of $k = 2$, we get $Q_3(1, 1) = 1$, $Q_3(2, 2) = 0.5$ and $Q_3(1, 2) = -\frac{5}{18}$ and $Q_3(2, 2) = 0$.

For $k = 3$, we have

$$(\pi_4, o_3) = \arg \max_{\pi \in \Pi} \min_{o \in O} Q_3(s, \pi(s), o) \quad (56)$$

We need to solve the problem in the Equation (54) to get the result of π_4 and the result is $\pi_4(1) = \frac{7}{16}$ and $\pi_4(2) = \frac{9}{16}$ and o_3 can be arbitrary, thus we assume that $o_3 = 1$. Since $a_3 = 2$, the Q-table update is given as

$$\begin{aligned} Q_4(2, 1) &= (1 - \alpha_3(2, 1))Q_3(2, 1) + \alpha_3(2, 1)(R(2, 1) + \gamma \mathbf{E}(Q_3(1, 1))) \\ &= \frac{3}{4} * 0 + \frac{1}{4} * (-1 + 0.5 * (\frac{7}{16} * Q_3(1, 1) + \frac{9}{16} * Q_3(2, 1))) = -\frac{3}{8} - \frac{1}{4} = -\frac{25}{128} \end{aligned} \quad (57)$$

At the end of $k = 3$, we get $Q_4(1, 1) = 1$, $Q_4(2, 2) = 0.5$ and $Q_4(1, 2) = -\frac{5}{18}$ and $Q_4(2, 1) = -\frac{25}{128}$.

Based on these steps, we can keep on computing the update for Q-table. However, the computation is hard to do manually, and involves random choice of actions based on policy π . Thus, we simulate the performance of the algorithm and the Q-values $Q_k(i, j)$ for iterations k are depicted in Fig. 3

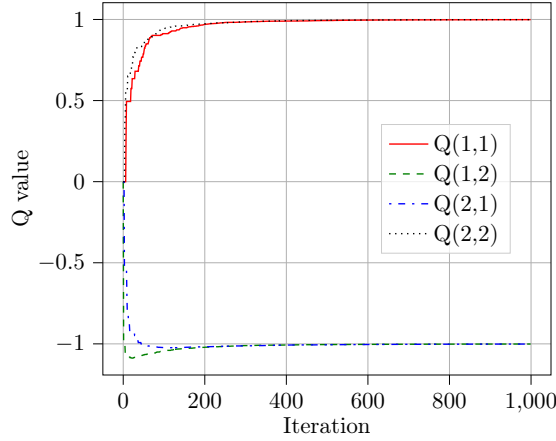


Figure 3: Convergence result for Example 1

We note that $Q_k(1, 1)$ and $Q_k(2, 2)$ converges to 1, while $Q_k(1, 2)$ and $Q_k(2, 1)$ converges to -1 . According to the optimal Bellman equation,

$$Q^*(s, a, o) = R(s, a, o) + \gamma \cdot \mathbf{E}(Q^*(s_+, \pi^*(s_+), o)) \quad (58)$$

we know $Q^*(1, 1) = 1 + 0.5 * [0.5 * Q^*(1, 1) + 0.5 * Q^*(2, 1)]$, which means

$$3Q^*(1, 1) = 4 + Q^*(2, 1) \quad (59)$$

Similarly, we have

$$3Q^*(2, 1) = -4 + Q^*(1, 1) \quad (60)$$

Combining these two equations, we have $Q^*(1, 1) = -Q^*(2, 1) = -1$. Similarly, $Q^*(2, 2) = -Q^*(1, 2) = -1$. Thus, we see that the algorithm successfully have the whole Q table converges to Q^* , which shows the correctness of the theorem. Moreover,

$$\pi^* = \arg \max_{\pi \in \Pi} \min_{o \in O} Q^*(s, \pi(s), o) \quad (61)$$

which gives $\pi^*(1|s) = \pi^*(2|s) = 0.5$ and we know this is the only feasible policy. Thus, we see that the Q-values of the proposed algorithm converges to that of the optimal policy and the policy converges to the only feasible policy in this example.

H.2 Static Process Example 2

We consider a static process (that is, the state is constant) and an agent that takes action from the action set $A = \{1, 2, 3\}$. There are three objectives given by the reward functions r_1, r_2 , and r_3 defined as

$$r^j(a) = \begin{cases} \frac{1}{2} & \text{if } a = j \\ 0 & \text{otherwise} \end{cases}$$

Note that we have dropped the dependence of the reward functions r_j on the state s as the state s is assumed to be constant. Let the discount factor be $\gamma = \frac{1}{2}$ and let

$$\alpha_0 = \alpha_1 = \alpha_2 = \alpha = \frac{1}{3}.$$

The agent would then be looking for a probability distribution over the set A , $\mathbf{Pr}(a)$ for $a \in A$, that simultaneously satisfies the objectives

$$\mathbf{E} \left(\sum_{k=0}^{\infty} \gamma^k r^j(a_k) \right) \geq \frac{1}{3}, \quad j = 1, 2, 3.$$

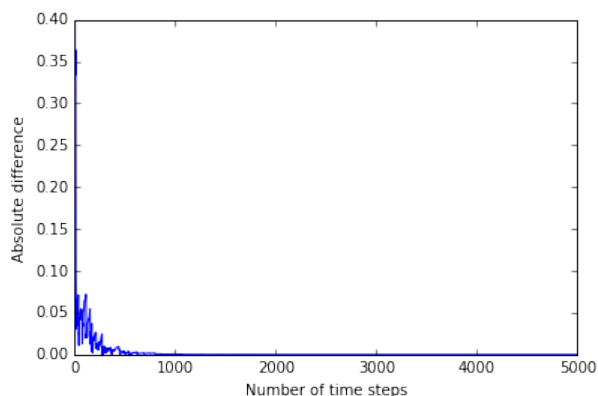


Figure 4: A plot of the maximum of $|p_1 - \hat{p}_1| + |p_2 - \hat{p}_2| + |p_3 - \hat{p}_3|$ over 1000 iterations, as a function of the number of time steps.

Now suppose that the agent takes action $a_k = 1$ with probability p_1 . Then we have that

$$\mathbf{E} \left(\sum_{k=0}^{\infty} \gamma^k r^1(a_k) \right) = p_1.$$

Similarly, we find that if the agent takes the action $a_k = j$ with probability p_j , $j = 2, 3$, then

$$\mathbf{E} \left(\sum_{k=0}^{\infty} \gamma^k r^j(a_k) \right) = p_j.$$

Without loss of generality, suppose that $p_1 \leq p_2 \leq p_3$. Now the equality $p_1 + p_2 + p_3 = 1$ together with the Arithmetic-Geometric Mean Inequality imply that

$$\frac{1}{3} = \frac{p_1 + p_2 + p_3}{3} \geq \sqrt[3]{p_1 p_2 p_3} \geq p_1$$

with equality if and only if $p_1 = p_2 = p_3 = \frac{1}{3}$. Thus, in order to satisfy all of the three objectives, the agent's mixed strategy is unique and given by $p_1 = p_2 = p_3 = \frac{1}{3}$.

We have run 1000 iterations of a simulation of the learning algorithm as given by Theorem 4 over 5000 time steps (with respect to the time index k). As the above calculations showed, the probability distribution of the optimal policy is given by $p_1 = p_2 = p_3 = \frac{1}{3}$. Let $\hat{p}_1, \hat{p}_2, \hat{p}_3$ be the estimated probabilities based on the Q -learning algorithm given by Theorem 4. In Figure 4 we see a plot of the maximum of the total error

$$|p_1 - \hat{p}_1| + |p_2 - \hat{p}_2| + |p_3 - \hat{p}_3|$$

over all iterations, as a function of the number of time steps. We see that it converges after 1000 time steps and stays stable for the rest of the simulation.