```
Algorithm 2 Infinite Q-learning with UCB-Hoeffding
    Initialized: \(Q(x, a) \leftarrow \frac{1}{1-\gamma}\) and \(N(x, a) \leftarrow 0\) for all \((x, a) \in \mathcal{S} \times \mathcal{A}\).
    Define \(\iota(k) \leftarrow \log (S A T(k+1)(k+2)), H \leftarrow \frac{\ln \left(2 /(1-\gamma) \Delta_{\min }\right)}{\ln (1 / \gamma)}, \alpha_{k}=\frac{H+1}{H+k}\).
    for step \(t \in[T]\) do
        Take action \(a_{t} \leftarrow \operatorname{argmax}_{a^{\prime}} Q\left(x_{t}, a^{\prime}\right)\), observe \(x_{t+1}\).
        \(k=N\left(x_{t}, a_{t}\right) \leftarrow N\left(x_{t}, a_{t}\right)+1\),
        \(b_{k} \leftarrow \frac{c_{2}}{1-\gamma} \sqrt{H \iota(k) / k}, \quad \triangleright c_{2}\) is a constant that can be set to \(4 \sqrt{2}\).
        \(\widehat{V}\left(x_{t+1}\right) \leftarrow \max _{a^{\prime} \in \mathcal{A}} \widehat{Q}\left(x_{t+1}, a^{\prime}\right)\),
        \(Q\left(x_{t}, a_{t}\right) \leftarrow\left(1-\alpha_{k}\right) Q\left(x_{t}, a_{t}\right)+\alpha_{k}\left[r\left(x_{t}, a_{t}\right)+b_{k}+\gamma \widehat{V}\left(x_{t+1}\right)\right]\),
        \(\widehat{Q}\left(x_{t}, a_{t}\right) \leftarrow \min \left\{\widehat{Q}\left(x_{t}, a_{t}\right), Q\left(x_{t}, a_{t}\right)\right\}\).
```


## 6 Algorithm for Discounted MDP

The pseudocode is listed in Algorithm 2. We acknowledge that Algorithm 2 relies on knowing a lower bound on $\Delta_{\text {min }}$, and we leave it an open problem to develop a parameter-free algorithm.

## 7 Proofs for Discounted MDP

Notations Let $Q^{t}(s, a), \widehat{Q}^{t}(s, a), V^{t}(s), \widehat{V}^{t}(s), N^{t}(s, a)$ denote the value of $Q(s, a), \widehat{Q}(s, a), V(s), \widehat{V}(s), N(s, a)$ right before the $t$-th step, respectively. Let $\tau(s, a, i):=\max \left\{t: N^{t}(s, a)=i-1\right\}$ be the step $t$ at which $\left(x^{t}, a^{t}\right)=$ $(x, a)$ for the $i$-th time. We will abbreviate $N^{t}\left(x^{t}, a^{t}\right)$ for $n^{t}$ when no confusion can arise. $\alpha_{t}^{i}$ is defined same as that in the finite-horizon episodic setting.

Proof of Theorem 3.2 We shall decompose the regret of each step as the expected sum of discounted gaps using the exact same argument as $\mathrm{Eq}(1)$, where the expect runs over all the possible infinite-length trajectories ${ }^{5}$ taken by Algorithm 2:

$$
\begin{equation*}
\left(V^{*}-V^{\pi_{t}}\right)\left(s_{t}\right)=\mathbb{E}\left[\sum_{h=0}^{\infty} \gamma^{h} \Delta\left(x_{t+h}, a_{t+h}\right) \mid a_{t+h}=\pi_{t+h}\left(s_{t+h}\right)\right] \tag{18}
\end{equation*}
$$

Based on this expression, the expected total regret over first $T$ steps can be rewritten as

$$
\begin{align*}
\mathbb{E}[\operatorname{Regret}(T)] & =\mathbb{E}\left[\sum_{t=1}^{T}\left(V^{*}-V^{\pi_{t}}\right)\left(x_{t}\right)\right]=\mathbb{E}\left[\sum_{t=1}^{T} \mathbb{E}\left[\sum_{h=0}^{\infty} \gamma^{h} \Delta\left(x_{t+h}, a_{t+h}\right)\right]\right] \\
& =\mathbb{E}\left[\sum_{t=1}^{T} \sum_{h^{\prime}=t}^{\infty} \gamma^{h^{\prime}-t} \Delta\left(x_{h^{\prime}}, a_{h^{\prime}}\right)\right] \tag{19}
\end{align*}
$$

Our next lemma is borrowed from Wang et al. (2019), which shows that Algorithm 2 satisfies optimism and bounded learning error with high probability. By abuse of notation, we still use $\mathcal{E}_{\text {conc }}$ to denote the successful concentration event in this setting. Recall that Algorithm 2 specifies $\iota(t)=\log (S A T(t+1)(t+2))$ and $\beta_{t}=$ $\frac{c_{3}}{1-\gamma} \sqrt{\frac{H \iota(t)}{t}}$.
Lemma 7.1 (Bounded Learning Error). Under Algorithm 2, event $\mathcal{E}_{\text {conc }}$ occurs w.p. at least $1-\frac{1}{2 T}$ :

$$
\begin{aligned}
& \mathcal{E}_{\mathrm{conc}}:=\left\{\forall(x, a, t) \in \mathcal{S} \times \mathcal{A} \times \mathbb{N}_{+}: 0 \leq\left(\widehat{Q}^{t}-Q^{*}\right)(x, a) \leq\left(Q^{t}-Q^{*}\right)(x, a)\right. \\
&\left.\leq \frac{\alpha_{n^{t}}^{0}}{1-\gamma}+\sum_{i=1}^{n^{t}} \gamma \alpha_{n^{t}}^{i}\left(\widehat{V}^{\tau(x, a, i)}-V^{*}\right)\left(x_{\tau(x, a, i)}\right)+\beta_{n^{t}}\right\}
\end{aligned}
$$

[^0]Then we proceed to present an analog of Lemma 4.3 that bounds the weighted sum of learning error in the discounted setting.
Lemma 7.2 (Weighted Sum of Learning Errors). Under $\mathcal{E}_{\mathrm{conc}}$, for every $(C, w)$-sequence $\left\{w_{t}\right\}_{t \geq 1}$, the following holds.

$$
\begin{equation*}
\sum_{t \geq 1} w_{t}\left(\widehat{Q}^{t}-Q^{*}\right)\left(x_{t}, a_{t}\right) \leq \frac{\gamma^{H} C}{1-\gamma}+\mathcal{O}\left(\frac{\sqrt{w S A H C \iota(C)}+w S A}{(1-\gamma)^{2}}\right) \tag{20}
\end{equation*}
$$

Proof. Recall that Lemma 7.1 bounds the learning error $\left(\widehat{Q}^{t}-Q^{*}\right)\left(x_{t}, a_{t}\right)$ on $\mathcal{E}_{\text {conc }}$. Thus we have:

$$
\begin{align*}
\sum_{t \geq 1} w_{t} \frac{\alpha_{n^{t}}^{0}}{1-\gamma} & \leq \sum_{t \geq 1} \mathbb{I}\left[n^{t}=0\right] \cdot \frac{w}{1-\gamma}=\frac{S A w}{1-\gamma} ;  \tag{21}\\
\sum_{t \geq 1} w_{t} \beta_{n^{t}} & =\sum_{s, a} \sum_{i} w_{\tau(s, a, i)} \beta_{i}=\frac{c_{3} \sqrt{H}}{1-\gamma} \sum_{s, a} \sum_{i} w_{\tau(s, a, i)} \sqrt{\frac{\iota(i)}{i}} \\
& \leq \frac{c_{3} \sqrt{H}}{1-\gamma} \sum_{s, a} \sum_{i=1}^{C_{s, a / a} w} w \sqrt{\frac{\iota(C)}{i}} \leq \frac{2 c_{3} \sqrt{H}}{1-\gamma} \sum_{s, a} \sqrt{C_{s, a} w \iota(C)} \quad \quad\left(C_{s, a}:=\sum_{i \geq 1} w_{\tau(s, a, i)}\right) \\
& \leq \frac{2 c_{3}}{1-\gamma} \sqrt{S A H C w \iota(C)} ; \quad \text { (Cauchy-Schwartz inequality) } \tag{22}
\end{align*}
$$

Moreover,

$$
\begin{align*}
& \sum_{t \geq 1} w_{t} \sum_{i=1}^{n^{t}} \gamma \alpha_{n^{t}}^{i}\left(\widehat{V}^{\tau(x, a, i)}-V^{*}\right)\left(x_{\tau(x, a, i)}\right) \\
= & \gamma \sum_{t \geq 1}\left(\widehat{V}^{t}-V^{*}\right)\left(x_{t+1}\right) \sum_{i \geq n^{t}+1} w_{\tau\left(x_{t}, a_{t}, i\right)} \alpha_{i}^{n^{t}+1} \\
= & \gamma \sum_{t \geq 1}\left(\widehat{V}^{t+1}-V^{*}\right)\left(x_{t+1}\right) \sum_{i \geq n^{t}+1} w_{\tau\left(x_{t}, a_{t}, i\right)} \alpha_{i}^{n^{t}+1}+\gamma \sum_{t \geq 1} \sum_{i \geq n^{t}+1} w_{\tau\left(x_{t}, a_{t}, i\right)} \alpha_{i}^{n^{t}+1}\left(\widehat{V}^{t}-\widehat{V}^{t+1}\right)\left(x_{t}\right) . \tag{23}
\end{align*}
$$

We let

$$
\widetilde{w}_{t+1}:=\sum_{i \geq n^{t}+1} w_{\tau\left(x_{t}, a_{t}, i\right)} n_{i}^{n^{t}+1}
$$

and further simplify (23) to be

$$
\begin{equation*}
\gamma \sum_{t \geq 2} \widetilde{w}_{t}\left(\widehat{V}^{t}-V^{*}\right)\left(x_{t}\right)+\gamma \sum_{t \geq 1} \widetilde{w}_{t+1}\left(\widehat{V}^{t}-\widehat{V}^{t+1}\right)\left(x_{t}\right) \tag{24}
\end{equation*}
$$

For the first term of (24), we claim that $\left\{\widetilde{w}_{t}\right\}_{t \geq 2}$ is a $(C,(1+1 / H) w)$-sequence. This can be verified by a similar argument to Ineq (15). We also have

$$
\left(\widehat{V}^{t}-V^{*}\right)\left(x_{t}\right)=\widehat{V}^{t}\left(x_{t}\right)-V^{*}\left(x_{t}\right)=\widehat{Q}^{t}\left(x_{t}, a_{t}\right)-V^{*}\left(x_{t}\right) \leq \widehat{Q}^{t}\left(x_{t}, a_{t}\right)-Q^{*}\left(x_{t}, a_{t}\right)=\left(\widehat{Q}^{t}-Q^{*}\right)\left(x_{t}, a_{t}\right)
$$

Therefore, the first term can be upper bounded by

$$
\begin{equation*}
\gamma \sum_{t \geq 2} \widetilde{w}_{t}\left(\widehat{V}^{t}-V^{*}\right)\left(x_{t}\right) \leq \gamma \sum_{t \geq 2} \widetilde{w}_{t}\left(\widehat{Q}^{t}-Q^{*}\right)\left(x_{t}, a_{t}\right) \tag{25}
\end{equation*}
$$

For the second term of (24), we have the following observation:

$$
\begin{align*}
\gamma \sum_{t \geq 1} \widetilde{w}_{t+1}\left(\widehat{V}^{t}-\widehat{V}^{t+1}\right)\left(x_{t}\right) & \leq \gamma(1+1 / H) w \sum_{s} \sum_{t \geq 1}\left(\widehat{V}^{t}-\widehat{V}^{t+1}\right)(s) \\
& \leq \gamma(1+1 / H) w \sum_{s} \widehat{V}^{1}(s) \leq \frac{\gamma(1+1 / H) w S}{1-\gamma} \tag{26}
\end{align*}
$$

Plugging (25) and (26) back into (24), we obtain

$$
\begin{equation*}
\sum_{t \geq 1} w_{t} \sum_{i=1}^{n^{t}} \gamma \alpha_{n^{t}}^{i}\left(\widehat{V}^{\tau(x, a, i)}-V^{*}\right)\left(x_{\tau(x, a, i)}\right) \leq \gamma \sum_{t \geq 2} \widetilde{w}_{t}\left(\widehat{Q}^{t}-Q^{*}\right)\left(x_{t}, a_{t}\right)+\frac{\gamma(1+1 / H) w S}{1-\gamma} \tag{27}
\end{equation*}
$$

Finally, combining (21), (22) and (27) and Lemma 7.1, we conclude that

$$
\begin{align*}
& \sum_{t \geq 1} w_{t}\left(\widehat{Q}^{t}-Q^{*}\right)\left(x_{t}, a_{t}\right) \\
\leq & \sum_{t \geq 1} w_{t}\left(\frac{\alpha_{n^{t}}^{0}}{1-\gamma}+\beta_{n^{t}}+\gamma \sum_{i=1}^{n^{t}} \alpha_{n^{t}}^{i}\left(\widehat{V}^{\tau(s, a, i)}-V^{*}\right)\left(x_{\tau(s, a, i)+1}\right)\right) \\
\leq & \frac{S A w}{1-\gamma}+\frac{2 c_{3}}{1-\gamma} \sqrt{\operatorname{SAHCw\iota (C)}}+\frac{\gamma(1+1 / H) w S}{1-\gamma}+\gamma \sum_{t \geq 2} \widetilde{w}_{t}\left(\widehat{Q}^{t}-Q^{*}\right)\left(x_{t}, a_{t}\right) \tag{28}
\end{align*}
$$

(Lemma 7.1)

Note that the last term in Ineq (28) is another weighted sum of learning errors starting from step 2, where the weights form a $(C,(1+1 / H) w)$-sequence. We can therefore repeat this unrolling argument for $H$ times. Our choice of $H$ in Algorithm 2 guarantees not only the bounded blow-up factor of weights, but also sufficiently small contribution of learning error after step $H$. In particular, we define a family of weights: when $h=0$, $\left\{w_{t}^{(h)}\right\}_{t \geq h+1}=\left\{w_{t}\right\}_{t \geq 1}$ is a $(C, w)$ sequence; $\forall h \in[H]\left\{w_{t}^{(h)}\right\}_{t \geq h+1}$ is a $\left(C,(1+1 / H)^{h} w \leq e w\right)$ sequence. Note that our previous definition of $\widetilde{w}$ is exactly $w_{t}^{(1)}$.

$$
\begin{align*}
& \sum_{t \geq 1} w_{t}\left(\widehat{Q}^{t}-Q^{*}\right)\left(x_{t}, a_{t}\right) \\
\leq & \sum_{h=0}^{H} \gamma^{h} \mathcal{O}\left(\frac{\sqrt{(1+1 / H)^{h} w S A H C \iota(C)}+(1+1 / H)^{h} w S A}{1-\gamma}\right)+\gamma^{H} \sum_{t \geq H+1} w_{t}^{(H)}\left(\widehat{Q}^{t}-Q^{*}\right)\left(x_{t}, a_{t}\right) \\
\leq & \mathcal{O}\left(\frac{\sqrt{w S A H C \iota(C)}+w S A}{(1-\gamma)^{2}}\right)+\frac{\gamma^{H}}{1-\gamma} \sum_{t \geq H+1} w_{t}^{(H)} . \tag{29}
\end{align*}
$$

Using the fact that the weights after $H$ unrolling $\left\{w_{t}^{(H)}\right\}_{t \geq H+1}$ is a $\left(C,(1+1 / H)^{H} w \leq e w\right)$-sequence completes the proof.

Note that we have clarified in Ineq (3) that on $\mathcal{E}_{\text {conc }}$ where optimism holds, sub-optimality gaps can be bounded by clipped learning error of $Q$-function. Again we divide its range $\left[\Delta_{\min }, \frac{1}{1-\gamma}\right]$ into disjoint subintervals and bound the sum inside each subinterval independently.
Lemma 7.3. Let $N=\left\lceil\log _{2}\left(1 / \Delta_{\min }(1-\gamma)\right)\right\rceil$. On $\mathcal{E}_{\text {conc }}$, for every $n \in[N]$,

$$
\begin{aligned}
& C^{(n)}:=\left|\left\{t \in \mathbb{N}_{+}:\left(\widehat{Q}^{t}-Q^{*}\right)\left(x_{t}, a_{t}\right) \in\left[2^{n-1} \Delta_{\min }, 2^{n} \Delta_{\min }\right)\right\}\right| \\
& \leq \mathcal{O}\left(\frac{S A}{4^{n} \Delta_{\min }^{2}(1-\gamma)^{5}} \ln \left(\frac{S A T}{(1-\gamma) \Delta_{\min }}\right)\right)
\end{aligned}
$$

Again, based on Lemma 7.1, we prove Lemma 7.3 by choosing a particular sequence of weights.
Proof. For every $n \in[N]$, let

$$
\begin{equation*}
w_{t}^{(n)}:=\mathbb{I}\left[\left(\widehat{Q}^{t}-Q^{*}\right)\left(x_{t}, a_{t}\right) \in\left[2^{n-1} \Delta_{\min }, 2^{n} \Delta_{\min }\right)\right], \tag{30}
\end{equation*}
$$

then $C^{(n)}=\sum_{t=1}^{\infty} w_{t}^{(n)}$ and $\left\{w_{t}^{(n)}\right\}_{t \geq 1}$ is a $\left(C^{(n)}, 1\right)$-sequence. According to Lemma 7.2,

$$
\begin{aligned}
\left(2^{n-1} \Delta_{\min }\right) \cdot C^{(n)} & \leq \sum_{t \geq 1} w_{t}^{(n)}\left(\widehat{Q}^{t}-Q^{*}\right)\left(x_{t}, a_{t}\right) \\
& \leq \frac{\gamma^{H} C^{(n)}}{1-\gamma}+\mathcal{O}\left(\frac{\sqrt{S A H C^{(n)} \iota\left(C^{(n)}\right)}+S A}{(1-\gamma)^{2}}\right) \\
& =\frac{\Delta_{\min }}{2} C^{(n)}+\mathcal{O}\left(\frac{\sqrt{S A H C^{(n)} \iota\left(C^{(n)}\right)}+S A}{(1-\gamma)^{2}}\right) \cdot \quad\left(H=\frac{\ln \left(\frac{2}{\Delta_{\min }(1-\gamma)}\right)}{\ln (1 / \gamma)}\right)
\end{aligned}
$$

Now we proceed to solve the above inequality for $C^{(n)}$. For simplicity, let $\delta=2^{n-2} \Delta_{\min }$ and $C^{(n)}=S A C^{\prime}$. Then we have the following:

$$
\begin{align*}
\delta \cdot S A C^{\prime} & \leq\left(2^{n-1}-\frac{1}{2}\right) \Delta_{\min } C^{(n)} \leq \mathcal{O}\left(S A \frac{\sqrt{H C^{\prime} \iota\left(S A C^{\prime}\right)}+1}{(1-\gamma)^{2}}\right) \\
\delta C^{\prime} & \stackrel{(1}{\leq} \mathcal{O}\left(\frac{\sqrt{C^{\prime}}}{(1-\gamma)^{5 / 2}} \sqrt{\ln \left(S A T C^{\prime}\right) \ln \frac{1}{\Delta_{\min }(1-\gamma)}}\right) \\
C^{\prime} & \leq \mathcal{O}\left(\frac{1}{\delta^{2}(1-\gamma)^{5}} \ln \frac{1}{\Delta_{\min }(1-\gamma)} \ln \left(S A T C^{\prime}\right)\right) \tag{31}
\end{align*}
$$

where (1) comes from the definition $H=\frac{\ln \left(2 / \Delta_{\min }(1-\gamma)\right)}{\ln (1 / \gamma)}$. Solving Ineq (31) yields

$$
C^{\prime} \leq \frac{1}{\delta^{2}(1-\gamma)^{5}} \ln \left(\frac{S A T}{\Delta_{\min }(1-\gamma)}\right)
$$

Finally, substituting $C^{(n)}=S A C^{\prime}$ and $\delta=2^{n-2} \Delta_{\text {min }}$ yields the desired formula.
Proof of Theorem 3.2. We continue the calculation based on the regret decomposition in Eq (19). For every infinite-length trajectory traj $\in \mathcal{E}_{\text {conc }}$,

$$
\begin{align*}
\sum_{t=1}^{T} \sum_{h^{\prime}=t}^{\infty} \gamma^{h^{\prime}-t} \Delta\left(x_{h^{\prime}}, a_{h^{\prime}} \mid \text { traj }\right) & \stackrel{(2}{=} \sum_{h=1}^{\infty} \Delta\left(x_{h}, a_{h}\right) \sum_{t=1}^{\min \{T, h\}} \gamma^{t} \leq \frac{1}{1-\gamma} \sum_{h=1}^{\infty} \Delta\left(x_{h}, a_{h} \mid \text { traj }\right) \\
& \stackrel{(3)}{\leq} \frac{1}{1-\gamma} \sum_{t \geq 1} \operatorname{cip}\left[\left(\widehat{Q}^{t}-Q^{*}\right)\left(x_{t}, a_{t} \mid \text { traj }\right) \mid \Delta_{\min }\right] \\
& \stackrel{4}{\leq} \frac{1}{1-\gamma} \sum_{n=1}^{N} 2^{n} \Delta_{\min } C^{(n)} \\
& \stackrel{(5)}{\leq} \mathcal{O}\left(\frac{S A}{\Delta_{\min }(1-\gamma)^{6}} \ln \left(\frac{S A}{p \epsilon(1-\gamma) \Delta_{\min }}\right)\right) \tag{32}
\end{align*}
$$

For the above inequalities, (2) comes from an interchange of summations, (3) is by optimism of estimated $Q$-values, (4) is because we can add an outer summation over subintervals $n \in[N]$ and bound each of them by their maximum value times the number of steps inside. Finally, (5) follows directly from Lemma 7.3.

On the other hand, for trajectories outside of $\mathcal{E}_{\text {conc }}$, since sub-optimality gaps are upper bounded by $1 / 1-\gamma$, we have:

$$
\begin{equation*}
\sum_{t=1}^{T} \sum_{h^{\prime}=t}^{\infty} \gamma^{h^{\prime}-t} \Delta\left(x_{h^{\prime}}, a_{h^{\prime}} \mid \text { traj }\right) \leq \sum_{t=1}^{T} \sum_{h^{\prime}=t}^{\infty} \gamma^{h^{\prime}-t} \frac{1}{1-\gamma} \leq \frac{T}{(1-\gamma)^{2}} \tag{33}
\end{equation*}
$$

Therefore, combining Ineq (32) and (33) gives us

$$
\begin{align*}
\mathbb{E}[\operatorname{Regret}(T)] & =\mathbb{E}\left[\sum_{t=1}^{T} \sum_{h^{\prime}=t}^{\infty} \gamma^{h^{\prime}-t} \Delta\left(x_{h^{\prime}}, a_{h^{\prime}}\right)\right] \\
& \leq \mathbb{P}\left(\mathcal{E}_{\text {conc }}\right) \cdot \mathcal{O}\left(\frac{S A}{\Delta_{\min }(1-\gamma)^{6}} \ln \left(\frac{S A T}{(1-\gamma) \Delta_{\min }}\right)\right)+\mathbb{P}\left(\overline{\mathcal{E}_{\text {conc }}}\right) \cdot \frac{T}{(1-\gamma)^{2}} \\
& \leq \mathcal{O}\left(\frac{S A}{\Delta_{\min }(1-\gamma)^{6}} \ln \left(\frac{S A T}{(1-\gamma) \Delta_{\min }}\right)\right) \tag{34}
\end{align*}
$$

where the last step is comes from $\mathbb{P}\left(\overline{\mathcal{E}_{\text {conc }}}\right) \leq 1 / 2 T$. Ineq (34) is precisely the assertion of Theorem 3.2.

## 8 Difficulty in Applying Optimistic Surplus

The closest related work is by Simchowitz and Jamieson (2019) who proved the logarithmic regret bound for a model-based algorithm. Simchowitz and Jamieson (2019) introduced a novel property characterizing optimistic algorithms, which is called optimistic surplus and defined as

$$
\begin{equation*}
E_{k, h}(x, a):=Q_{h}^{k}(x, a)-\left[r_{h}(x, a)+P_{h}(x, a)^{\top} V_{h+1}^{k}\right] . \tag{35}
\end{equation*}
$$

Under model-based algorithm with bonus term $b_{h}^{k}$, surplus can be decomposed as follows, where $\widehat{P}$ is the estimated transition probability:

$$
E_{k, h}(x, a)=\left(\widehat{P}_{h}^{\top}(x, a)-P_{h}^{\top}(x, a)\right) V_{h+1}^{*}+\left(\hat{P}_{h}^{\top}(x, a)-P_{h}^{\top}(x, a)\right)\left(V_{h+1}^{k}-V_{h+1}^{*}\right)+b_{h}^{k}
$$

The analysis of model-based algorithms is to first bound the regret ( $V^{*}-V^{\pi_{k}}$ ) by a sum over surpluses that are clipped to zero whenever being smaller than some $\Delta$-related quantities, then combine the concentration argument and properties of specially-designed bonus terms $b_{h}^{k}$ to provide high probability bound for surpluses. However, for model-free algorithms, estimates of transition probabilities are no longer maintained, so $\widehat{P}_{h}$ is a one-hot vector reflecting only the current step's empirical sample drawn from the real next-state distribution. In this scenario, concentration argument of $(\widehat{P}-P)$ cannot give us $\log T$ regret.
Following the update rule of $Q$-learning with learning rate $\alpha_{i}$ and upper confidence bound $b_{i}$, the surplus becomes

$$
E_{k, h}(x, a)=\alpha_{t}^{0} H+\left(\sum_{i=1}^{t} \alpha_{t}^{i} V_{h+1}^{k_{i}}\left(x_{h+1}^{k_{i}}\right)-P_{h}(x, a)^{\top} V_{h+1}^{k}\right)+\sum_{i=1}^{t} \alpha_{t}^{i} b_{i}
$$

in which $t=n_{h}^{k}(x, a)$ is the number of times $(x, a)$ has been visited, and $\alpha_{t}^{i}=\alpha_{i} \prod_{j=i+1}^{t}\left(1-\alpha_{j}\right)$ is the equivalent weight associated with the $i$-th visit of pair $(x, a)$. This indicates that the surplus of an episode is closely correlated with estimates of value functions during previous episodes. The correlation makes the analysis more difficult. Therefore, we use a very different approach to analyze $Q$-learning in this paper.


[^0]:    ${ }^{5}$ For the convenience of analysis, when proving the upper bound we remove the constraint $t \in[T]$ in the for-loop of line 3. Instead, we allow the algorithm to take as many steps as we need, even yielding infinite-length trajectories.

