# Supplementary Material: Online Kernel Learning with a Near Optimal Sparsity Bound

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### A. Proof of Theorem 2

We here prove a lower bound on the number of support vectors to achieve the optimal regret bound.

First, we construct a set of *n* examples  $\mathcal{T}_1 = \{(\mathbf{x}_i, y_i)\}_{i=1}^n$ , where  $\langle \kappa(\mathbf{x}_i, \cdot), \kappa(\mathbf{x}_{j}, \cdot) \rangle_{\mathcal{H}_{\kappa}} = \delta_{ij}$  and  $y_i \in \{1, -1\}$ . To make the construction, consider the degree-d polynomial kernel  $\kappa(\mathbf{x}, \mathbf{y}) = (\mathbf{x}^T \mathbf{y})^d$  and an Euclidean space  $\mathbb{R}^m$ where m > n. Since m > n, we can find a set of orthonormal vectors  $\{\mathbf{x}_1, \ldots, \mathbf{x}_n\}$  in  $\mathbb{R}^m$  such that  $\mathbf{x}_i^T \mathbf{x}_j = 0$ when  $i \neq j$  and  $\mathbf{x}_i^T \mathbf{x}_i = 1$ . It is easy to verify that this construction satisfies our assumption  $\kappa(\mathbf{x}_i, \mathbf{x}_j) = \delta_{ij}$ . For the Gaussian kernel, when the distance between  $\mathbf{x}_i$  and  $\mathbf{x}_j$  are large enough, we also have  $\kappa(\mathbf{x}_i, \mathbf{x}_j) \approx \delta_{ij}$ .

Based on  $\mathcal{T}_1$ , we construct another set  $\mathcal{T}_2$ :  $(\mathbf{z}, u) \in \mathcal{T}_2$  if there exist an index  $j \in [n]$  and a function  $\xi \in \mathcal{H}_{\kappa}$  such that

$$\kappa(\mathbf{z}, \cdot) = \kappa(\mathbf{x}_j, \cdot) + \xi, \ u = y_j, \text{ and } \langle \xi, \kappa(\mathbf{x}_i, \cdot) \rangle_{\mathcal{H}_{\kappa}} = 0, \forall i \in [n].$$
(13)

Thus, for each  $(\mathbf{z}, u) \in \mathcal{T}_2$ , there is a corresponding  $(\mathbf{x}_j, y_j) \in \mathcal{T}_1$  such that the relationships in (13) hold. The existence of  $\mathcal{T}_2$  can be proved in a similar way as that of  $\mathcal{T}_1$ .

Second, we select T distinct training examples  $(\mathbf{z}_1, u_1), \ldots, (\mathbf{z}_T, u_T)$  from  $\mathcal{T}_2$  such that, for each  $(\mathbf{x}, y) \in \mathcal{T}_1$  there are T/n examples constructed from it. Taking logit loss  $\ell(y,z) = \ln(1 + \exp(-yz))$  as an example. From the above constructions, it is easy to check that

$$f_* = \frac{R}{\sqrt{n}} \sum_{i=1}^n y_i \kappa(\mathbf{x}_i, \cdot)$$

minimizes the cumulative loss on the T training examples, i.e.,

$$f_* = \operatorname*{argmin}_{\|f\|_{\mathcal{H}_{\kappa}} \leq R} \sum_{i=1}^T \ell(u_i, f(\mathbf{z}_i)) = \operatorname*{argmin}_{\|f\|_{\mathcal{H}_{\kappa}} \leq R} \sum_{i=1}^T \ell(u_i, \langle f, \kappa(\mathbf{z}_i, \cdot) \rangle_{\mathcal{H}_{\kappa}}) = \operatorname*{argmin}_{\|f\|_{\mathcal{H}_{\kappa}} \leq R} \sum_{j=1}^n \sum_{k=1}^{T/n} \ell(y_j, \langle f, \kappa(\mathbf{x}_j, \cdot) + \xi_{jk} \rangle_{\mathcal{H}_{\kappa}}),$$

and the minimal loss is given by

$$\epsilon = T \ln(1 + \exp(-R/\sqrt{n})).$$

Choosing

$$R = \sqrt{n} \ln \frac{T}{n},$$

Xiaofei He

we have

$$\epsilon \le T \exp(-R/\sqrt{n}) = n$$

Thus, the optimal solution  $f_*$  has n support vectors and the associated loss is O(n).

Third, we examine the performance of Algorithm 1. From Theorem 1, we know that both the regret and the number of support vectors of OSKL is on the order of

$$O(\epsilon + R^2) = O(n[\ln T]^2).$$

Finally, we consider any algorithm that outputs a sequence of kernel classifiers  $f'_1, \ldots, f'_T$  with no more than n-1 support vectors. By our construction, this algorithm must misclassify at least T/n training examples, and the cumulative loss  $\sum_{i=1}^{T} \ell(u_i, f'_i(\mathbf{z}_i))$  must be *larger than* 

$$\frac{T}{n}\ln 2 = \Omega\left(\frac{T}{n}\right).$$

Recall that the cumulative loss of  $f_*$  is O(n). So, the regret of this algorithm is also larger than  $\Omega(T/n)$ , which is significantly worse than  $O(n[\ln T]^2)$  for large T.

### B. Proof of Lemma 1

We first state the Bernstein's inequality for martingales (Cesa-Bianchi & Lugosi, 2006), which lays the foundation of the main results.

**Theorem 3.** (Bernstein's inequality for martingales). Let  $X_1, \ldots, X_n$  be a bounded martingale difference sequence with respect to the filtration  $\mathcal{F} = (\mathcal{F}_i)_{1 \leq i \leq n}$  and with  $|X_i| \leq K$ . Let

$$S_i = \sum_{j=1}^i X_j$$

be the associated martingale. Denote the sum of the conditional variances by

$$\Sigma_n^2 = \sum_{t=1}^n \operatorname{E}\left[X_t^2 | \mathcal{F}_{t-1}\right].$$

Then, for all constants  $t, \nu > 0$ ,

$$\Pr\left[\max_{i=1,\dots,n} S_i > t \text{ and } \Sigma_n^2 \le \nu\right] \le \exp\left(-\frac{t^2}{2(\nu + Kt/3)}\right),$$

and therefore,

$$\Pr\left[\max_{i=1,\ldots,n} S_i > \sqrt{2\nu t} + \frac{2}{3}Kt \text{ and } \Sigma_n^2 \le \nu\right] \le e^{-t}.$$

Proof. Define martingale difference

$$X_t = GZ_t - |\ell'(y_t, f_t(\mathbf{x}_t))|,$$

and martingale  $\Lambda_T = \sum_{t=1}^T X_t$ . Define

$$K = \max_{t} |X_t| \le G.$$

Define the conditional variance  $\Sigma_T^2$  as

$$\Sigma_T^2 = \sum_{t=1}^T \mathcal{E}_{t-1} \left[ (GZ_t - |\ell'(y_t, f_t(\mathbf{x}_t))|)^2 \right] \le \sum_{t=1}^T G|\ell'(y_t, f_t(\mathbf{x}_t))| = GA_T.$$

Since  $A_T \leq 1$ , we have  $\Sigma_T^2 \leq \frac{G}{T}$ . Following Theorem 3, with probability at least  $1 - \delta$ , we have

$$\Lambda_T = \sum_{t=1}^T GZ_t - |\ell'(y_t, f_t(\mathbf{x}_t))| \le \sqrt{2\frac{G}{T}\ln\frac{1}{\delta}} + \frac{2}{3}G\ln\frac{1}{\delta} \le G\ln\frac{1}{\delta},$$

where the last inequality follows from the fact  $T \ge 18/[G\ln(1/\delta)]$ .

## C. Proof of Lemma 2

We use the same definitions of  $X_t$ ,  $\Lambda_T$ , K and  $\Sigma_T^2$  in Appendix B. Notice that  $A_T$  in the upper bound for  $\Sigma_T^2$  is a random variable, thus we cannot direct apply Theorem 3. To handle this challenge, we make use of the peeling process described in (Bartlett et al., 2005), and have

$$\Pr\left(\Lambda_T \ge 2\sqrt{GA_T\tau} + \frac{2}{3}K\tau\right)$$

$$= \Pr\left(\Lambda_T \ge 2\sqrt{GA_T\tau} + \frac{2}{3}K\tau, A_T \le G_1T\right)$$

$$= \Pr\left(\Lambda_T \ge 2\sqrt{GA_T\tau} + \frac{2}{3}K\tau, \Sigma_T^2 \le GA_T, A_T \le G_1T\right)$$

$$\leq \Pr\left(\Lambda_T \ge 2\sqrt{GA_T\tau} + \frac{2}{3}K\tau, \Sigma_T^2 \le GA_T, A_T \le \frac{1}{T}\right)$$

$$+ \sum_{i=1}^m \Pr\left(\Lambda_T \ge 2\sqrt{GA_T\tau} + \frac{2}{3}K\tau, \Sigma_T^2 \le GA_T, A_T \le \frac{1}{T}\right)$$

$$\leq \Pr\left(A_T \le 2\sqrt{GA_T\tau} + \frac{2}{3}K\tau, \Sigma_T^2 \le GA_T, \frac{2^{i-1}}{T} < A_T \le \frac{2^i}{T}\right)$$

$$\leq \Pr\left(A_T \le \frac{1}{T}\right) + \sum_{i=1}^m \Pr\left(\Lambda_T \ge \sqrt{2\frac{G2^i}{T}\tau} + \frac{2}{3}K\tau, \Sigma_T^2 \le \frac{G2^i}{T}\right)$$

$$\leq \Pr\left(A_T \le \frac{1}{T}\right) + me^{-\tau},$$

where  $m = \lceil \log_2(G_1T^2) \rceil$ , and the last step follows the Bernstein's inequality for martingales. We complete the proof by setting  $\tau = \ln(m/\delta)$ , and using the assumption  $A_T > 1/T$ .