

## 7. Proofs

### 7.1. Proof of Theorem 1

Observe that  $Y^* \in \mathcal{S}_{\text{psd}} \subset \mathcal{S}_{\text{nuclear}}$ , so it suffices to prove the theorem assuming  $\mathcal{S} = \mathcal{S}_{\text{nuclear}}$  in (3).

We need some additional notation. Suppose the size of the  $i$ -th cluster is  $K_i$ , and the rank- $r$  SVD of  $Y^*$  is  $U\Sigma U^\top$ . Note that  $UU^\top$  is a block diagonal matrix with  $r$  blocks such that the  $i$ -th block has size  $K_i \times K_i$  with all entries equal to  $\frac{1}{K_i} \leq \frac{1}{K}$ . We define the projections  $\mathcal{P}_T$  and  $\mathcal{P}_{T^\perp}$  by

$$\mathcal{P}_T Z = UU^\top Z + ZUU^\top - UU^\top ZUU^\top$$

and

$$\mathcal{P}_{T^\perp} Z = Z - \mathcal{P}_T Z.$$

Define the matrix  $W_{ij} := ((2A_{ij} - 1)B_{ij})_{i,j=1}^n \in \mathbb{R}^{n \times n}$  and the quantities  $\theta := \mathbb{E}[(1 - 2P_{ij})B_{ij}]$  and  $\rho := \mathbb{E}[B_{ij}^2] - \theta^2$ . Note that

$$\mathbb{E}W_{ij} = \mathbb{E}[\mathbb{E}[(2A_{ij} - 1)B_{ij}|P]] = (2Y_{ij}^* - 1) \mathbb{E}[(1 - 2P_{ij})B_{ij}] = (2Y_{ij}^* - 1) \theta$$

and

$$\text{Var}[W_{ij}] = \mathbb{E}[W_{ij}^2] - (\mathbb{E}W_{ij})^2 = \mathbb{E}[(2A_{ij} - 1)^2 B_{ij}^2] - (2Y_{ij}^* - 1)^2 \theta^2 = \rho.$$

Our proof requires two standard concentration results for the random matrix  $W$ .

**Lemma 1.** *If  $0 \leq W_{ij} \leq b_0$  almost surely for all  $i, j$  and the condition (6) holds, then with high probability, we have*

$$\|W - \mathbb{E}[W]\| \leq c_2 \left( b \log n + \sqrt{\rho n \log n} \right) \quad (12)$$

and

$$\|UU^\top (W - \mathbb{E}[W])\|_\infty \leq c_3 \frac{\sqrt{b^2 \log^2 n + \rho K \log n}}{K} \quad (13)$$

for some universal constants  $c_2, c_3$ .

We prove the lemma in Section 7.1.1 to follow. We now prove Theorem 1 assuming the two inequalities (12) and (13) in the lemma hold.

For any matrix  $Y$ , we define  $\Delta(Y) := \langle Y^* - Y, W \rangle$ . To prove the theorem, it suffices to show that  $\Delta(Y) > 0$  for all feasible  $Y$  of the program 2-4 with  $Y \neq Y^*$ . We rewrite  $\Delta(Y)$  as

$$\Delta(Y) = \langle \mathbb{E}W, Y^* - Y \rangle + \langle W - \mathbb{E}W, Y^* - Y \rangle. \quad (14)$$

We bound the two terms above. For any feasible  $Y$  obeying the constraint 4, the first term in (14) can be written as

$$\begin{aligned} \langle \mathbb{E}W, Y^* - Y \rangle &= \sum_{i,j} (2Y_{ij}^* - 1) \theta \cdot (Y_{ij}^* - Y_{ij}) \\ &= \theta \|Y^* - Y\|_1, \end{aligned} \quad (15)$$

where the last equality follows from  $0 \leq Y_{ij} \leq 1, \forall i, j$ .

On the other hand, if we let  $\lambda := c_2 (\log n + \sqrt{\rho n \log n})$ , then by (12) we have

$$\left\| \frac{1}{\lambda} \mathcal{P}_{T^\perp} (W - \mathbb{E}W) \right\| \leq \left\| \frac{1}{\lambda} (W - \mathbb{E}W) \right\| \leq 1.$$

This means  $UU^\top + \frac{1}{\lambda} \mathcal{P}_{T^\perp} (W - \mathbb{E}W)$  is a subgradient of the function  $f(X) = \|X\|_*$  at  $X = Y^*$ . Therefore, for any feasible  $Y$  we have

$$0 \geq \|Y\|_* - \|Y^*\|_* \geq \langle UU^\top + \frac{1}{\lambda} \mathcal{P}_{T^\perp} (W - \mathbb{E}W), Y - Y^* \rangle,$$

which means

$$\langle W - \mathbb{E}W, Y^* - Y \rangle \geq \langle \mathcal{P}_T (W - \mathbb{E}W) - \lambda UU^\top, Y^* - Y \rangle \quad (16)$$

We substitute (15) and (16) into (14) to obtain that for all feasible  $Y$ ,

$$\begin{aligned} \Delta(Y) &\geq \theta \|Y^* - Y\|_1 + \langle \mathcal{P}_T(W - \mathbb{E}W) - \lambda U U^\top, Y^* - Y \rangle \\ &\stackrel{(a)}{\geq} (\theta - \lambda \|U U^\top\|_\infty - \|\mathcal{P}_T(W - \mathbb{E}W)\|_\infty) \|Y^* - Y\|_1 \\ &\stackrel{(b)}{\geq} \left( \theta - \frac{\lambda}{K} - \|\mathcal{P}_T(W - \mathbb{E}W)\|_\infty \right) \|Y^* - Y\|_1, \end{aligned}$$

where (a) follows from the Holder's inequality and (b) follows from the structure of  $U$ . But by definition of  $\mathcal{P}_T$ , we have

$$\begin{aligned} \|\mathcal{P}_T(W - \mathbb{E}W)\|_\infty &\leq \|U U^\top (W - \mathbb{E}W)\|_\infty + \|(W - \mathbb{E}W) U U^\top\|_\infty + \|U U^\top (W - \mathbb{E}W) U U^\top\|_\infty \\ &\leq 3 \|U U^\top (W - \mathbb{E}W)\|_\infty \leq 3c_3 \frac{\sqrt{b^2 \log^2 n + K\rho \log n}}{K}, \end{aligned}$$

where the last inequality follows from (13). It follows that

$$\Delta(Y) \geq \left( \theta - \frac{c_2 (b \log n + \sqrt{\rho n \log n})}{K} - 3c_3 \frac{\sqrt{b^2 \log^2 n + K\rho \log n}}{K} \right) \|Y^* - Y\|_1.$$

If the condition 6 in the theorem holds, then the quantity inside the parenthesis is positive (note that  $\rho \leq \mathbb{E}[B_{ij}^2]$ ). This means  $\Delta(Y) > 0$  for all  $Y \neq Y^*$ , which proves the theorem.

### 7.1.1. PROOF OF LEMMA 1

Let  $e_i$  be the  $i$ -th standard basis vector in  $\mathbb{R}^n$ . For the first inequality in the lemma, note that

$$W - \mathbb{E}W = \sum_{i,j} (W_{ij} - \mathbb{E}W_{ij}) e_i e_j^\top,$$

which is the sum of  $n^2$  i.i.d. zero-mean matrix. We compute

$$\|(W_{ij} - \mathbb{E}W_{ij}) e_i e_j^\top\| = |W_{ij} - \mathbb{E}W_{ij}| \leq b$$

for all  $(i, j)$  and

$$\begin{aligned} \left\| \mathbb{E} \sum_{i,j} (W_{ij} - \mathbb{E}W_{ij})^2 e_j e_i^\top e_i e_j^\top \right\| &= \left\| \mathbb{E} \sum_{i,j} (W_{ij} - \mathbb{E}W_{ij})^2 e_i e_j^\top e_j e_i^\top \right\| \\ &= \rho \left\| \sum_{i,j} e_i e_i^\top \right\| = \rho n. \end{aligned}$$

Applying the matrix Bernstein inequality (Tropp, 2012) gives that w.h.p.

$$\|W - \mathbb{E}W\| \leq c_2 \left( b \log n + \sqrt{\rho n \log n} \right)$$

for some constant  $c_2$ .

We prove the second inequality. Fix  $(i, j)$ . Assume node  $i$  belongs to the cluster  $k$ . Then

$$(U U^\top (W - \mathbb{E}W))_{ij} = \frac{1}{K_k} \sum_{i' \in C_k^*} (W - \mathbb{E}W)_{i'j},$$

which is the average of  $K_k$  independent zero-mean random variables taking values in  $[-b, b]$  with variance bounded by  $\rho$ . Therefore, by standard Bernstein inequality, we know that for some constant  $c_3$  that

$$\left| (U U^\top (W - \mathbb{E}W))_{ij} \right| \leq \frac{1}{K_k} c_3 \left( b \log n + \sqrt{\rho K_k \log n} \right) \leq c_3 \frac{\sqrt{b^2 \log^2 n + K\rho \log n}}{K}, \text{ w.h.p.}$$

where the last inequality follows from  $K_k \geq K$ . The lemma follows from a union bound over all  $(i, j)$ .

## 7.2. Proof of Corollary 1

For the first part of the corollary, we only need to show that the condition (6) in Theorem 1 is satisfied. Take  $b_0 := 10 \log \frac{1}{\epsilon}$ . Note that under the assumption of the theorem, we have almost surely

$$B_{ij} = B_{ij}^{\text{MLE}} = \log \frac{1 - \bar{P}_{ij}}{\bar{P}_{ij}} \leq \min \left\{ \frac{1 - 2\bar{P}_{ij}}{\bar{P}_{ij}}, \log \frac{1}{\epsilon} \right\} \leq 10 (1 - 2\bar{P}_{ij}) \log \frac{1}{\epsilon}, \quad (17)$$

so  $B_{ij} \leq b_0$ . The condition (8) in the corollary statement implies that

$$\mathbb{E} \left[ \left( \frac{1}{2} - P_{ij} \right) B_{ij} \right] \geq c_1 \cdot \frac{n \log n}{K^2} \cdot \frac{b_0}{10} \geq c_1 \frac{b_0 \log n}{10 K}$$

since  $\bar{P}_{ij} \geq P_{ij}$  and  $K \leq n$ . On the other hand, the second term in the RHS of (6) can be upper bounded as follows:

$$\begin{aligned} c_0 \sqrt{\mathbb{E} [B_{ij}^2]} \frac{\sqrt{n \log n}}{K} &= c_0 \sqrt{\mathbb{E} \left[ \log \frac{1 - \bar{P}_{ij}}{\bar{P}_{ij}} \cdot \log \frac{1 - \bar{P}_{ij}}{\bar{P}_{ij}} \right]} \frac{\sqrt{n \log n}}{K} \\ &\stackrel{(a)}{\leq} 10c_0 \sqrt{\mathbb{E} \left[ (1 - 2\bar{P}_{ij}) \log \frac{1 - \bar{P}_{ij}}{\bar{P}_{ij}} \right]} \frac{\sqrt{n \log(1/\epsilon) \log n}}{K}, \\ &\stackrel{(b)}{\leq} \frac{c_0}{2} \mathbb{E} \left[ (1 - 2P_{ij}) \log \frac{1 - \bar{P}_{ij}}{\bar{P}_{ij}} \right] \\ &= \frac{c_0}{2} \mathbb{E} [(1 - 2P_{ij}) B_{ij}], \end{aligned}$$

where the inequality (a) follows from the previous bound (17), and (b) follows from the condition (8) in the corollary statement and  $\bar{P}_{ij} \geq P_{ij}$ . Combining the last two display equations proves that (6) is satisfied.

For the second part of the corollary, we note that  $\bar{P}_{ij} := \max \left\{ \frac{1}{16}, P_{ij} \right\} \geq \epsilon := \frac{1}{16}$ . The RHS of (8) is upper bounded by  $\log 16 \cdot c_1 \frac{n}{K^2} \log n$ . Because  $\log \frac{1-x}{x} \geq \frac{1}{10}(1-2x) = \frac{1}{5}(\frac{1}{2}-x)$  for all  $x \leq \frac{1}{2}$ , we have  $(\frac{1}{2} - \bar{P}_{ij}) \log \frac{1 - \bar{P}_{ij}}{\bar{P}_{ij}} \geq \frac{1}{5} (\frac{1}{2} - \bar{P}_{ij})^2 \geq \frac{1}{10} (\frac{1}{2} - P_{ij})^2$  almost surely. It follows that the LHS of (8) is lower bounded by  $\frac{1}{10} \mathbb{E} \left[ (\frac{1}{2} - P_{ij})^2 \right]$ . Under the condition 9, we conclude that (8) is satisfied.

## 7.3. Proof of Theorem 2

We prove the lemma using Fano's inequality. By Stirling's formula we have

$$|\mathcal{Y}| = \binom{n}{n/2} \geq 2^{n/2}.$$

Now suppose  $Y^*$  is sampled uniformly at random from  $\mathcal{Y}$ , and then  $P$  and  $A$  are generated according to our model. We have

$$\begin{aligned} I(A, P; Y^*) &= H(A, P) - H(A, P|Y^*) \\ &\leq \binom{n}{2} [H(A_{11}, P_{11}) - H(A_{11}, P_{11}|Y_{11}^*)] \\ &\leq n^2 I(A_{11}, P_{11}; Y_{11}^*), \end{aligned}$$

where the first inequality follows from symmetry. Let  $a = A_{11}, p = P_{11}$  and  $y = Y_{11}^*$ . We now compute

$$\begin{aligned} I(a, p; y) &= \mathbb{E}_{a,p,y} \left[ \log \frac{\mathbb{P}(a, p|y)}{\mathbb{P}(a, p)} \right] = \mathbb{E}_{a,p,y} \left[ \log \frac{\mathbb{P}(a|y, p) \mathbb{Q}(p)}{\mathbb{P}(a|p) \mathbb{Q}(p)} \right] \\ &= \mathbb{E}_{a,p,y} [\log \mathbb{P}(a|y, p)] - \mathbb{E}_{a,p,y} [\log \mathbb{P}(a|p)] \\ &= \mathbb{E}_{a,p,y} [\log \mathbb{P}(a|y, p)] - \mathbb{E}_p \left[ p \log \frac{1}{2} \right] - \mathbb{E}_p \left[ (1-p) \log \frac{1}{2} \right]. \end{aligned}$$

One verifies that

$$\mathbb{E}_{a,p,y} [\log \mathbb{P}(a|y,p)] = \mathbb{E}_{a,p,y} [\mathbb{E} [\log \mathbb{P}(a|y,p) | a, y]] = \mathbb{E}_p [p \log p] + \mathbb{E}_p [(1-p) \log(1-p)].$$

Combining the last three display equations gives

$$\begin{aligned} I(A, P; Y^*) &\leq n^2 \mathbb{E}_p [p \log 2p + (1-p) \log 2(1-p)] \\ &\leq n^2 \mathbb{E}_p [p(2p-1) + (1-p)(1-2p)] \\ &= n^2 \mathbb{E}_p [(1-2p)^2]. \end{aligned}$$

It follows that under the condition (10) of the theorem, we have

$$I(A, P; Y^*) \leq n^2 \cdot \frac{c'}{n} = c'n \leq \frac{1}{4} \log |\mathcal{Y}|$$

provided  $c'$  is sufficiently small. Applying Fano's inequality, we obtain that for any  $\hat{Y}$ ,

$$\mathbb{P} [\hat{Y}(A, P) \neq Y^*] \geq 1 - \frac{I(A, P; Y^*) + \log 2}{\log |\mathcal{Y}|} \geq \frac{1}{2},$$

where the probability is w.r.t. the randomness in  $Y^*$ ,  $P$  and  $A$ . Because the supremum is lower bounded by the average, we conclude that

$$\sup_{Y^* \in \mathcal{Y}} \mathbb{P} [\hat{Y}(A, P) \neq Y^*] \geq \frac{1}{2}.$$

Taking the infimum over all  $\hat{Y}$  proves the theorem.