A Theorem 2

The proof of this section and the lemmas of the next section follow from the proofs of Gabillon et al. [2012]. The modifications we have made to this proof correspond to the introduction of the function g_k which bounds the uncertainty s_k in order to make it simpler to introduce other models. We also introduce a sufficient condition on this bound, i.e. that it is monotonically decreasing in N in order to bound the arm pulls with respect to g_k^{-1} . Ultimately, this form of the theorem reduces the problem of of proving a regret bound to that of checking a few properties of the uncertainty model.

Theorem 2. Consider a bandit problem with horizon T and K arms. Let $U_k(t)$ and $L_k(t)$ be upper and lower bounds that hold for all times $t \leq T$ and all arms $k \leq K$ with probability $1 - \delta$. Finally, let g_k be a monotonically decreasing function such that $s_k(t) \leq g_k(N_k(t-1))$ and $\sum_k g_k^{-1}(H_{k\epsilon}) \leq T - K$. We can then bound the simple regret as

$$\Pr(R_{\Omega_T} \le \epsilon) \ge 1 - KT\delta. \tag{10}$$

Proof. We will first define the event \mathcal{E} such that on this event every mean is bounded by its associated bounds for all times t. More precisely we can write this as

$$\mathcal{E} = \{ \forall k \le K, \forall t \le T, L_k(t) \le \mu_k \le U_k(t) \}.$$

By definition, these bounds are given such that the probability of deviating from a single bound is δ . Using a union bound we can then bound the probability of remaining within all bounds as $\Pr(\mathcal{E}) \geq 1 - KT\delta$.

We will next condition on the event \mathcal{E} and assume regret of the form $R_{\Omega_T} > \epsilon$ in order to reach a contradiction. Upon reaching said contradiction we can then see that the simple regret must be bounded by ϵ with probability given by the probability of event \mathcal{E} , as stated above. As a result we need only show that a contradiction occurs.

We will now define $\tau = \arg \min_{t \leq T} B_{J(t)}(t)$ as the time at which the recommended arm attains the minimum bound, i.e. $\Omega_T = J(\tau)$ as defined in (8). Let $t_k \leq T$ be the last time at which arm k is pulled. Note that each arm must be pulled at least once due to the initialization phase. We can then show the following sequence of inequalities:

$$\min(0, s_k(t_k) - \Delta_k) + s_k(t_k) \ge B_{J(t_k)}(t_k)$$
 (a)

$$\geq B_{\Omega_T}(\tau)$$
 (b)

$$\geq R_{\Omega_T}$$
 (c)

$$> \epsilon$$
. (d)

Of these inequalities, (a) holds by Lemma B3, (c) holds by Lemma B1, and (d) holds by our assumption on the simple regret. The inequality (b) holds due to the definition Ω_T and time τ . Note, that we can also write the preceding inequality as two cases

$$s_k(t_k) > 2s_k(t_k) - \Delta_k > \epsilon, \quad \text{if } \Delta_k > s_k(t_k);$$

$$2s_k(t_k) - \Delta_k \ge s_k(t_k) > \epsilon, \quad \text{if } \Delta_k \le s_k(t_k)$$

This leads to the following bound on the confidence diameter,

$$s_k(t_k) > \max(\frac{1}{2}(\Delta_k + \epsilon), \epsilon) = H_{k\epsilon}$$

which can be obtained by a simple manipulation of the above equations. More precisely we can notice that in each case, $s_k(t_k)$ upper bounds both ϵ and $\frac{1}{2}(\Delta_k + \epsilon)$, and thus it obviously bounds their maximum.

Now, for any arm k we can consider the final number of arm pulls, which we can write as

$$N_k(T) = N_k(t_k - 1) + 1 \le g^{-1}(s_k(t_k)) + 1$$

$$< g^{-1}(H_{k\epsilon}) + 1.$$

This holds due to the definition of g as a monotonic decreasing function, and the fact that we pull each arm at least once during the initialization stage. Finally, by summing both sides with respect to k we can see that $\sum_{k} g^{-1}(H_{k\epsilon}) + K > T$, which contradicts our definition of g in the Theorem statement.

B Lemmas

In order to simplify notation in this section, we will first introduce $B(t) = \min_k B_k(t)$ as the minimizer over all gap indices for any time t. We will also note that this term can be rewritten as

$$B(t) = B_{J(t)}(t) = U_{j(t)}(t) - L_{J(t)}(t),$$

which holds due to the definitions of j(t) and J(t).

Lemma B1. For any sub-optimal arm $k \neq k^*$, any time $t \in \{1, \ldots, T\}$, and on event \mathcal{E} , the immediate regret of pulling that arm is upper bounded by the index quantity, i.e. $B_k(t) \geq R_k$.

Proof. We can start from the definition of the bound and expand this term as

$$B_k(t) = \max_{i \neq k} U_i(t) - L_k(t)$$
$$\geq \max_{i \neq k} \mu_i - \mu_k = \mu^* - \mu_k = R_k.$$

The first inequality holds due to the assumption of event \mathcal{E} , whereas the following equality holds since we are only considering sub-optimal arms, for which the best alternative arm is obviously the optimal arm. \Box

Lemma B2. For any time t let $k = a_t$ be the arm pulled, for which the following statements hold:

if
$$k = j(t)$$
, then $L_{j(t)}(t) \le L_{J(t)}(t)$,
if $k = J(t)$, then $U_{j(t)}(t) \le U_{J(t)}(t)$.

Proof. We can divide this proof into two cases based on which of the two arms is selected.

Case 1: let k = j(t) be the arm selected. We will then assume that $L_{j(t)}(t) > L_{J(t)}(t)$ and show that this is a contradiction. By definition of the arm selection rule we know that $s_{j(t)}(t) \ge s_{J(t)}(t)$, from which we can easily deduce that $U_{j(t)}(t) > U_{J(t)}(t)$ by way of our first assumption. As a result we can see that

$$B_{j(t)}(t) = \max_{\substack{j \neq j(t) \\ j \neq J(t)}} U_j(t) - L_{j(t)}(t)$$

$$< \max_{\substack{j \neq J(t) \\ j \neq J(t)}} U_j(t) - L_{J(t)}(t) = B_{J(t)}(t).$$

This inequality holds due to the fact that arm j(t) must necessarily have the highest upper bound over all arms. However, this contradicts the definition of J(t) and as a result it must hold that $L_{j(t)}(t) \leq L_{J(t)}(t)$.

Case 2: let k = J(t) be the arm selected. The proof follows the same format as that used for k = j(t). \Box

Corollary B2. If arm $k = a_t$ is pulled at time t, then the minimum index is bounded above by the uncertainty of arm k, or more precisely

$$B(t) \le s_k(t)$$

Proof. We know that k must be restricted to the set $\{j(t), J(t)\}$ by definition. We can then consider the case that k = j(t), and by Lemma B2 we know that this imposes an order on the lower bounds of each possible arm, allowing us to write

$$B(t) \le U_{j(t)}(t) - L_{j(t)}(t) = s_{j(t)}(t)$$

from which our corollary holds. We can then easily see that a similar argument holds for k = J(t) by ordering the upper bounds, again via Lemma B2.

Lemma B3. On event \mathcal{E} , for any time $t \in \{1, \ldots, T\}$, and for arm $k = a_t$ the following bound holds on the minimal gap,

$$B(t) \le \min(0, s_k(t) - \Delta_k) + s_k(t)$$

Proof. In order to prove this lemma we will consider a number of cases based on which of $k \in \{j(t), J(t)\}$ is selected and whether or not one or neither of these arms corresponds to the optimal arm k^* . Ultimately, this results in six cases, the first three of which we will present are based on selecting arm k = j(t). **Case 1:** consider $k^* = k = j(t)$. We can then see that the following sequence of inequalities holds,

$$\mu_{(2)} \stackrel{(a)}{\geq} \mu_{J(t)}(t) \stackrel{(b)}{\geq} L_{J(t)}(t) \stackrel{(c)}{\geq} L_{j(t)}(t) \stackrel{(d)}{\geq} \mu_k - s_k(t).$$

Here (b) and (d) follow directly from event \mathcal{E} and (c) follows from Lemma B2. Inequality (a) follows trivially from our assumption that $k = k^*$, as a result J(t) can only be as good as the 2nd-best arm. Using the definition of Δ_k and the fact that $k = k^*$, the above inequality yields

$$s_k(t) - (\mu_k - \mu_{(2)}) = s_k(t) - \Delta_k \ge 0$$

Therefore the min in the result of Lemma B3 vanishes and the result follows from Corollary B2.

Case 2: consider k = j(t) and $k^* = J(t)$. We can then write

$$B(t) = U_{j(t)}(t) - L_{J(t)}(t)$$

$$\leq \mu_{j(t)}(t) + s_{j(t)}(t) - \mu_{J(t)}(t) + s_{J(t)}(t)$$

$$\leq \mu_k - \mu^* + 2s_k(t)$$

where the first inequality holds from event \mathcal{E} , and the second holds because by definition the selected arm must have higher uncertainty. We can then simplify this as

$$= 2s_k(t) - \Delta_k$$

$$\leq \min(0, s_k(t) - \Delta_k) + s_k(t),$$

where the last step evokes Corollary B2.

Case 3: consider $k = j(t) \neq k^*$ and $J(t) \neq k^*$. We can then write the following sequence of inequalities,

$$\mu_{j(t)}(t) + s_{j(t)}(t) \stackrel{(a)}{\geq} U_{j(t)}(t) \stackrel{(b)}{\geq} U_{k^*}(t) \stackrel{(c)}{\geq} \mu^*.$$

Here (a) and (c) hold due to event \mathcal{E} and (b) holds since by definition j(t) has the highest upper bound other than J(t), which in turn is not the optimal arm by assumption in this case. By simplifying this expression we obtain $s_k(t) - \Delta_k \ge 0$, and hence the result follows from Corollary B2 as in Case 1.

Cases 4–6: consider k = J(t). The proofs for these three cases follow the same general form as the above cases and is omitted. Cases 1 through 6 cover all possible scenarios and prove Lemma B3.

Lemma B4. Consider a normally distributed random variable $X \sim \mathcal{N}(\mu, \sigma^2)$ and $\beta \geq 0$. The probability that X is within a radius of $\beta\sigma$ from its mean can then be written as

$$\Pr(|X - \mu| \le \beta\sigma) \ge 1 - e^{-\beta^2/2}.$$

Proof. Consider $Z \sim \mathcal{N}(0, 1)$. The probability that Z exceeds some positive bound c > 0 can be written

$$\Pr(Z > c) = \frac{e^{-c^2/2}}{\sqrt{2\pi}} \int_c^\infty e^{(c^2 - z^2)/2} dz$$
$$= \frac{e^{-c^2/2}}{\sqrt{2\pi}} \int_c^\infty e^{-(z-c)^2/2 - c(z-c)} dz$$
$$\leq \frac{e^{-c^2/2}}{\sqrt{2\pi}} \int_c^\infty e^{-(z-c)^2/2} dz = \frac{1}{2} e^{-c^2/2}$$

The inequality holds due to the fact that $e^{-c(z-c)} \leq 1$ for $z \geq c$. Using a union bound we can then bound both sides as $\Pr(|Z| > c) \leq e^{-c^2/2}$. Finally, by setting $Z = (X - \mu)/\sigma$ and $c = \beta$ we obtain the bound stated above.