
Supplementary Material: A Hybrid Approach for Probabilistic Inference using Random Projections

Appendix

Proposition. For any graphical model G and choice of parameters $M \geq 1$, $m \geq 1$, $f \in (0, \frac{1}{2})$, $p \in [0, 1]$, $\frac{1}{M} \sum_{k=1}^M \left(\frac{2}{p+1}\right)^m Z^{(k)}$ is an unbiased estimator of the partition function Z . Furthermore, the variance of this estimator is bounded from above by the following expressions:

$$\begin{aligned} \text{Var} \left[\frac{1}{M} \sum_{k=1}^M \left(\frac{2}{p+1}\right)^m Z^{(k)} \right] &\leq \frac{1}{M} \left(\frac{4(p + \frac{1}{2}(1-f)(1-p)^2)}{(p+1)^2} \right)^m Z^2 + \frac{1}{M} \left(\frac{2(p^2+1)}{(p+1)^2} \right)^m \sum_{x \in \mathcal{X}} w(x)^2 - \frac{Z^2}{M} \\ &\leq \frac{2}{M} \left(\frac{2(p^2+1)}{(p+1)^2} \right)^m Z^2 \end{aligned}$$

Preliminaries

Before we prove the proposition, let us start off by proving some preliminaries. The key idea here is to relate the expected value of certain quantities involving the smoothed ($p > 0$) XOR factor potentials $\phi_{\{A_i^{(k)}, b_i^{(k)}\}}(x_{\{A_i^{(k)}\}})$ to the hard ($p = 0$) XOR factor potentials $\psi_{\{A_i^{(k)}, b_i^{(k)}\}}(x_{\{A_i^{(k)}\}})$.

Define the random variable $y_k(x) = \prod_{i=1}^m \phi_{\{A_i^{(k)}, b_i^{(k)}\}}(x_{\{A_i^{(k)}\}}) = \prod_{i=1}^m \phi_i(x)$ to be the product over all of the smoothed ($p > 0$) XOR factor potentials, where we use $\phi_i(x)$ as short-hand notation. Let $\psi_i(x)$ be the hard XOR factor potential corresponding to $\phi_i(x)$ for $p = 0$. By construction, we have that $\phi_i(x) = p + (1-p)\psi_i(x)$.

Recall that we proved $\mathbb{E}[y_k(x)] = \left(\frac{p+1}{2}\right)^m$ in Section 3.1 as follows: $\mathbb{E}[y_k(x)] = \mathbb{E}[\prod_{i=1}^m \phi_i(x)] = \prod_{i=1}^m \mathbb{E}[p + (1-p)\psi_i(x)] = \prod_{i=1}^m (p + (1-p)\frac{1}{2}) = \left(\frac{p+1}{2}\right)^m$.

Similarly, we can prove that $\mathbb{E}[y_k(x)^2] = \left(\frac{p^2+1}{2}\right)^m$ as follows: $\mathbb{E}[y_k(x)^2] = \mathbb{E}[\prod_{i=1}^m \phi_i(x)^2] = \prod_{i=1}^m \mathbb{E}[\phi_i(x)^2] = \prod_{i=1}^m \mathbb{E}[(p + (1-p)\psi_i(x))^2] = \prod_{i=1}^m (p^2 + 2p(1-p)\frac{1}{2} + (1-p)^2\frac{1}{2}) = \left(\frac{p^2+1}{2}\right)^m$.

Given two elements $a, b \in \mathcal{X}$ with $a \neq b$, we can derive an expression for $\mathbb{E}[y_k(a)y_k(b)]$ as follows:

$$\begin{aligned} \mathbb{E}[y_k(a)y_k(b)] &= \mathbb{E} \left[\prod_{i=1}^m \phi_i(a)\phi_i(b) \right] \\ &= \prod_{i=1}^m \mathbb{E}[\phi_i(a)\phi_i(b)] \\ &= \prod_{i=1}^m \mathbb{E} \left[(p + (1-p)\psi_i(a))(p + (1-p)\psi_i(b)) \right] \\ &= \prod_{i=1}^m \left(p + (1-p)^2 \mathbb{E}[\psi_i(a)\psi_i(b)] \right) \\ &= \sum_{t=0}^m \binom{m}{t} p^{m-t} (1-p)^{2t} \mathbb{E} \left[\prod_{\text{any } t \text{ distinct indices}} \psi_{i_1}(a)\psi_{i_1}(b) \dots \psi_{i_t}(a)\psi_{i_t}(b) \right] \end{aligned}$$

where we used the fact that the expected value of the product over any t distinct indices $\{i_1, \dots, i_t\}$ is equal by symmetry, which allows us to simplify the last expression.

Proof

Proof. For convenience, define a scaling factor $r = \frac{2}{p+1}$.

The partition function of our randomly projected $G^{(k)}$ is $Z^{(k)} = \sum_{x \in \mathcal{X}} w(x) \prod_{i=1}^m \phi_{\{A_i^{(k)}, b_i^{(k)}\}}(x_{\{A_i^{(k)}\}}) = \sum_{x \in \mathcal{X}} w(x) y_k(x)$ with $\mathbb{E}[Z^{(k)}] = \sum_{x \in \mathcal{X}} w(x) \mathbb{E}[y_k(x)] = \left(\frac{p+1}{2}\right)^m Z$.

Using variance and expectation formulas, we have the following equalities:

$$\begin{aligned}
 \text{Var} \left[\frac{1}{M} \sum_{k=1}^M r^m Z^{(k)} \right] &= \frac{r^{2m}}{M^2} \sum_{k=1}^M \text{Var}[Z^{(k)}] \\
 &= \frac{r^{2m}}{M^2} \sum_{k=1}^M \text{Var} \left[\sum_{x \in \mathcal{X}} w(x) y_k(x) \right] \\
 &= \frac{r^{2m}}{M^2} \sum_{k=1}^M \left(\sum_{x \in \mathcal{X}} w(x)^2 \text{Var}[y_k(x)] + \sum_{x_i \in \mathcal{X}} \sum_{x_j \neq x_i} w(x_i) w(x_j) \text{Cov}[y_k(x_i), y_k(x_j)] \right) \\
 &= \frac{r^{2m}}{M^2} \sum_{k=1}^M \left(\sum_{x \in \mathcal{X}} w(x)^2 (\mathbb{E}[y_k(x)^2] - \mathbb{E}[y_k(x)]^2) \right. \\
 &\quad \left. + \sum_{x_i \in \mathcal{X}} \sum_{x_j \neq x_i} w(x_i) w(x_j) (\mathbb{E}[y_k(x_i) y_k(x_j)] - \mathbb{E}[y_k(x_i)] \mathbb{E}[y_k(x_j)]) \right) \\
 &= \frac{r^{2m}}{M^2} \sum_{k=1}^M \left(\left(\frac{p^2+1}{2} \right)^m \sum_{x \in \mathcal{X}} w(x)^2 + \sum_{x_i \in \mathcal{X}} \sum_{x_j \neq x_i} w(x_i) w(x_j) \mathbb{E}[(y_k(x_i) y_k(x_j))] \right) - \frac{Z^2}{M} \\
 &= \frac{r^{2m}}{M^2} \sum_{k=1}^M \sum_{x_i \in \mathcal{X}} \sum_{x_j \neq x_i} w(x_i) w(x_j) \mathbb{E}[(y_k(x_i) y_k(x_j))] + \frac{r^{2m}}{M} \left(\frac{p^2+1}{2} \right)^m \sum_{x \in \mathcal{X}} w(x)^2 - \frac{Z^2}{M}
 \end{aligned}$$

Let us focus on bounding the first term in the variance expression above, V_1 . Using the lemma proved earlier with $a = x_i, b = x_j$, and $x_j \neq x_i$, we have that

$$\begin{aligned}
 V_1 &= \frac{r^{2m}}{M^2} \sum_{k=1}^M \sum_{x_i \in \mathcal{X}} \sum_{x_j \neq x_i} w(x_i) w(x_j) \mathbb{E}[(y_k(x_i) y_k(x_j))] \\
 &= \frac{r^{2m}}{M} \sum_{x_i \in \mathcal{X}} \sum_{x_j \neq x_i} w(x_i) w(x_j) \left(\sum_{t=0}^m \binom{m}{t} p^{m-t} (1-p)^{2t} \mathbb{E} \left[\prod_{\text{any } t \text{ distinct indices}} \psi_{i_1}(x_i) \psi_{i_1}(x_j) \dots \psi_{i_t}(x_i) \psi_{i_t}(x_j) \right] \right) \\
 &= \frac{r^{2m}}{M} \sum_{t=0}^m \binom{m}{t} p^{m-t} (1-p)^{2t} \left(\sum_{x_i \in \mathcal{X}} \sum_{x_j \neq x_i} w(x_i) w(x_j) \mathbb{E} \left[\prod_{\text{any } t \text{ distinct indices}} \psi_{i_1}(x_i) \psi_{i_1}(x_j) \dots \psi_{i_t}(x_i) \psi_{i_t}(x_j) \right] \right)
 \end{aligned}$$

Hashing with Low-Density (Sparse) Parity Constraints

Recall that in the construction of our hash functions in Section 3.1, we let $A_{i,j} = 1$ if node j is included in the i^{th} XOR constraint and $A_{i,j} = 0$ otherwise, and we let b_i be the parity bit generated for the i^{th} XOR constraint. The process for randomly generating A, b ($A_{i,j}^{(k)} \stackrel{iid}{\sim} \text{Bernoulli}(f)$ and $b_i^{(k)} \stackrel{iid}{\sim} \text{Bernoulli}(\frac{1}{2})$) is designed to give us sparse parity constraints. Given the set of m XOR constraints, a configuration $x \in \mathcal{X}$ satisfies the set of XOR constraints if and only if $Ax \equiv b \pmod{2}$.

Adapting the proof of Theorem 3 in *Low-density Parity Constraints for Hashing-Based Discrete Integration* (Ermon et al., 2014) and letting $A_{\{i_1, \dots, i_t\}}$ and $b_{\{i_1, \dots, i_t\}}$ be the t rows of A and b , respectively, corresponding to the t parity constraints inside the expectation, the expression we have can be written as:

$$\begin{aligned}
 & \sum_{x_i \in \mathcal{X}} \sum_{x_j \neq x_i} w(x_i)w(x_j) \mathbb{E} \left[\prod_{\text{any } t \text{ distinct indices}} \psi_{i_1}(x_i)\psi_{i_1}(x_j) \dots \psi_{i_t}(x_i)\psi_{i_t}(x_j) \right] \\
 = & \sum_{x_i \in \mathcal{X}} w(x_i) \sum_{x_j \neq x_i} w(x_j) \mathbb{P}[A_{\{i_1, \dots, i_t\}}x_i + b_{\{i_1, \dots, i_t\}} = \vec{0}_t, A_{\{i_1, \dots, i_t\}}x_j + b_{\{i_1, \dots, i_t\}} = \vec{0}_t] \\
 = & 2^{-t} \sum_{x_i \in \mathcal{X}} w(x_i) \sum_{s=1}^n W_{x_i}(s) r^{(s,f)}(0, 0) \\
 = & 2^{-2t} \sum_{x_i \in \mathcal{X}} w(x_i) \sum_{s=1}^n W_{x_i}(s) \left(1 + (1 - 2f)^s \right)^t \\
 \leq & 2^{-2t} (2 - 2f)^t \sum_{x_i \in \mathcal{X}} w(x_i) \sum_{s=1}^n W_{x_i}(s) \\
 \leq & 2^{-t} (1 - f)^t Z^2
 \end{aligned}$$

where $W_{x_i}(s)$ is the sum of $w(x_j)$ for all x_j at Hamming distance s from x_i and we used the fact that $r^{(s,f)}(0, 0) \leq r^{(1,f)}(0, 0)$.

Proof (continued)

The bound for our first term is

$$\begin{aligned}
 V_1 &= \frac{r^{2m}}{M} \sum_{t=0}^m \binom{m}{t} p^{m-t} (1-p)^{2t} \left(\sum_{x_i \in \mathcal{X}} \sum_{x_j \neq x_i} w(x_i)w(x_j) \mathbb{E} \left[\prod_{\text{any } t \text{ distinct indices}} \psi_{i_1}(x_i)\psi_{i_1}(x_j) \dots \psi_{i_t}(x_i)\psi_{i_t}(x_j) \right] \right) \\
 &\leq \frac{r^{2m}}{M} \sum_{t=0}^m \binom{m}{t} p^{m-t} (1-p)^{2t} \left(2^{-t} (1-f)^t Z^2 \right) \\
 &= Z^2 \frac{r^{2m}}{M} \sum_{t=0}^m \binom{m}{t} p^{m-t} \left(\frac{1}{2} (1-f)(1-p)^2 \right)^t \\
 &= Z^2 \frac{r^{2m}}{M} \left(p + \frac{1}{2} (1-f)(1-p)^2 \right)^m \\
 &\leq Z^2 \frac{r^{2m}}{M} \left(p + \frac{1}{2} (1-p)^2 \right)^m \\
 &= Z^2 \frac{r^{2m}}{M} \left(\frac{p^2 + 1}{2} \right)^m
 \end{aligned}$$

The final bound for our variance becomes

$$\begin{aligned}
 \text{Var} \left[\frac{1}{M} \sum_{k=1}^M r^m Z^{(k)} \right] &\leq Z^2 \frac{r^{2m}}{M} \left(\frac{p^2 + 1}{2} \right)^m + \frac{r^{2m}}{M} \left(\frac{p^2 + 1}{2} \right)^m \sum_{x \in \mathcal{X}} w(x)^2 - \frac{Z^2}{M} \\
 &\leq Z^2 \frac{2}{M} r^{2m} \left(\frac{p^2 + 1}{2} \right)^m \\
 &= Z^2 \frac{2}{M} \left(\frac{2(p^2 + 1)}{(p + 1)^2} \right)^m
 \end{aligned}$$

□

References

Ermon, Stefano, Gomes, Carla P., Sabharwal, Ashish, and Selman, Bart. Low-density parity constraints for hashing-based discrete integration. In *Proc. of the 31st International Conference on Machine Learning (ICML)*, pp. 271–279, 2014.