

# Fast Approximate Spectral Clustering for Dynamic Networks: Supplementary Material

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## 1 Approximating HR as $h(\mathbf{L})\mathbf{R}$ .

We compute the product with each column  $\mathbf{r}_i$  of  $\mathbf{R}$  independently. To achieve this using Chebyshev polynomials (Shuman et al., 2011), one employs the equation

$$\mathbf{H}\mathbf{r}_i \approx h(\mathbf{L})\mathbf{r}_i = \sum_{m=0}^c a_m \mathcal{T}_m(\mathbf{L})\mathbf{r}_i,$$

where each  $\mathcal{T}_m(\mathbf{L})\mathbf{r}_i$  is computed based on the recursion

$$\mathcal{T}_m(\mathbf{L})\mathbf{r}_i = \left( \frac{4}{\lambda_{\max}} \mathbf{L} - 2\mathbf{I} \right) \mathcal{T}_{m-1}(\mathbf{L})\mathbf{r}_i - \mathcal{T}_{m-2}(\mathbf{L})\mathbf{r}_i,$$

having as initial conditions

$$\mathcal{T}_0(\mathbf{L})\mathbf{r}_i = \mathbf{r}_i \quad \text{and} \quad \mathcal{T}_1(\mathbf{L})\mathbf{r}_i = \left( \frac{2}{\lambda_{\max}} \mathbf{L} - \mathbf{I} \right) \mathbf{r}_i.$$

The constant  $a_m$  should be selected as  $a_m = \frac{2-\delta(m)}{c+1} \sum_{j=0}^c s\left(\frac{\lambda_{\max}}{2}(1 + \cos(\pi \frac{2j+1}{2(c+1)}))\right) \cos(c\pi \frac{2j+1}{2(c+1)})$ , where  $s(x) = \mathbf{1}_{\{x \leq \lambda_k\}}$  is a step-function. The total computational complexity amounts to  $dc$  matrix-vector multiplications with a sparse matrix containing  $m$  non-zero elements.

## 2 Proof of Lemma 3.1

*Proof.* Let  $\mathbf{X}_\Phi$  and  $\mathbf{X}_\Psi$  be respectively the SC and CSC clustering assignments. Moreover, we denote for compactness the additive error term by  $\mathbf{E} = \Psi - \Phi \mathbf{I}_{k \times d} \mathbf{Q}$ . We have that

$$\begin{aligned} C_\Psi &= \|\Phi - \mathbf{X}_\Psi \mathbf{X}_\Psi^\top \Phi\|_F \\ &= \|(\mathbf{I} - \mathbf{X}_\Psi \mathbf{X}_\Psi^\top)(\Psi - \mathbf{E})\|_F \\ &\leq \|(\mathbf{I} - \mathbf{X}_\Psi \mathbf{X}_\Psi^\top)\Psi\|_F + \|(\mathbf{I} - \mathbf{X}_\Psi \mathbf{X}_\Psi^\top)\mathbf{E}\|_F \\ &\leq \|(\mathbf{I} - \mathbf{X}_\Psi \mathbf{X}_\Psi^\top)\Psi\|_F + \|\mathbf{E}\|_F \\ &\leq \|(\mathbf{I} - \mathbf{X}_\Phi \mathbf{X}_\Phi^\top)\Psi\|_F + \|\mathbf{E}\|_F \\ &= \|(\mathbf{I} - \mathbf{X}_\Phi \mathbf{X}_\Phi^\top)(\Phi \mathbf{I}_{k \times d} \mathbf{Q} + \mathbf{E})\|_F + \|\mathbf{E}\|_F \\ &\leq \|(\mathbf{I} - \mathbf{X}_\Phi \mathbf{X}_\Phi^\top)\Phi \mathbf{I}_{k \times d} \mathbf{Q}\|_F + 2\|\mathbf{E}\|_F \\ &= C_\Phi + 2\|\Psi - \Phi \mathbf{I}_{k \times d} \mathbf{Q}\|_F \end{aligned} \tag{1}$$

The lower bound comes from that  $\mathbf{X}_\Phi$  in eq. (1) defines the argmin of our cost functions, and thus  $C_\Phi \leq C_\Psi$ .  $\square$

### 3 Proof of Theorem 3.2

*Proof.* Let us start by noting that, by the unitary invariance of the Frobenius norm, for any  $k \times k$  matrix  $\mathbf{M}$

$$\|\Phi\mathbf{M}\|_F = \|\mathbf{U}\mathbf{I}_{n \times k}\mathbf{M}\|_F = \|\mathbf{I}_{n \times k}\mathbf{M}\|_F = \|\mathbf{M}\|_F. \quad (2)$$

We can thus rewrite the feature error as

$$\begin{aligned} \|\Psi - \Phi\mathbf{I}_{k \times d}\mathbf{Q}\|_F &= \|\Phi\Phi^\top\mathbf{R} - \Phi\mathbf{I}_{k \times d}\mathbf{Q}\|_F \\ &= \|\Phi^\top\mathbf{R} - \mathbf{I}_{k \times d}\mathbf{Q}\|_F \\ &= \|\mathbf{I}_{k \times n}\mathbf{U}^\top\mathbf{R} - \mathbf{I}_{k \times d}\mathbf{Q}\|_F \\ &= \|\mathbf{R}' - \mathbf{I}_{k \times d}\mathbf{Q}\|_F. \end{aligned} \quad (3)$$

We claim that there is a unitary matrix  $\mathbf{Q}$  that satisfies eq. (9). We describe this matrix as follows. Let  $\mathbf{R}' = \mathbf{Q}_L\Sigma\mathbf{Q}_R^\top$  be the singular value decomposition of  $\mathbf{R}'$  and set

$$\mathbf{Q} = \begin{pmatrix} \mathbf{Q}_L & 0 \\ 0 & \mathbf{I}_{d-k} \end{pmatrix} \mathbf{Q}_R^\top. \quad (4)$$

Substituting this to the feature error, we have that

$$\begin{aligned} \|\mathbf{R}' - \mathbf{I}_{k \times d}\mathbf{Q}\|_F &= \|\mathbf{Q}_L\Sigma\mathbf{Q}_R^\top - \mathbf{I}_{k \times d}\mathbf{Q}\|_F \\ &= \|\Sigma - \mathbf{Q}_L^\top\mathbf{I}_{k \times d}\mathbf{Q}\mathbf{Q}_R\|_F \\ &= \|\Sigma - \mathbf{Q}_L^\top\mathbf{I}_{k \times d} \begin{pmatrix} \mathbf{Q}_L & 0 \\ 0 & \mathbf{I}_{d-k} \end{pmatrix} \mathbf{Q}_R^\top\mathbf{Q}_R\|_F \\ &= \|\Sigma - \mathbf{Q}_L^\top \begin{pmatrix} \mathbf{Q}_L & 0 \end{pmatrix}\|_F \\ &= \|\Sigma - \mathbf{I}_{k \times d}\|_F, \end{aligned} \quad (5)$$

which is the claimed result.  $\square$

### 4 Proof of Corollary 3.2

*Proof.* To obtain the following extremal inequality for the singular values of  $\mathbf{R}'$ , we note that  $\mathbf{R}'$  is composed of i.i.d. Gaussian random variables with zero mean and variance  $1/d$ , and thus use Cor. 3.1 setting  $\mathbf{R}' = \mathbf{N}/d$  providing for every  $i$ ,

$$\sigma_i(\mathbf{R}') = \sigma_i(\mathbf{N})/\sqrt{d} \leq 1 + \frac{\sqrt{k} + \varepsilon}{\sqrt{d}}. \quad (6)$$

By simple algebraic manipulation, we then find that

$$\begin{aligned} \|\Sigma - \mathbf{I}_{k \times d}\|_F^2 &= \sum_{i=1}^k (\sigma_i(\mathbf{R}') - 1)^2 \\ &\leq k \left( \frac{\sqrt{k} + \varepsilon}{\sqrt{d}} \right)^2 = \frac{k}{d} (\sqrt{k} + \varepsilon)^2, \end{aligned} \quad (7)$$

which, after taking a square root, matches the claim.  $\square$

## 5 Relation Between Edge Similarity and Spectral Similarity

**Corollary 5.1** (adapted from Cor. 4 (Hunter and Strohmer, 2010)). *Let  $\mathbf{H}_{t-1}$  and  $\mathbf{H}_t$  be the orthogonal projection on to the span of  $[\mathbf{U}_k]_{t-1}(= \Phi_{t-1})$  and  $[\mathbf{U}_k]_t(= \Phi_t)$ . If there exists an  $\alpha > 0$  such that  $\alpha \leq \lambda_{k+1}^{(t-1)} - \lambda_k^t$  and  $\alpha \leq \lambda_k^t$ , then,*

$$\|\mathbf{H}_t - \mathbf{H}_{t-1}\|_F \leq \frac{\sqrt{2}}{\alpha} \|\mathbf{L}_t - \mathbf{L}_{t-1}\|_F. \quad (8)$$

Note that the bounds on  $\alpha$  are those described in their Thm. 3.

## References

- Hunter, B. and Strohmer, T. (2010). Performance analysis of spectral clustering on compressed, incomplete and inaccurate measurements. *arXiv preprint arXiv:1011.0997*.
- Shuman, D. I., Vandergheynst, P., and Frossard, P. (2011). Chebyshev polynomial approximation for distributed signal processing. In *2011 International Conference on Distributed Computing in Sensor Systems and Workshops (DCOSS)*, pages 1–8. IEEE.