## A. Supplement to "A Listwise Approach to Collaborative Ranking"

## A.1. Proofs in Theory section

*Proof of Theorem 1.* Notice that  $\Pi_{i1}$  is the argument, k, that minimizes  $Y_{ik}$ , and  $\mathbb{P}\{\Pi_{i1} = k\} = \mathbb{P}\{Y_{ik} \leq \min\{Y_{ij}\}_{j \neq k}\}$ . Furthermore,  $\min\{Y_{ij}\}_{j \neq k}$  is exponential with rate parameter  $\sum_{j \neq k} \phi(X_{ij})$  and is independent of  $Y_{ik}$ . Hence,

$$\mathbb{P}\left\{Y_{ik} \leq \min\{Y_{ij}\}_{j \neq k}\right\}$$
$$= \int_0^\infty \phi(X_{ik}) e^{-u\phi(X_{ik})} e^{-\sum_{j \neq k} u\phi(X_{ij})} du$$
$$= \frac{\phi(X_{ik})}{\sum_j \phi(X_{ij})}.$$

Furthermore,

$$\mathbb{P}\{\Pi_i | \Pi_{i1}\} = \mathbb{P}\{Y_{\Pi_{i_2}} \le \dots \le Y_{\Pi_{im}} | Y_{\Pi_{ij}} \ge Y_{\Pi_{i1}}, \forall j\}$$
  
=  $\mathbb{P}\{Y_{\Pi_{i2}} - Y_{\Pi_{i1}} \le \dots \le Y_{\Pi_{im}} - Y_{\Pi_{i1}} | Y_{\Pi_{ij}} \ge Y_{\Pi_{i1}}, \forall j\}$ 

By the memorylessness property, we have that the joint distribution of  $Y_{\Pi_{i2}} - Y_{\Pi_{i1}}, \ldots, Y_{\Pi_{im}} - Y_{\Pi_{i1}} | Y_{\Pi_{ij}} \ge Y_{\Pi_{i1}}, \forall j > 1$  is equivalent to the distribution of  $Y_{\Pi_{i2}}, \ldots, Y_{\Pi_{im}}$ . Hence, we can apply induction with the previous argument, and the tower property of conditional probability.  $\Box$ 

Proof of Lemma 1. By optimality,

$$\frac{1}{n}\sum_{i=1}^{n} -\log P_{\hat{X}_{i}}(\Pi_{i}) \leq \frac{1}{n}\sum_{i=1}^{n} -\log P_{X_{i}^{\star}}(\Pi_{i}).$$

Which is equivalent to

$$\frac{1}{n}\sum_{i=1}^{n} -\log \frac{P_{X_{i}^{\star}}(\Pi_{i})}{P_{\hat{X}_{i}}(\Pi_{i})} \geq 0.$$

Thus, we can subtract the expectation,

$$\begin{split} &\frac{1}{n}\sum_{i=1}^{n} \mathbb{E}\log\frac{P_{X_{i}^{\star}}(\Pi_{i})}{P_{\hat{X}_{i}}(\Pi_{i})} \\ &\leq -\frac{1}{n}\sum_{i=1}^{n}\left(\log\frac{P_{X_{i}^{\star}}(\Pi_{i})}{P_{\hat{X}_{i}}(\Pi_{i})} - \mathbb{E}\log\frac{P_{X_{i}^{\star}}(\Pi_{i})}{P_{\hat{X}_{i}}(\Pi_{i})}\right) \end{split}$$

where the expectation  $\mathbb{E}$  is with respect to the draw of  $\Pi_i$  conditional on X.

**Lemma 2.** Let  $\pi$  be a permutation vector and x be a score vector each of length m. Suppose that  $|\log \phi(x_j)| \le C$  for all j = 1, ..., m. Define the relative loss function,

$$L_{x,x'}(\pi) := \log \frac{P_x(\pi)}{P_{x'}(\pi)}$$

Consider translating an item in position  $\ell$  to position  $\ell'$ in the permutation  $\pi$ , thus forming  $\pi'$  where  $\pi'_{\ell'} = \pi_{\ell}$ . Specifically,  $\pi'_k = \pi_k$  if  $k < \min\{\ell, \ell'\}$  or  $k > \max\{\ell, \ell'\}$ ; if  $\ell < \ell'$  then  $\pi'_k = \pi_{k+1}$  for  $k = \ell, \ldots, \ell' - 1$ ; if  $\ell' < \ell$ then  $\pi'_k = \pi_{k-1}$  for  $k = \ell' + 1, \ldots, \ell$ . The relative loss function has bounded differences in the sense that

$$|L_{x,x'}(\pi) - L_{x,x'}(\pi')| \le C_0 \|\log \phi(x) - \log \phi(x')\|_{\infty},$$

where  $\phi(x)$  is applied elementwise and  $C_0 = 2 + e^{2C} \ln(m + 1)$ .

*Proof.* Suppose that  $\ell < \ell'$ , and define the following shorthand,

$$\lambda_j = \phi(x_{\pi_j}), \quad \Lambda_j = \sum_{k=j}^m \lambda_k,$$

and let  $\lambda'_j$ ,  $\Lambda'_j$  be similarly defined with x'. Then by replacing the permutation  $\pi$  with  $\pi'$  causes the  $\Lambda_j$  to be replaced with  $\Lambda_j - \lambda_j + \lambda_\ell$  for  $j = \ell + 1, \dots, \ell'$ . Hence,

$$\log Q_x(\pi) - \log Q_x(\pi')$$

$$= \sum_{j=\ell+1}^{\ell'} \log (\Lambda_j - \lambda_j + \lambda_\ell) - \log (\Lambda_j)$$

$$= \sum_{j=\ell+1}^{\ell'} \log (\Lambda_{j-1} + \lambda_\ell) - \log (\Lambda_j)$$

$$= \sum_{j=\ell}^{\ell'-1} \log \left(1 + \frac{\lambda_\ell}{\Lambda_j}\right) + \log \Lambda_\ell - \log \Lambda_{\ell'}$$

So we can bound the difference,

$$|L_{x,x'}(\pi) - L_{x,x'}(\pi')| \le \left|\log\frac{\Lambda_{\ell}}{\Lambda_{\ell}'} - \log\frac{\Lambda_{\ell'}}{\Lambda_{\ell'}'}\right| + \sum_{j=\ell}^{\ell'-1} \left|\log\left(1 + \frac{\lambda_{\ell}}{\Lambda_j}\right) - \log\left(1 + \frac{\lambda_{\ell}'}{\Lambda_j'}\right)\right|.$$

Suppose that for each j,  $|\log \lambda_j - \log \lambda'_j| \le \delta$  and that  $|\log \lambda_j| \le C$ . Then we have that

$$\log \frac{\Lambda_{\ell}}{\Lambda'_{\ell}} \le \max_{j \ge \ell} \left| \log \frac{\lambda_j}{\lambda'_j} \right| \le \delta.$$

The same equation can be made for this term with  $\ell'$ . Let

$$\alpha_j = \max\left\{\frac{\lambda_\ell}{\Lambda_j}, \frac{\lambda'_\ell}{\Lambda'_j}\right\}$$

Then we have

$$\begin{split} & \left| \log \left( 1 + \frac{\lambda_{\ell}}{\Lambda_j} \right) - \log \left( 1 + \frac{\lambda_{\ell}'}{\Lambda_j'} \right) \right| \\ & \leq \left| \log(1 + \alpha_j) - \log(1 + e^{-\delta} \alpha_j) \right| \leq \left| 1 - e^{-\delta} \right| |\alpha_j| \end{split}$$

Furthermore, because  $\Lambda_j \geq (m-j+1)e^{-C}$  then  $|\alpha_j| \leq (m-j+1)^{-1}e^{2C}$  and

$$\begin{aligned} |L_{x,x'}(\pi) - L_{x,x'}(\pi')| &\leq 2\delta \\ &+ \sum_{j=\ell}^{\ell'-1} |1 - e^{-\delta}| \frac{e^{2C}}{m - j + 1} \\ &\leq 2\delta + |1 - e^{-\delta}| e^{2C} H_m \\ &\leq \delta(2 + e^{2C} \ln(m + 1)). \end{aligned}$$

In the above equation,  $H_m$  is the *m*th harmonic number. A similar calculation can be made when  $\ell' > \ell$ . Setting  $\delta = \|\log \phi(x) - \log \phi(x')\|_{\infty}$  concludes the proof.  $\Box$ 

*Proof of Theorem 2.* Define the empirical process function to be

$$\rho_n(x) := \frac{1}{n} \sum_{i=1}^n \left( \log \frac{P_{X_i^\star}(\Pi_i)}{P_{X_i}(\Pi_i)} - \mathbb{E} \log \frac{P_{X_i^\star}(\Pi_i)}{P_{X_i}(\Pi_i)} \right).$$

By the listwise representation theorem,  $\rho_n(x)$  is a function of  $n \times m$  independent exponential random variables. Moreover, if we were to change the value of a single element  $y_{ik}$ then this would result in a change of permutation of the type described in Lemma 2. Notice that the bound on the change in the relative loss is  $C_0 \|\log \phi(X_i) - \log \phi(X'_i)\|_{\infty}$ , where  $C_0 = 2 + e^{2C} \ln(m+1)$ , and notice that the sum of squares of these bounds are,

$$\sum_{i,k} C_0^2 \|\log \phi(X_i) - \log \phi(X'_i)\|_{\infty}^2$$
  
=  $mC_0^2 \sum_{i=1}^n \|\log \phi(X_i) - \log \phi(X'_i)\|_{\infty}^2$   
=  $mC_0^2 \|Z - Z'\|_{\infty,2}^2$ ,

where Z, Z' are  $\log \phi$  applied elementwise to X, X' respectively. By Lemma 2 and McDiarmid's inequality,

$$\mathbb{P}\{n(\rho_n(x) - \rho_n(x')) > \epsilon\} \le \exp\left(-\frac{2\epsilon^2}{mC_0^2 \|Z - Z'\|_{\infty,2}^2}\right).$$

Hence, the stochastic process  $\{n\rho_n(X) : X \in \mathcal{X}\}$  is a subGaussian field with canonical distance,

$$d(X, X') := \sqrt{m}C_0 \|Z - Z'\|_{\infty, 2}$$

The result follows by Dudley's chaining (Talagrand, 2006).

**Lemma 3.** If  $\log \phi$  is 1-Lipschitz then we have that  $g(\mathcal{Z}) \leq g(\mathcal{X})$ .

*Proof.* Let  $X, X' \in \mathcal{X}$ , and  $Z = \log \phi(X), Z' = \log \phi(X')$ . Then

$$|z_{ij} - z'_{ij}| \le |x_{ij} - x'_{ij}|$$

by the Lipschitz property. Hence,  $||Z - Z'||_{\infty,2} \le ||X - X'||_{\infty,2}$ , and so

$$\mathcal{N}(u, \mathcal{Z}, \|.\|_{\infty, 2}) \leq \mathcal{N}(u, \mathcal{X}, \|.\|_{\infty, 2}).$$

*Proof of Corollary 1.* Consider two matrices X, X' in the model (7) then

$$||x_i - x'_i||_{\infty} = \max_j |\beta^\top z_{ij} - \beta'^\top z_{ij}| \le ||\beta - \beta'||_2 ||Z_i||_{2,\infty}.$$

Let  $\zeta = \max_i \|Z_i\|_{2,\infty}$  then

$$||X - X'||_{\infty,2} \le \zeta ||\beta - \beta'||_2.$$

The covering number of  $\mathcal{X}$  is therefore bounded by

$$\mathcal{N}(\mathcal{X}, u, \|.\|_{\infty, 2}) \le \mathcal{N}(B_{c_b}, u/\zeta, \|.\|_2) \le \left(\frac{\zeta C_0 c_b}{u}\right)^s,$$

for an absolute constant  $C_0$ , where  $B_{c_b}$  is the  $\ell_2$  ball of radius  $c_b$ . The result follows by Lemma 3 and Theorem 2.

Proof of Corollary 2. Let  $X = UV^{\top}$  and  $X' = U'V'^{\top}$  such that U, V, U', V' bounded by 1 in Frobenius norm (modifying the Frobenius norm bound does not change the substance of the proof). Consider

$$\begin{aligned} |u_i^{\top} v_j - u_i'^{\top} v_j'| &\leq |u_i^{\top} v_j - u_i'^{\top} v_j| + |u_i'^{\top} v_j - u_i'^{\top} v_j'| \\ &\leq ||u_i - u_i'|| ||v_j|| + ||u_i'|| ||v_j - v_j'||. \end{aligned}$$

Maximizing this over the selection of j,

$$\begin{split} \max_{j} &|u_{i}^{\top}v_{j} - u_{i}^{\prime\top}v_{j}^{\prime}| \\ &\leq \|u_{i} - u_{i}^{\prime}\|_{2}\|V\|_{2,\infty} + \|u_{i}\|_{2}\|V - V^{\prime}\|_{2,\infty} \end{split}$$

Hence,

$$\begin{aligned} \|X - X'\|_{\infty,2} \\ &\leq \|U - U'\|_F \|V\|_{2,\infty} + \|U'\|_F \|V - V'\|_{2,\infty} \\ &\leq \|U - U'\|_F + \|V - V'\|_{2,\infty}. \end{aligned}$$

Consider the vectorization mapping from the  $m \times r$  matrix to the mr dimensional vectors. The Frobenius norm is mapped to the  $\ell_2$  norm, and we can consider the  $2, \infty$  norm to be, the norm  $||x||_{\rho} = \max_j ||(x_{jr+1}, \ldots, x_{j(r+1)})||_2$ . The  $\rho$ -norm unit ball  $(B_{\rho})$  is just the Cartesian product of the  $\ell_2$  norm ball in K dimensions. The volume of a d-dimensional ball,  $V_d$ , is bounded by

$$C_l \le \frac{V_d}{(e\pi)^{d/2}} d^{\frac{d}{2}} \le C_u,$$

where  $C_l < C_u$  are universal constants. So the volume ratio between the  $\ell_2$  norm ball and the  $\rho$  norm ball is bounded by

$$\frac{V(B_2)}{V(B_{\rho})} \le C \left(\frac{r^{r/2}}{(e\pi)^{r/2}}\right)^m / \left(\frac{(rm)^{rm/2}}{(e\pi)^{rm/2}}\right) < Cm^{-rm/2},$$

where  $C = C_u/C_l$ .

$$\mathcal{N}(\epsilon, B_2, \|.\|_{\rho}) \le C_r \left(\frac{2}{\epsilon} + 1\right)^{rm} m^{-rm/2}$$
$$\le C \left(\frac{3}{\epsilon\sqrt{m}}\right)^{rm},$$

for  $\epsilon \leq 1$ . This is also the covering number of the Frobenius norm ball in the 2,  $\infty$  norm. Moreover, we know that the covering number of the unit Frobenius norm ball in  $n \times K$ matrices  $(B_F)$  in the Frobenius norm is

$$\mathcal{N}(\epsilon, B_F, \|.\|_F) \le \left(\frac{c}{\epsilon}\right)^{nr}$$

for some constant c. Consider covering the space  $\mathcal{X}$ , by selecting centers U, V from the  $\epsilon/2$ -coverings of  $B_F$  in the F-norm and  $2, \infty$  norm respectively. By the above norm bound, this produces an  $\epsilon$ -covering in the  $\infty, 2$  norm. Dudley's entropy bound is thus

$$\int_{0}^{\infty} \sqrt{\log \mathcal{N}(\epsilon, B_F, \|.\|_F) + \log \mathcal{N}(\epsilon, B_F, \|.\|_{2,\infty})} d\epsilon$$
$$\leq \int_{0}^{c} \sqrt{nr \log(c/\epsilon)} d\epsilon + \int_{0}^{3/\sqrt{m}} \sqrt{-mr \log(\epsilon\sqrt{m}/3)} d\epsilon$$

So that

$$\int_0^c \sqrt{nr \log(c/\epsilon)} \mathrm{d}\epsilon$$
$$\leq c' \sqrt{nr}$$

for some absolute constant c' and

$$\int_{0}^{3/\sqrt{m}} \sqrt{-mr \log(\epsilon \sqrt{m}/3)} \mathrm{d}\epsilon \le \int_{0}^{3} \sqrt{r \log(u/3)} \mathrm{d}u$$
$$\le c'\sqrt{r}.$$

Hence, for  $g(\mathcal{X}) \leq c' \sqrt{nr}$  and we have the result.

## A.2. Algorithms

Algorithm 3 Compute gradient for V when U fixed

**Input:**  $\Pi$ , U, V,  $\lambda$ ,  $\rho$ **Output:**  $g \{ g \in \mathbb{R}^{r \times m} \text{ is the gradient for } f(V) \}$  $g = \lambda \cdot V$ for i = 1 to n do Precompute  $h_t = u_i^T v_{\Pi_{it}}$  for  $1 \le t \le \bar{m}$  {For implicit feedback, it should be  $(1 + \rho) \cdot \tilde{m}$  instead of  $\tilde{m}$ , since  $\rho \cdot \tilde{m}$  0's are appended to the back} Initialize total = 0, tt = 0for  $t = \bar{m}$  to 1 do  $total += \exp(h_t)$ tt += 1/totalend for Initialize c[t] = 0 for  $1 \le t \le \overline{m}$ for  $t = \bar{m}$  to 1 do  $c[t] += h_t \cdot (1 - h_t)$  $c[t] += \exp(h_t) \cdot h_t \cdot (1-h_t) \cdot tt$  $total += \exp(h_t)$ tt = 1/totalend for for t = 1 to  $\bar{m}$  do  $q[:,\Pi_{it}] += c[t] \cdot u_i$ end for end for Return g

**Algorithm 4** Gradient update for V (Same procedure for updating U)

**Input:** *V*, *ss*, *rate* {*rate* refers to the decaying rate of the step size *ss*}

**Output:** V Compute gradient g for V {see alg 3}  $V = ss \cdot g$ ss \*= rate**Return** V



Figure 3. Comparing implicit feedback methods.



Figure 4. Effectiveness of Stochastic Queuing Process.



Figure 5. Effectiveness of using full lists.