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# Composable Core-sets for Determinant Maximization: A Simple Near-Optimal Algorithm

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## A. Missing Proofs

*Proof of Claim 5.5 of the paper.* We prove the claim by induction on  $t$ , and show that for any  $j$  s.t.  $j > t$ , the point  $\Pi(\mathcal{G}'_t)(q'_j)$  can be written as the sum  $\sum_{i \leq t} \alpha_i q'_i$  such that  $|\alpha_i| \leq 3^t$ .

**Base Case.** First, we prove the base case of induction, i.e.,  $t = 1$ . Recall that by our assumption,  $\|q_1\| > 3^k x$ , and thus by triangle inequality, we have that  $\|q'_1\| \geq \|q_1\| - x/k \geq 3^k x - x/k \geq 2x$ . Therefore, since  $q_1$  is the vector with largest norm in  $P$ , using triangle inequality again, we have that for any  $j > 1$ ,

$$\|q'_j\| \leq \|q_j\| \leq \|q_1\| \leq \|q'_1\| + x/k \leq (1 + \frac{1}{2k})\|q'_1\|$$

Therefore we can write  $\Pi(\mathcal{G}'_1)(q'_j) = \alpha_1 q'_1$  where  $|\alpha_1| \leq 2$ .

**Inductive step.** Now, lets assume that the hypothesis holds for  $\mathcal{G}'_t$ . In particular this means that we can write  $\Pi(\mathcal{G}'_t)(q'_{t+1}) = \sum_{i \leq t} \beta_i q'_i$  where  $|\beta_i| \leq 3^t$ , and that for a given  $j > t + 1$ , we can write  $\Pi(\mathcal{G}'_t)(q'_j) = \sum_{i \leq t} \gamma_i q'_i$  where  $|\gamma_i|$ 's are at most  $3^t$ . Now let  $\ell = \text{dist}(q'_{t+1}, \mathcal{G}'_t)$ . By

triangle inequality, we get that

$$\text{dist}(q_{t+1}, \mathcal{G}_t) \leq \text{dist}(q_{t+1}, q'_{t+1}) \tag{1}$$

$$+ \text{dist}(q'_{t+1}, \Pi(\mathcal{G}'_t)(q'_{t+1})) \tag{2}$$

$$\text{dist}(\Pi(\mathcal{G}'_t)(q'_{t+1}), \mathcal{G}_t)$$

$$\leq x/k + \ell + \text{dist}(\sum_{i \leq t} \beta_i q'_i, \sum_{i \leq t} \beta_i q_i)$$

$$\leq x/k + \ell + \sum_{i \leq t} |\beta_i| x/k$$

$$\leq \ell + 3^t x. \tag{3}$$

Now we consider two case. If  $\ell \leq 3^t x$  then using the above

$$\text{dist}(q_{t+1}, \mathcal{G}_t) \leq 2 \cdot 3^t x \leq 3^k x,$$

which contradicts our assumption of  $\text{dist}(q_{t+1}, \mathcal{G}_t) > 3^k x$ . Otherwise,

$$\begin{aligned} \text{dist}(\Pi(\mathcal{G}'_{t+1})(q'_j), \mathcal{G}'_t) &\leq \text{dist}(q'_j, \mathcal{G}'_t) \leq \text{dist}(q_j, \mathcal{G}_t) \\ &\leq \text{dist}(q_{t+1}, \mathcal{G}_t) \leq 2\ell, \end{aligned}$$

where the last inequality follows from Equation 1 . Therefore, we can write  $\Pi(\mathcal{G}'_{t+1})(q'_j) = \alpha_{t+1} q'_{t+1} - \alpha_{t+1} \Pi(\mathcal{G}'_t)(q'_{t+1}) + \Pi(\mathcal{G}'_t)(q'_j)$  where  $|\alpha_{t+1}| \leq 2$ .

By the hypothesis, we can write  $\Pi(\mathcal{G}'_t)(q'_j) = \sum_{i \leq t} \gamma_i q'_i$ , where  $|\gamma_i| \leq 3^t$ . Since  $|\alpha_{t+1}| \leq 2$ , we can write

$$\begin{aligned} \Pi(\mathcal{G}'_{t+1})(q'_j) &= \alpha_{t+1} q'_{t+1} + \sum_{i \leq t} (\gamma_i - \alpha_{t+1} \beta_i) q'_i \\ &= \sum_{i \leq t+1} \alpha_i q'_i \quad \text{where } |\alpha_i| \leq 3^{t+1}. \end{aligned}$$

This completes the proof of the claim. □

## B. Details of Experiments on Local Search vs. Greedy as offline algorithms

Here, we compare the quality of Greedy and Local Search as centralized algorithms on the whole data sets. Figure 1 shows the improvement ratio of the determinant of the solution returned by the Local Search algorithm over the determinant of the solution returned by the Greedy algorithm.

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On average over all values of  $k$ , Local Search improves over Greedy by 13% for GENES data set and 5% for MNIST data set. Figure 2 shows the ratio of the time it takes to run the Local Search and Greedy algorithms as a function of  $k$  for both data sets. On average, it takes about 6.5 times more to run the Local Search algorithm.

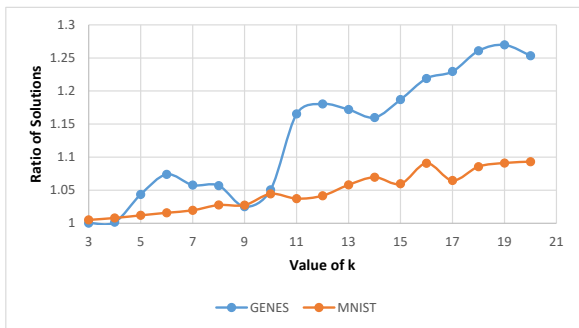


Figure 1. Average improvement of Local Search over Greedy as a function of  $k$ .

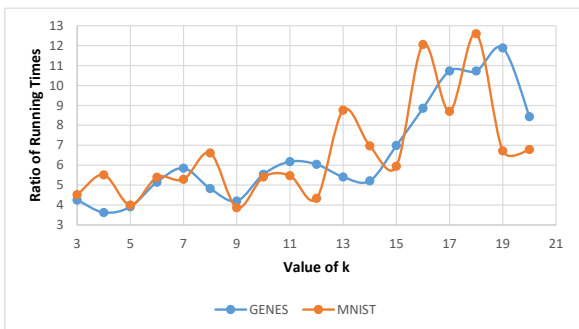


Figure 2. Average ratio of the run time of Local Search over Greedy as a function of  $k$ .

### C. Details of Experiments for Comparing Local Search vs. the LP-based Algorithm

In this section, we compare the performance of the Local Search algorithm and the LP-based algorithm of (Indyk et al., 2018) for constructing core-sets, i.e., we compare GD/LS with GD/LP. Figure 3 shows how much Local Search improves over the LP-based algorithm. On average this improvement is 7.3%, 1.8% and 1.4% for GENES, MNIST10 and MNIST50 respectively. Moreover, in 78% of all runs, Local Search performed better than Lp-based algorithm, and this improvement can go upto 63%. Figure 4 shows the average ratio of the time to construct core-sets using the LP-based algorithm vs. Local Search.

As it is clear from the graphs, our proposed Local Search algorithm performs better than even the LP-based algorithm which has almost tight approximation guarantees: while

picking fewer points in the core-set, in most cases it finds a better solution and runs faster.

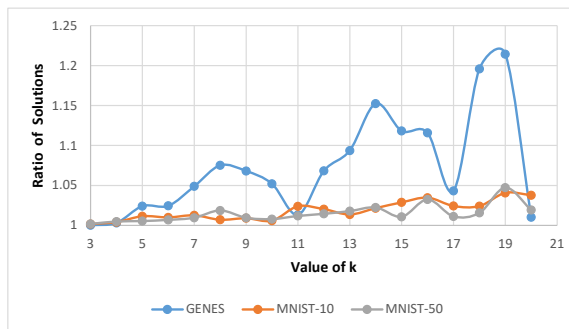


Figure 3. Average improvement of Local Search over LP-based algorithm for constructing core-sets as a function of  $k$ .

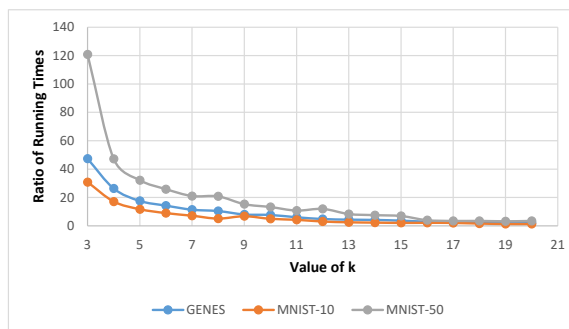


Figure 4. Average ratio of the run time of the optimal algorithm over local search as a function of  $k$ .

### References

Indyk, P., Mahabadi, S., Gharan, S. O., and Rezaei, A. Composable core-sets for determinant maximization problems via spectral spanners. *arXiv preprint arXiv:1807.11648*, 2018.