

# Exponential Convergence Time of Gradient Descent for One-Dimensional Deep Linear Neural Networks

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## Abstract

We study the dynamics of gradient descent on objective functions of the form  $f(\prod_{i=1}^k w_i)$  (with respect to scalar parameters  $w_1, \dots, w_k$ ), which arise in the context of training depth- $k$  linear neural networks. We prove that for standard random initializations, and under mild assumptions on  $f$ , the number of iterations required for convergence scales exponentially with the depth  $k$ . We also show empirically that this phenomenon can occur in higher dimensions, where each  $w_i$  is a matrix. This highlights a potential obstacle in understanding the convergence of gradient-based methods for deep linear neural networks, where  $k$  is large.

**Keywords:** Deep Learning, linear neural networks, gradient descent

## 1. Introduction

One of the biggest open problems in theoretical machine learning is to explain why deep artificial neural networks can be efficiently trained in practice, using simple gradient-based methods. Such training requires optimizing complex, highly non-convex objective functions, which seem intractable from a worst-case viewpoint. Over the past few years, much research has been devoted to this question, but it remains largely unanswered.

Trying to understand simpler versions of this question, significant attention has been devoted to *linear* neural networks, which are predictors mathematically defined as  $\mathbf{x} \mapsto \prod_{i=1}^k W_i \mathbf{x}$ , with  $W_1, \dots, W_k$  being a set of parameter matrices, and  $k$  being the depth parameter (e.g. [Saxe et al. \(2013\)](#); [Kawaguchi \(2016\)](#); [Hardt and Ma \(2016\)](#); [Lu and Kawaguchi \(2017\)](#); [Bartlett et al. \(2018\)](#); [Laurent and Brecht \(2018\)](#)). The optimization problem associated with training such networks can be formulated as

$$\min_{W_1, \dots, W_k} F(W_1, \dots, W_k) := f \left( \prod_{i=1}^k W_i \right) \quad (1)$$

for some function  $f$ . Although much simpler than general feedforward neural networks (which involve additional non-linear functions), it is widely believed that Eq. (1) captures important aspects of neural network optimization problems. Moreover, Eq. (1) has a simple algebraic structure, which makes it more amenable to analysis. In particular, it is known that when  $f$  is convex and differentiable, Eq. (1) has no local minima except global ones (see [Laurent and Brecht \(2018\)](#) and references therein). In other words, if an optimization algorithm converges to some local minimum, then it must converge to a global minimum.

Importantly, this no-local-minima result *does not* imply that gradient-based methods indeed solve Eq. (1) efficiently: Even when they converge to local minima (which is not always guaranteed, say in case the parameters diverge), the number of required iterations might be arbitrarily large. To study this question, [Bartlett et al. \(2018\)](#) recently considered the special case where  $F(W_1, \dots, W_k) := \frac{1}{2} \|\prod_{i=1}^k W_i - Y\|_F^2$  (where  $\|\cdot\|_F$  is the Frobenius norm) for square matrices  $W_1, \dots, W_k, Y$ , using gradient descent starting specifically from  $W_i = I$  for all  $i$ . In this setting, the authors prove a polynomial-time convergence guarantee when  $Y$  is positive semidefinite. On the other hand, when  $Y$  is symmetric and with negative eigenvalues, it is shown that gradient descent with this initialization will never converge. Although these results provide important insights, they crucially assume that each  $W_i$  is initialized *exactly* at the identity  $I$ . Since in practice parameters are initialized randomly, it is natural to ask whether such results hold with random initialization. Indeed, even though gradient descent might fail to converge with a specific initialization, it could be that even a tiny random perturbation is sufficient for polynomial-time convergence<sup>1</sup>. More recently, [Arora et al. \(2018a\)](#) considered gradient descent on a similar objective, and managed to prove strong polynomial-time convergence guarantees under certain assumptions about the initialization. However, as the authors discuss (in section 3.2.1), these assumptions are not generally satisfied for standard initialization approaches. In another recent related work, [Ji and Telgarsky \(2018\)](#) show that for certain classification problems on linearly separable data (corresponding to a suitable choice of  $f$  in Eq. (1)), gradient descent asymptotically converges to a globally optimal objective value. However, the result only applies to particular choices of  $f$ , and more importantly, is asymptotic and hence does not imply a finite-time convergence guarantee. Thus, analyzing the finite-time convergence of gradient descent on Eq. (1), with standard random initializations, remains a challenging open problem.

In this paper, we consider a simpler special case of Eq. (1), where the matrices  $W_1, \dots, W_k$  are all scalars:

$$\min_{\mathbf{w} \in \mathbb{R}^k} F(\mathbf{w}) := f\left(\prod_{i=1}^k w_i\right). \quad (2)$$

Our main and perhaps surprising result is that even in this relatively simple setting, gradient descent with random initialization can require  $\exp(\Omega(k))$  iterations to converge. This holds under mild conditions on the function  $f$ , and with standard initializations (including Xavier initialization and any reasonable initialization close to  $(1, \dots, 1)$ ). We complement this by showing that  $\exp(\tilde{O}(k)) \cdot \max\{1, \log(1/\epsilon)\}$  iterations are also sufficient for convergence to an  $\epsilon$ -optimal point. Moreover, in Sec. 4 we present experiments which strongly suggest that this phenomenon is not unique to one-dimensional networks, and at least in some cases, the same exponential dependence can also occur in multi-dimensional networks (i.e., Eq. (1) where each  $W_i$  is a  $d \times d$  matrix,  $d > 1$ ). The take-home message is that even if we focus on linear neural networks, natural objective functions without any spurious local minima, and random initializations, the associated optimization problems can sometime be intractable for gradient descent to solve, when the depth is large.

Before continuing, we emphasize that our results do not imply that gradient-based methods cannot learn deep linear networks. What they do imply is that one would need to make additional

1. For example, consider the objective  $F(w_1, w_2) = (w_1 w_2 + 1)^2$  where  $w_1, w_2 \in \mathbb{R}$ . It is an easy exercise to show that gradient descent starting from any  $w_1 = w_2 > 0$  (and sufficiently small step sizes) will converge to the suboptimal saddle point  $(0, 0)$ . On the other hand, polynomial-time convergence holds with random initialization (see [Du et al. \(2018\)](#)).

assumptions or algorithmic modifications to circumvent these negative results: For example, explicitly using the fact that the matrix sizes are larger than 1 – something which is not clear how to do with current analyses – or having a fine-grained dependency on the variance of the random initialization, as further discussed in Sec. 4. Alternatively, our results might be circumvented using other gradient-based algorithms (for example, by adding random noise to the gradient updates or using adaptive step sizes), or other initialization strategies. However, that would not explain why plain gradient descent with standard random initializations is often practically effective on these problems. Overall, we believe our results point to a potential obstacle in understanding the convergence of gradient-based methods for linear networks: At the very least, one would have to rule out one-dimensional layers, or consider algorithms other than plain gradient descent with standard initializations, in order to establish polynomial-time convergence guarantees for deep linear networks.

Finally, we note that our results provide a possibly interesting contrast to the recent work of Arora et al. (2018b), which suggests that increasing depth can sometimes accelerate the optimization process. Here we show that at least in some cases, the opposite occurs: Adding depth can quickly turn a trivial optimization problem into an intractable one for gradient descent.

## 2. Preliminaries

**Notation.** We use bold-faced letters to denote vectors. Given a vector  $\mathbf{w}$ ,  $w_j$  refers to its  $j$ -th coordinate.  $\|\cdot\|$ ,  $\|\cdot\|_1$  and  $\|\cdot\|_\infty$  refer to the Euclidean norm, the 1-norm and the infinity norm respectively. We let  $\prod_{i=1}^k w_i$  and  $\prod_i w_i$  be a shorthand for  $w_1 \cdot w_2 \cdots w_k$ . Also, we define a product over an empty set as being equal to 1. Since our main focus is to study the dependence on the network depth  $k$ , we use the standard notation  $\mathcal{O}(\cdot)$ ,  $\Omega(\cdot)$ ,  $\Theta(\cdot)$  to hide constants independent of  $k$ , and  $\tilde{\mathcal{O}}(\cdot)$ ,  $\tilde{\Omega}(\cdot)$ ,  $\tilde{\Theta}(\cdot)$  to hide constants and factors logarithmic in  $k$ .

**Gradient Descent.** We consider the standard gradient descent method for unconstrained optimization of functions  $F$  in Euclidean space, which given an initialization point  $\mathbf{w}(1)$ , performs repeated iterations of the form  $\mathbf{w}(t+1) := \mathbf{w}(t) - \eta \nabla F(\mathbf{w}(t))$  for  $t = 1, 2, \dots$  (where  $\nabla F(\cdot)$  is the gradient, and  $\eta > 0$  is a step size parameter). For objectives as in Eq. (2), we have  $\frac{\partial}{\partial w_j} F(\mathbf{w}) = f'(\prod_i w_i) \prod_{j \neq i} w_i$ , and gradient descent takes the form

$$\forall j, w_j(t+1) = w_j(t) - \eta f' \left( \prod_i w_i(t) \right) \prod_{j \neq i} w_i(t).$$

**Random Initialization.** One of the most common initialization methods for neural networks is *Xavier* initialization (Glorot and Bengio, 2010), which in the setting of Eq. (1) corresponds to choosing each entry of each  $d \times d$  matrix  $W_i$  independently from a zero-mean distribution with variance  $1/d$  (usually uniform or Gaussian). This ensures that the variance of the network outputs (with respect to the initialization) is constant irrespective of the network size. Motivated by residual networks, Hardt and Ma (2016) and Bartlett et al. (2018) consider initializing each  $W_i$  independently at  $I$ , possibly with some random perturbation. In this paper we denote such an initialization scheme as a *near-identity* initialization. Since we focus here on the case  $d = 1$  as in Eq. (2), Xavier initialization corresponds to choosing each  $w_i$  independently from a zero-mean, unit-variance distribution, and near-identity initialization corresponds to choosing each  $w_i$  close to 1.

### 3. Exponential Convergence Time for Gradient Descent

For our negative results, we impose the following mild conditions on the function  $f$  in Eq. (2):

**Assumption 1**  $f : \mathbb{R} \rightarrow \mathbb{R}$  is differentiable, Lipschitz continuous and strictly monotonically increasing on any interval  $[-\frac{1}{2}, z)$  where  $z > 0$ . Moreover,  $\inf_{p \in [-\frac{1}{2}, \infty)} f(p) - \inf_{p \in \mathbb{R}} f(p) > 0$ .

Here, we assume that  $f$  is fixed, and our goal is to study the convergence time of gradient descent on Eq. (2) as a function of the depth  $k$ . Some simple examples satisfying Assumption 1 in the context of machine learning include  $f(x) = (x + 1)^2$  and  $f(x) = \log(1 + \exp(x))$  (e.g., squared loss and logistic loss with respect to the input/output pair  $(1, -1)$ , respectively). We note that this non-symmetry with respect to positive/negative values is completely arbitrary, and one can prove similar results if their roles are reversed.

#### 3.1. Xavier Initialization

We begin with the case of Xavier initialization, where we initialize all coordinates of  $\mathbf{w}$  in Eq. (2) independently from a zero-mean, unit variance distribution. We will consider any distribution which satisfies the following:

**Assumption 2**  $w_1(1), \dots, w_k(1)$  are drawn i.i.d. from a zero-mean, unit variance distribution such that

1.  $\Pr(w_1(1) \in [-a, a]) \leq c_1 a$  for all  $a \geq 0$
2.  $\mathbb{E}[|w_1(1)|] \leq 1 - c_2$

where  $c_1, c_2 > 0$  are absolute constants independent of  $k$ .

The first part of the assumption is satisfied for any distribution with bounded density. As to the second part, the following lemma shows that it is satisfied for uniform and Gaussian distributions (with an explicit  $c_2$ ), and in fact for any non-trivial distribution (with a distribution-dependent  $c_2$ ):

**Lemma 1** *The following hold:*

- If  $w$  is drawn from a zero-mean, unit-variance Gaussian, then  $\mathbb{E}[|w|] < 0.8$ .
- If  $w$  is drawn from a zero-mean, unit-variance uniform distribution, then  $\mathbb{E}[|w|] < 0.9$ .
- If  $w$  is drawn from any zero-mean, unit variance distribution not supported on a single value, then  $\mathbb{E}[|w|] < 1$ .

**Proof** The first two parts follow from standard results on Gaussian and uniform distributions. As to the third part, by Jensen's inequality and the fact that  $\sqrt{\cdot}$  is a strictly concave function,  $\mathbb{E}[|w|] = \mathbb{E}[\sqrt{w^2}] < \sqrt{\mathbb{E}[w^2]} = 1$ . ■

With such an initialization, we now show that gradient descent is overwhelmingly likely to take at least exponential time to converge:

**Theorem 2** *The following holds for some positive constants  $c, c'$  independent of  $k$ : Under Assumptions 1 and 2, if gradient descent is ran with any step size  $\eta \leq \exp(ck)$ , then with probability at least  $1 - \exp(-\Omega(k))$  over the initialization, the number of iterations required to reach suboptimality less than  $c'$  is at least  $\exp(\Omega(k))$ .*

In the above,  $\Omega(\cdot)$  hides dependencies on the absolute constants in the theorem statement and the assumptions. The proof is presented in Sec. 5.

The intuition behind the theorem is quite simple: Under our assumptions, it is easy to show that the product of any  $\Omega(k)$  coordinates from  $w_1(1), \dots, w_k(1)$  is overwhelmingly likely to be exponentially small in  $k$ . Since the derivative of our objective w.r.t. any  $w_j$  has the form  $f'(\prod_i w_i) \prod_{i \neq j} w_i$ , it follows that the gradient is exponentially small in  $k$ . Moreover, we show that the gradient is exponentially small at any point within a bounded distance from the initialization (which is the main technical challenge of the proof, since the gradient is by no means Lipschitz). As a result, gradient descent will only make exponentially small steps. Assuming we start from a point bounded away from a global minimum, it follows that the number of required iterations must be exponentially large in  $k$ .

We note that the observation that Xavier initialization leads to highly skewed values in deep enough networks is not new (see Saxe et al. (2013); Pennington et al. (2017)), and has motivated alternative initializations such as orthogonal initialization<sup>2</sup>. Our contribution here is to rigorously analyze how this affects the optimization process for our setting.

### 3.2. Near-Identity Initialization

We now turn to consider initializations where each  $w_i$  is initialized close to 1. Here, it will be convenient to make deterministic rather than stochastic assumptions on the initialization point (which are satisfied with high probability for reasonable distributions):

**Assumption 3** *For some absolute constants  $c_1, c_2, c_3 > 0$  independent of  $k$ , gradient descent is initialized at a point  $\mathbf{w}(1)$  which satisfies  $\max_j |w_j(1) - 1| \leq k^{-c_1}$  and  $c_2 \leq \prod_i w_i(1) \leq c_3$ .*

To justify this assumption, note that if  $w_1(1), \dots, w_k(1)$  are chosen i.i.d. and not in the range of  $1 \pm k^{-c_1}$  for some  $c_1 > 0$ , then their product is likely to explode or vanish with  $k$ .

**Theorem 3** *The following holds for some positive constants  $c, c'$  independent of  $k$ : Under Assumptions 1 and 3, if gradient descent is ran with any positive step size  $\eta \leq c$ , then the number of iterations required to reach suboptimality less than  $c'$  is at least  $\exp(\Omega(k))$ .*

As before,  $\Omega(\cdot)$  hides dependencies on the absolute constants in the theorem statement, as well as those in the assumptions.

The formal proof appears in the appendix. To help explain its intuition, we provide in Figure 1 the actual evolution of  $w_j(t)$  for a typical run of gradient descent, when  $F(\mathbf{w}) = F(w_1, \dots, w_7) = \frac{1}{2}(\prod_{i=1}^7 w_i + 1)^2$  and we initialize all coordinates reasonably close to 1. Recall that for any  $w_j(t)$ ,

2. It is interesting to note that in our setting, orthogonal initialization amounts to choosing each  $w_i$  in  $\{-1, +1\}$ , which can easily cause non-convergence, e.g. for  $F(w_1, \dots, w_k) = (\prod_i w_i - y)^2$  when  $y \prod_i w_i(1) < 0$  and small enough step sizes.

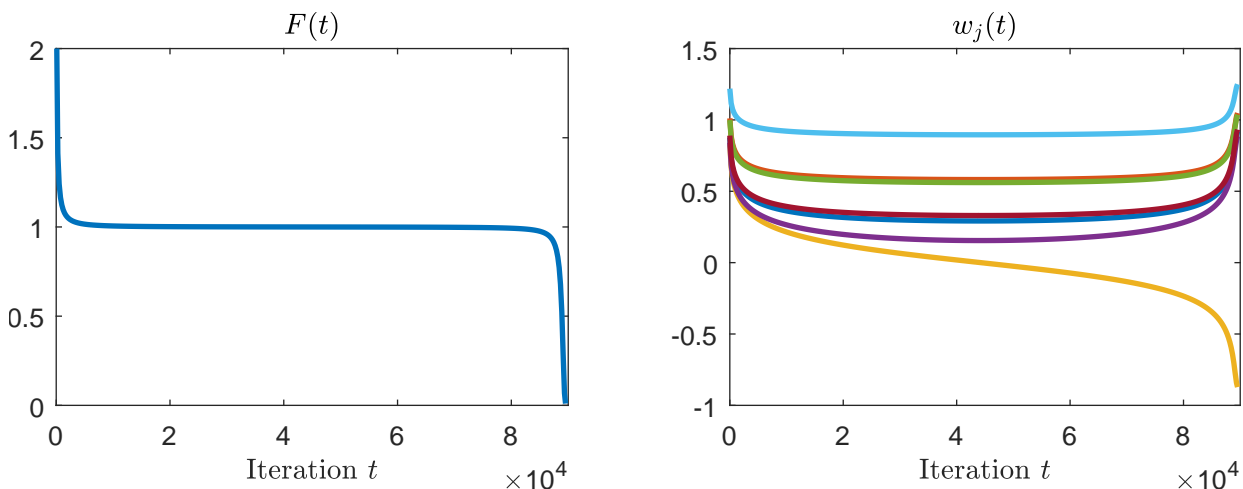


Figure 1: The left figure plots  $F(\mathbf{w}(t))$  as a function of iteration  $t$ , for  $F(\mathbf{w}) = (\prod_{i=1}^7 w_i + 1)^2$ . The right figure plots  $w_1(t), w_2(t), \dots, w_7(t)$  as a function of  $t$ . Best viewed in color.

the gradient descent updates take the form

$$\forall j, w_j(t+1) = w_j(t) - \eta \left( \prod_i w_i(t) + 1 \right) \prod_{i \neq j} w_i(t),$$

where  $\prod_i w_i(1) > 0$ . Thus, initially, all parameters  $w_j(t)$  decrease with  $t$ , as to be expected. However, as their value fall to around or below 1, their product decreases rapidly to  $\exp(-\Omega(k))$ . Since the gradient of each  $w_j(t)$  scales as  $\prod_{i \neq j} w_i(t)$ , the magnitude of the gradients becomes very small, and the algorithm makes only slow progress. Eventually, one of the parameters becomes negative, in which case all other parameters start increasing, and the algorithm converges. However, by a careful analysis, the length of the slow middle phase can be shown to be exponential in the depth / number of parameters  $k$ .

### 3.3. A Positive Result

Having established that the number of iterations is at least  $\exp(\Omega(k))$ , we now show that this is nearly tight. Specifically, we prove that gradient descent indeed converges in the settings studied so far, with a number of iterations scaling as  $\exp(\tilde{O}(k))$  (this can be interpreted as a constant for any constant  $k$ ). For simplicity, we prove this in the case where  $f(\prod_i w_i) = \frac{1}{2}(\prod_i w_i - y)^2$ , but the technique can be easily generalized to other convex  $f$  under mild conditions. We note that the case of  $y > 0$  and each  $w_i$  initialized to 1 is covered by the results in [Bartlett et al. \(2018\)](#). However, here we show a convergence result for other values of  $y$ , and even if  $w_i$  are not all initialized at 1.

We will use the following assumptions on our objective and parameters:

**Assumption 4** *The following hold for some absolute positive constants  $c_1, c_2, c_3, c_4$  independent of  $k$ :*

- $y = -c_1 < 0$
- *The initialization  $w_1(1), \dots, w_k(1)$  satisfies the following:*
  - $|w_i(1)| \leq c_2$  and  $\prod_i w_i(1) > y$
  - $\min_{j \neq j'} |w_j(1)| - |w_{j'}(1)| \geq k^{-c_4}$
  - $\max_{j, j'} \left| \prod_{i \notin \{j, j'\}} w_i(1) \right| \leq c_4$

The assumptions  $y < 0$  and  $\prod_i w_i(1) > y$  ensure that the objective satisfies the conditions of our negative results, for both Xavier and near-identity initializations (the other cases can be studied using similar techniques).

**Theorem 4** *Consider the objective  $F(\mathbf{w}) = \frac{1}{2} (\prod_i w_i - y)^2$ . Under Assumption 4, for any step size  $\eta = k^{-c}$  for some large enough constant  $c > 0$ , and for any  $\epsilon > 0$ , the number of gradient descent iterations  $t$  required for  $F(\mathbf{w}_t) \leq \epsilon$  is at most  $\exp(\tilde{O}(k)) \cdot \max\{1, \log(1/\epsilon)\}$ .*

The proof of the theorem appears in the appendix. Intuitively, it considers a similar situation as in Thm. 3, but with a careful analysis, shows that the algorithm still makes continuous progress (albeit exponentially small at certain phases).

#### 4. Multi-Dimensional Networks

So far, we showed that for one-dimensional linear neural networks, gradient descent can easily require exponentially many iterations (in the depth of the network) to converge. However, these results are specific to the case where the parameter matrix  $W_i$  of each layer is one-dimensional, and do not necessarily extend to higher dimensions. A possibly interesting exception is when  $F(W_1, \dots, W_k) = \left\| \prod_i W_i - Y \right\|_F^2$ , and both  $Y$  and the initialization  $W_1(1), \dots, W_k(1)$  are diagonal matrices. In that case, it is easy to show that the matrices produced by gradient descent remain diagonal, and the objective can be rewritten as a sum of independent one-dimensional problems for which our results would apply. However, this reasoning fails for non-diagonal initializations and target matrices  $Y$ .

In this section, we study experimentally whether our theoretical results for one-dimensional networks might also extend to multi-dimensional ones. In particular, we consider the multi-dimensional generalization of the objective function studied earlier:

$$F(W_1, \dots, W_k) = \frac{1}{2} \left\| \prod_{i=1}^k W_i - Y \right\|_F^2,$$

where  $W_1, \dots, W_k$  are  $d \times d$  square matrices (for  $d = 25$ ),  $Y = -I$  ( $I$  being the identity matrix), and  $\|\cdot\|_F$  is the Frobenius norm. We ran gradient descent on this objective using three initialization strategies:

1. *Xavier initialization:* Each entry of each matrix  $W_i$  was initialized independently from a zero-mean Gaussian with variance  $\frac{1}{d}$ .



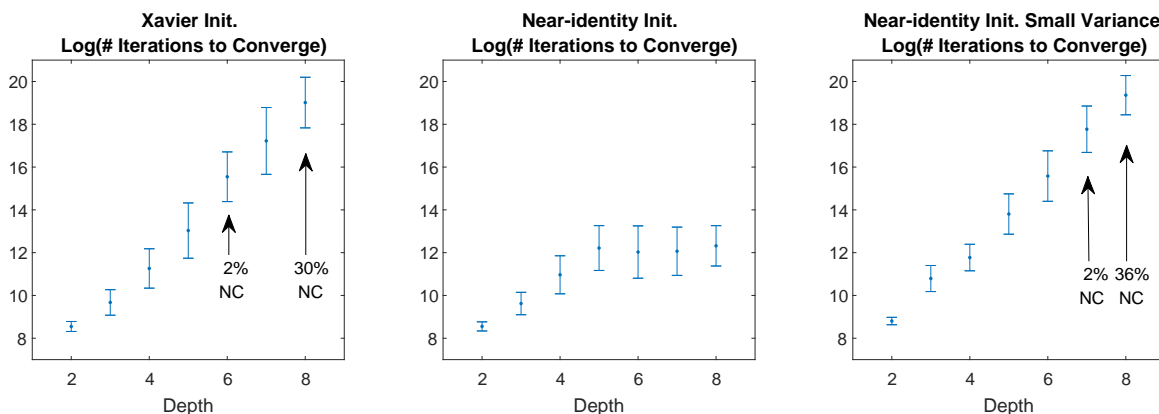


Figure 2: Mean and Standard Deviation of the log number of iterations required for convergence, over 50 trials, for each initialization strategy and depth parameter  $k \in \{2, 3, \dots, 8\}$ . ‘NC’ refers to the percentage of runs (for a given initialization and depth) which did not converge after  $10^9$  iterations, if any. Note that when some trials did not converge, the bars actually under-estimate the mean convergence time if all trials were ran till convergence (since they only represent runs which took a log number of iterations less than  $\log(10^9) = 20.72\dots$ ).

2. *Near-Identity initialization*: Each  $W_i$  was initialized as  $I + M$  where each entry of  $M$  was sampled independently from a zero-mean Gaussian with variance  $\frac{1}{dk}$ . Up to numerical constants, this is the largest variance which ensures that  $\mathbb{E}[(\prod_{i=1}^k W_i)(\prod_{i=1}^k W_i)^\top]$  remains bounded independent of  $d, k$ . To see this, note that had we used variance  $\frac{c}{dk}$  for some constant  $c$ , then  $\mathbb{E}[W_i W_i^\top] = (1 + \frac{c}{k}) I$  and thus  $\mathbb{E}[(\prod_{i=1}^k W_i)(\prod_{i=1}^k W_i)^\top] = (1 + \frac{c}{k})^k I \approx \exp(c)I$ .
3. *Near-Identity initialization with smaller variance*: Each  $W_i$  was initialized as above, except that the variance of each entry in the matrix  $M$  was  $\frac{1}{(dk)^2}$ .

For each random initialization strategy, and for depth parameter  $k \in \{2, 3, \dots, 8\}$ , we ran 50 trials of gradient descent, with a step size<sup>3</sup> of  $10^{-3}$ , until either one of the following two stopping conditions occurred:

- The objective value dropped below 0.1 (or equivalently,  $\|\prod_{i=1}^k W_i - Y\|_F \leq \sqrt{1/5}$ , a rather mild requirement).
- The number of iterations exceeded  $10^9$  iterations, in which case the algorithm was deemed to have failed to converge (note that from a practical viewpoint, one billion iterations is exceedingly large considering our problem size).

3. Our results did not seem to change significantly by taking other bounded step sizes.



In Figure 4, we plot the mean and standard deviation for the *logarithm* of the number of iterations required to make the objective value less than 0.1 (among the 50 trials which converged). We also point out the percentage of trials which did not converge, if any.

The figure strongly suggests that using both Xavier initialization and near-identity initialization with small variance, the required runtime scales exponentially with the depth (recall that the  $y$ -axis is in log scale). This indicates that the phenomenon of exponential scaling with depth is not just an artifact of one-dimensional networks, and can also occur in multi-dimensional networks, even with reasonable random initializations. On the flip side, when performing near-identity initialization with a large enough variance, we did not observe such an exponential scaling (as evidenced in the middle plot in the figure). Moreover, based on some additional experiments with other objective functions, it appears that although gradient descent can sometime require exponential time to converge, this phenomenon is not particularly common. A possible explanation to this is that in one dimension,  $\prod_i w_i$  had to change sign, and hence pass through zero (see Figure 1). This brought the iterates to a “flat” region with exponentially small gradients. In contrast, in multiple dimensions, to continuously change  $\prod_i W_i$  from a matrix to some other matrix, it is always possible to go “around” any particular point. Our experiments suggest that gradient descent indeed avoids problematic flat regions in many cases, but not always. Overall, it seems quite possible that for multi-dimensional networks, the exponential runtime dependence on the depth can be avoided under reasonable assumptions – however, some such assumptions would be necessary, and would need to exclude either objectives of the type we studied here, or some of the initializations. For example, such an analysis might need to explicitly separate between one-dimensional and multi-dimensional networks, or between near-identity initialization with variance  $1/dk$  and with variance  $1/(dk)^2$  (which are both polynomially large in  $d, k$ ), and how to do so with existing analyses is currently unclear.

## 5. Proof of Thm. 2

In this section, we provide the proof of Thm. 2. The proofs for all our other theorems (which are longer and hence do not fit in the main paper) are provided in the appendix.

The proof is based on the following two lemmas:

**Lemma 5** *Suppose  $w_1, \dots, w_k$  are drawn i.i.d. from a distribution such that  $\mathbb{E}[|w_1|] \leq a$  for some  $a > 0$ . Then*

$$\Pr \left( \max_j \left| \prod_{i \neq j} w_i \right| \geq ka^{(k-1)/2} \right) \leq a^{(k-1)/2}$$

**Proof** For any fixed  $j$ , by Markov’s inequality and the i.i.d. assumption,

$$\Pr \left( \left| \prod_{i \neq j} w_i \right| \geq ka^{(k-1)/2} \right) \leq \frac{\mathbb{E} \left[ \left| \prod_{i \neq j} w_i \right| \right]}{ka^{(k-1)/2}} = \frac{(\mathbb{E}[|w_1|])^{k-1}}{ka^{(k-1)/2}} \leq \frac{a^{k-1}}{ka^{(k-1)/2}} = \frac{1}{k} a^{(k-1)/2}.$$

Taking a union bound over all  $j = 1, 2, \dots, k$ , the result follows.  $\blacksquare$

**Lemma 6** *Let  $\alpha, \beta, \delta > 0$  be fixed. Let  $\mathbf{w} \in \mathbb{R}^k$  such that  $\max_j \left| \prod_{i \neq j} w_i \right| \leq \alpha$  and  $\min_i |w_i| \geq \delta$ . Then for any  $\mathbf{v}$  such that  $\|\mathbf{v} - \mathbf{w}\| \leq \frac{\delta}{\sqrt{k-1}} \log(\beta/\alpha)$ , it holds that  $\left| \prod_i v_i \right| \leq \beta \|\mathbf{v}\|_\infty$  as well as  $\|\nabla F(\mathbf{v})\| \leq \sup_{p: |p| \leq \beta \|\mathbf{v}\|_\infty} |f'(p)| \cdot \sqrt{k} \beta$ .*

**Proof** We claim that it is enough to prove the following:

$$\forall \mathbf{w}, \mathbf{v} \in \mathbb{R}^k \text{ s.t. } \max_j \left| \prod_{i \neq j} w_i \right| \leq \alpha, \quad \min_i |w_i| \geq \delta, \quad \max_j \left| \prod_{i \neq j} v_i \right| > \beta$$

it holds that  $\|\mathbf{v} - \mathbf{w}\| > \frac{\delta}{\sqrt{k-1}} \log(\beta/\alpha)$ . (3)

Indeed, this would imply that for any  $\mathbf{w}$  satisfying the conditions above, and any  $\mathbf{v}$  s.t.  $\|\mathbf{v} - \mathbf{w}\| \leq \frac{\delta}{\sqrt{k-1}} \log(\beta/\alpha)$ , we must have  $\max_j \left| \prod_{i \neq j} v_i \right| \leq \beta$ , and therefore  $\left| \prod_i v_i \right| \leq \beta \|\mathbf{v}\|_\infty$ , as well as  $\|\nabla F(\mathbf{v})\| = \sup_{p: |p| \leq \beta \|\mathbf{v}\|_\infty} |f'(p)| \cdot \left\| \left( \prod_{i \neq 1} v_i, \dots, \prod_{i \neq k} v_i \right) \right\| \leq \sup_{p: |p| \leq \beta \|\mathbf{v}\|_\infty} |f'(p)| \sqrt{k} \beta$  by definition of  $F$ , as required.

To prove Eq. (3), we first state and prove the following auxiliary result:

$$\forall \mathbf{w}, \mathbf{v} \in \mathbb{R}^{k-1} \text{ s.t. } \forall i v_i \geq w_i \geq 0, \quad \prod_i w_i \leq \alpha, \quad \min_i w_i \geq \delta, \quad \prod_i v_i > \beta$$

it holds that  $\|\mathbf{v} - \mathbf{w}\| > \frac{\delta}{\sqrt{k-1}} \log(\beta/\alpha)$ . (4)

This statement holds by the following calculation:

$$\begin{aligned} \|\mathbf{v} - \mathbf{w}\| &\geq \frac{1}{\sqrt{k-1}} \|\mathbf{v} - \mathbf{w}\|_1 = \frac{1}{\sqrt{k-1}} \cdot \sum_i (v_i - w_i) \\ &\stackrel{(*)}{\geq} \frac{1}{\sqrt{k-1}} \sum_i w_i (\log(v_i) - \log(w_i)) \geq \frac{\delta}{\sqrt{k-1}} \sum_i (\log(v_i) - \log(w_i)) \\ &= \frac{\delta}{\sqrt{k-1}} \log \left( \frac{\prod_i v_i}{\prod_i w_i} \right) > \frac{\delta}{\sqrt{k-1}} \log \left( \frac{\beta}{\alpha} \right), \end{aligned}$$

where  $(*)$  is due to the fact that  $\log(\cdot)$  is  $1/z$ -Lipschitz in  $[z, \infty)$ , and the assumption that  $v_i \geq w_i \geq 0$ .

It remains to explain how Eq. (4) implies Eq. (3). Indeed, let  $\mathbf{w}, \mathbf{v}$  be any two vectors in  $\mathbb{R}^k$  which satisfy the conditions of Eq. (3). Now, suppose we transform them into vectors  $\mathbf{w}', \mathbf{v}' \in \mathbb{R}^{k-1}$  by the following procedure:

- Change the sign of every  $w_i$  and  $v_i$  to be positive
- For any  $i$  such that  $v_i < w_i$ , change  $v_i$  to equal  $w_i$ .
- Drop a coordinate  $j$  which maximizes  $\left| \prod_{i \neq j} v_i \right|$ .

It is easy to verify that the resulting vectors  $\mathbf{w}', \mathbf{v}'$  satisfy the conditions of Eq. (4), and  $\|\mathbf{v}' - \mathbf{w}'\| \leq \|\mathbf{v} - \mathbf{w}\|$ . Therefore, by Eq. (4),  $\|\mathbf{v} - \mathbf{w}\| \geq \|\mathbf{v}' - \mathbf{w}'\| \geq \frac{\delta}{\sqrt{k-1}} \log(\beta/\alpha)$  as required.  $\blacksquare$

With these two lemmas in hand, we turn to prove the theorem. By Lemma 5 and Assumption 2, we have

$$\Pr \left( \max_j \left| \prod_{i \neq j} w_i(1) \right| \geq \exp(-2Ck) \right) \leq \exp(-C'k).$$

for some fixed constants  $C, C' > 0$  and any large enough  $k$ . Moreover, again by Assumption 2, it holds for any  $i$  that  $\Pr(|w_i(1)| \leq \exp(-Ck)) \leq \mathcal{O}(\exp(-Ck))$ , so by a union bound,

$$\Pr(\min_i |w_i| < \exp(-Ck)) \leq \mathcal{O}(k \exp(-Ck)).$$

Finally, by Assumption 2, Markov's inequality and a union bound,

$$\Pr(\|\mathbf{w}(1)\|_\infty \geq \exp(Ck)) \leq k \exp(-Ck)$$

Combining the last three displayed equations with a union bound, and applying Lemma 6 (with  $\alpha = \exp(-2Ck)$ ,  $\beta = 2\alpha$ , and  $\delta = \exp(-Ck)$ ), we get the following: With probability at least  $1 - \exp(-C'k) - \mathcal{O}(k \exp(-Ck)) - k \exp(-Ck) = 1 - \exp(-\Omega(k))$  over the choice of  $\mathbf{w}(1)$ ,

- $\|\mathbf{w}(1)\|_\infty \leq \exp(Ck)$ .
- For any  $\mathbf{v}$  at a distance at most  $\exp(-Ck) \frac{\log(2)}{\sqrt{k-1}}$  from  $\mathbf{w}(1)$ , we have

$$\|\mathbf{v}\|_\infty \leq \|\mathbf{w}(1)\|_\infty + \exp(-Ck) \frac{\log(2)}{\sqrt{k-1}} \leq \mathcal{O}(\exp(Ck)),$$

$$\left| \prod_i v_i \right| \leq \beta \|\mathbf{v}\|_\infty = 2 \exp(-2Ck) \cdot \mathcal{O}(\exp(Ck)) = \mathcal{O}(\exp(-Ck))$$

and

$$\begin{aligned} \|\nabla F(\mathbf{v})\| &\leq \sup_{p: |p| \leq \beta \|\mathbf{v}\|_\infty} |f'(p)| \cdot \sqrt{k} \beta \leq \sup_{p: |p| \leq \mathcal{O}(\exp(-Ck))} |f'(p)| \cdot 2\sqrt{k} \exp(-2Ck) \\ &= \mathcal{O}\left(\sqrt{k} \exp(-2Ck)\right). \end{aligned}$$

This has two implications:

1. Since the gradient descent updates are of the form  $\mathbf{w}(t+1) = \mathbf{w}(t) - \eta \nabla F(\mathbf{w}(t))$ , and we can assume  $\eta \leq \exp(Ck/2)$  by the theorem's conditions, the number of iterations required to get to a distance larger than  $\exp(-Ck) \frac{\log(2)}{\sqrt{k-1}}$  from  $\mathbf{w}(1)$  is at least

$$\frac{\exp(-Ck) \frac{\log(2)}{\sqrt{k-1}}}{\exp(Ck/2) \cdot \mathcal{O}(\sqrt{k} \exp(-2Ck))} = \Omega\left(\frac{\exp(Ck/2)}{k}\right),$$

which is at least  $\exp(\Omega(k))$  iterations.

2. As long as we are at a distance smaller than the above,  $|\prod_i v_i| \leq \mathcal{O}(\exp(-Ck)) \leq \exp(-\Omega(k))$ . In particular,  $\prod_i v_i \geq -1/2$  for large enough  $k$ , so by Assumption 1 and definition of  $F$ , we have that  $F(\mathbf{v}) - \inf_{\mathbf{v}} F(\mathbf{v})$  is lower bounded by a constant independent of  $k$ .

Overall, we get that with probability at least  $1 - \exp(-\Omega(k))$ , we initialize at some region in which all points are at least  $\Omega(1)$  suboptimal, and at least  $\exp(\Omega(k))$  iterations are required to escape it. This immediately implies our theorem.

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## Appendix A. Additional Proofs

### A.1. Proof of Thm. 3

We begin with the following auxiliary lemma, and then turn to analyze the dynamics of gradient descent in our setting.

**Lemma 7** *For any positive scalars  $\alpha, w_1, \dots, w_k$  such that  $\min_i w_i > \alpha$ ,*

$$\prod_i (w_i - \alpha) \leq \left( \left( \prod_i w_i \right)^{1/k} - \alpha \right)^k.$$

**Proof** Taking the  $k$ -th root and switching sides, the inequality in the lemma is equivalent to proving

$$\left( \prod_i (w_i - \alpha) \right)^{1/k} + \alpha \leq \left( \prod_i w_i \right)^{1/k}.$$

Letting  $a_i = w_i - \alpha$ , and  $b_i = \alpha$  for all  $i$ , the above is equivalent to proving that

$$\left( \prod_i a_i \right)^{1/k} + \left( \prod_i b_i \right)^{1/k} \leq \left( \prod_i (a_i + b_i) \right)^{1/k},$$

namely that the sum of the geometric means of two positive sequences  $(a_i)$  and  $(b_i)$  is at most the geometric mean of their sum  $(a_i + b_i)$ . This follows from the superadditivity of the geometric mean (see [Steele \(2004, Exercise 2.11\)](#)) ■

**Lemma 8** *If  $\min_i w_i(t) \geq C$  and  $\prod_i w_i(t) \leq C'$  for some positive constants  $C, C'$ , then for any  $j, j'$ ,*

$$|w_j(t+1)^2 - w_{j'}(t+1)^2| \leq |w_j(t)^2 - w_{j'}(t)^2| + C'' \eta^2 \left( \prod_i w_i(t) \right)^2,$$

where  $C''$  is some constant dependent only on  $C, C'$  and the function  $f$ .

**Proof** By definition,

$$\begin{aligned} & w_j(t+1)^2 - w_{j'}(t+1)^2 \\ &= \left( w_j(t) - \eta f' \left( \prod_i w_i(t) \right) \prod_{i \neq j} w_i(t) \right)^2 - \left( w_{j'}(t) - \eta f' \left( \prod_i w_i(t) \right) \prod_{i \neq j'} w_i(t) \right)^2 \\ &= w_j(t)^2 - w_{j'}(t)^2 + \eta^2 f' \left( \prod_i w_i(t) \right)^2 \left( \left( \prod_{i \neq j} w_i(t) \right)^2 - \left( \prod_{i \neq j'} w_i(t) \right)^2 \right) \\ &= w_j(t)^2 - w_{j'}(t)^2 + \eta^2 \left( \prod_i w_i(t) \right)^2 \cdot f' \left( \prod_i w_i(t) \right) \left( \frac{1}{w_j(t)^2} - \frac{1}{w_{j'}(t)^2} \right). \end{aligned}$$

By assumption,  $0 \leq \prod_i w_i(t) \leq C'$  and  $\max_j \frac{1}{w_j(t)^2} \leq \frac{1}{C^2}$ . Therefore, by our assumptions on  $f$ , the displayed equation above implies that

$$|w_j(t+1)^2 - w_{j'}(t+1)^2| \leq |w_j(t)^2 - w_{j'}(t)^2| + C''\eta^2 \left( \prod_i w_i(t) \right)^2$$

for some constant  $C'' > 0$  dependent on  $C, C'$  and  $f$  as required.  $\blacksquare$

**Lemma 9** *Suppose that at some iteration  $t$ , for some constant  $C$  independent of  $k$ , it holds that  $\max_i w_i(t) \leq C$  and  $\prod_i w_i(t) \leq \beta$  for some  $\beta \in (0, C)$ . Then after at most  $\tau \leq 1 + \mathcal{O}(1) \cdot \frac{\beta^{1/k-1}}{\eta^k}$  iterations, if  $\min_j w_j(r) \geq 1/2$  for all  $r = t, t+1, \dots, t+\tau$ , then*

- Each  $w_i(r)$  as well as  $\prod_i w_i(r)$  monotonically decrease in  $r = t, t+1, \dots, t+\tau$
- For all  $r = t, t+1, \dots, t+\tau-1$ ,  $\max_j |w_j(r+1) - w_j(r)| \leq \mathcal{O}(1) \cdot \eta\beta$
- $\prod_i w_i(t+\tau) \leq \beta \cdot \exp(-1)$ .

In the above,  $\mathcal{O}(1)$  hides constants dependent only on  $C$  and the function  $f$ .

**Proof** If  $\prod_i w_i(t) \leq \beta \cdot \exp(-1)$ , we can pick  $\tau = 0$ , and the lemma trivially holds. Otherwise, let  $\tau$  be the smallest (positive) index such that  $\prod_i w_i(t) \leq \beta \cdot \exp(-1)$  (if no such index exists, take  $\tau = \infty$ , although the arguments below imply that  $\tau$  must be finite). Since we assume  $w_i(r)$  for all  $i$  are positive, and  $f$  is monotonically increasing,

$$w_j(r+1) = w_j(r) - \eta f' \left( \prod_i w_i(r) \right) \prod_{i \neq j} w_i(r) \leq w_j(r),$$

so  $w_j(r)$  monotonically decreases in  $r$ . Moreover, these are all positive numbers by assumption, so  $\prod_i w_i(r)$  monotonically decreases in  $r$  as well. This shows the first part of the lemma.

As to the second part, the displayed equation above, the fact that  $w_j(r)$  and  $\prod_i w_i(r)$  decrease in  $r$ , and our assumptions on  $f$  imply that for any  $r < t+\tau$ ,

$$\begin{aligned} w_j(r+1) &= w_j(r) - \eta f' \left( \prod_i w_i(r) \right) \prod_{i \neq j} w_i(r) = w_j(r) - \frac{\eta}{w_j(r)} f' \left( \prod_i w_i(r) \right) \prod_i w_i(r) \\ &= w_j(r) - \Theta(1) \cdot \eta\beta. \end{aligned}$$

where  $\Theta(1)$  hides constants dependent only on  $f$  and  $C$ . As to the third part of the lemma, fix some  $s < \tau$ , and repeatedly apply the displayed equation above for  $r = t, t+1, \dots, t+s$ , to get that that  $w_j(t+s) \leq w_j(t) - \Theta(1) \cdot \eta\beta s$  (which is still  $\geq 1/2$  by the lemma assumptions). In that case,

$$\begin{aligned} \prod_i w_i(t+s) &\leq \prod_i (w_i(t) - \Theta(1) \cdot \eta\beta s) \stackrel{(*)}{\leq} \left( \beta^{1/k} - \Theta(1) \cdot \eta\beta s \right)^k = \beta \left( 1 - \Theta(1) \cdot \eta\beta^{1-1/k} s \right)^k \\ &\leq \beta \exp \left( -\Theta(1) \cdot \eta\beta^{1-1/k} s k \right) \end{aligned}$$



where (\*) follows from Lemma 7 and the fact that  $\prod_i w_i(t) \leq \beta$ . The right hand side in turn is at most  $\beta \cdot \exp(-1)$  for any  $s \geq C' \beta^{1/k-1}/\eta k$  for some constant  $C'$ . In particular, if  $\tau > 1 + C' \beta^{1/k-1}/\eta k$ , then by choosing  $s$  s.t.  $\tau > s \geq C' \beta^{1/k-1}/\eta k$ , we get that  $\prod_i w_i(t+s) \leq \beta \cdot \exp(-1)$  even though  $s < \tau$ , which contradicts the definition of  $\tau$ . Hence  $\tau \leq 1 + C' \beta^{1/k-1}/\eta k$  as stated in the lemma.  $\blacksquare$

Combining Lemma 8 and Lemma 9, we have the following:

**Lemma 10** *For any constants  $C > 0$  and index  $T$ , if  $\prod_i w_i(1) \leq C$  and  $w_i(t) \geq \frac{1}{2}$  for all  $i = 1, \dots, k$  and  $t = 1, 2, \dots, T$ , then for all such  $t$ ,*

- Each  $w_i(t)$  as well as  $\prod_i w_i(t)$  monotonically decrease in  $t$ .
- $\max_j |w_j(t+1) - w_j(t)| \leq \mathcal{O}(1) \cdot \eta$
- $\max_{j,j'} |w_j(t) - w_{j'}(t)| \leq k^{-\Omega(1)} + \mathcal{O}(1) \cdot (\eta^2 + \frac{\eta}{k})$ .

In the above,  $\mathcal{O}(\cdot)$  hides constants dependent only on  $C$  and the constants in Assumptions 1 and 3.

**Proof** The first two parts of the lemma follow from Lemma 9 and the fact that by Assumption 3,  $w_i(1) \leq 1 + k^{-\Omega(1)} \leq \mathcal{O}(1)$ . As to the last part, define  $t_0 \leq t_1 \leq \dots \leq t_s$  (where  $t_0 = 1$ ) as the first indices  $\leq T$  such that for all  $r = 0, \dots, s$ ,  $\prod_i w_i(t_r) \leq (\prod_i w_i(1)) \exp(-r)$  (where  $s$  is taken to be as large as possible). By Lemma 9, we have the following:

- For all  $r = 0, \dots, s-1$ ,  $|t_{r+1} - t_r| \leq 1 + \mathcal{O}(1) \cdot \frac{\exp(-r)^{1/k-1}}{\eta k}$ .
- $|T - t_s| \leq 1 + \mathcal{O}(1) \cdot \frac{\exp(-s)^{1/k-1}}{\eta k}$ .
- For all  $r = 0, \dots, s-1$  and any  $t_r \leq t \leq t_{r+1}$ , we have  $\prod_i w_i(t) \leq \mathcal{O}(1) \cdot \exp(-r)$ .

Combining this with Lemma 8, it follows that for any  $j, j'$ , and any  $r = 0, \dots, s-1$ ,

$$\begin{aligned} |w_j(t_{r+1})^2 - w_{j'}(t_{r+1})^2| &\leq |w_j(t_r)^2 - w_{j'}(t_r)^2| + \mathcal{O}(1) \cdot \eta^2 \exp(-2r) \cdot \left(1 + \frac{\exp(-r)^{1/k-1}}{\eta k}\right) \\ &\leq |w_j(t_r)^2 - w_{j'}(t_r)^2| + \mathcal{O}(1) \cdot \left(\eta^2 \exp(-2r) + \frac{\eta \exp(-r)}{k}\right), \end{aligned}$$

as well as

$$|w_j(T)^2 - w_{j'}(T)^2| \leq |w_j(t_s)^2 - w_{j'}(t_s)^2| + \mathcal{O}(1) \cdot \left(\eta^2 \exp(-2s) + \frac{\eta \exp(-s)}{k}\right).$$

Repeatedly applying the last two displayed equations, and using Assumption 3, we get that

$$\begin{aligned} |w_j(T)^2 - w_{j'}(T)^2| &\leq |w_j(1)^2 - w_{j'}(1)^2| + \mathcal{O}(1) \cdot \left(\eta^2 \sum_{r=0}^s \exp(-2r) + \frac{\eta}{k} \sum_{r=0}^s \exp(-r)\right) \\ &\leq k^{-\Omega(1)} + \mathcal{O}(1) \cdot \left(\eta^2 + \frac{\eta}{k}\right). \end{aligned}$$

Since  $|w_j(T)^2 - w_{j'}(T)^2| = |w_j(T) + w_{j'}(T)| \cdot |w_j(T) - w_{j'}(T)| \geq |w_j(T) - w_{j'}(T)|$  (as we have  $\min_i w_i(T) \geq 1/2$  by assumption), we get that  $|w_j(T) - w_{j'}(T)| \leq k^{-\Omega(1)} + \mathcal{O}(1) \cdot (\eta^2 + \frac{\eta}{k})$  as required.  $\blacksquare$

With Lemma 10 in hand, we can now prove the theorem. Let  $T$  be the largest index such that  $\min_i w_i(t) \geq 1/2$  for all  $t = 1, 2, \dots, T$  (and  $\infty$  if this holds for all  $t$ ). It follows that  $\prod_i w_i(t) \geq 0$ , and therefore, by Assumption 1,  $F(\mathbf{w}(t)) - \inf_{\mathbf{w}} F(\mathbf{w})$  is at least a constant independent of  $k$  for all  $t = 1, 2, \dots, T$ . Thus, to prove the theorem, it is enough to show that if  $T < \infty$ , then  $T \geq \exp(\Omega(k))$ .

By Assumption 3 and Lemma 10, we have that  $w_1(1) \geq 1 - k^{-\Omega(1)}$ ,  $|w_1(t+1) - w_1(t)| \leq \mathcal{O}(1) \cdot \eta$ , and  $\max_j |w_j(t) - w_1(t)| \leq k^{-\Omega(1)} + \mathcal{O}(1) \cdot (\eta^2 + \frac{\eta}{k})$ . On the other hand, if  $T < \infty$ , then  $\min_i w_i(T+1) < 1/2$ . Therefore, if  $k$  is large enough and  $\eta$  is small enough, there exists some iteration  $t \leq T$  such that  $w_j(t) \in [2/3, 3/4]$  for all  $j$ . This means that  $\prod_i w_i(t) \leq (3/4)^k = \exp(-\Omega(k))$ . Thus, by Lemma 9 (with  $\beta = \exp(-\Omega(k))$ ), from iteration  $t$  till iteration  $T$ , each  $w_j$  decreases by at most  $\mathcal{O}(1) \cdot \eta\beta \leq \exp(-\Omega(k))$  at each iteration. By assumption, at iteration  $T+1$ , there is some  $w_j(T+1) < 1/2$ , so we must have  $T-t \geq (2/3 - 1/2) / \exp(-\Omega(k)) = \exp(\Omega(k))$  as required.

## A.2. Proof of Thm. 4

To prove the theorem, we first state and prove the following key lemma:

**Lemma 11** *For any initialization  $\mathbf{w}(1)$  and any  $(\sigma_1, \dots, \sigma_k) \in \{-1, +1\}^k$ , let  $\mathbf{v}(1), \mathbf{v}(2), \dots$  denote the iterates produced by gradient descent starting from  $\mathbf{v}(1) := (\sigma_1 w_1(1), \dots, \sigma_k w_k(1))$ , w.r.t. the function*

$$F_\sigma(\mathbf{v}) := \frac{1}{2} \left( \prod_i v_i - \sigma y \right)^2,$$

where  $\sigma := \prod_i \sigma_i$ . Then for any  $t \geq 1$ ,

$$\mathbf{v}(t) = (\sigma_1 w_1(t), \dots, \sigma_k w_k(t)) \quad \text{and} \quad F(\mathbf{w}(t)) = F_\sigma(\mathbf{v}(t)).$$

**Proof** We prove the lemma by induction. The base case ( $t = 1$ ) is immediate from the definitions and the fact that

$$F_\sigma(\mathbf{v}(1)) = \frac{1}{2} \left( \prod_i \sigma_i w_i(1) - \sigma y \right)^2 = \frac{1}{2} \left( \sigma \prod_i w_i(1) - \sigma y \right)^2 = F(\mathbf{w}(1)).$$

Assuming that the induction hypothesis holds for  $t$ , and recalling that  $\sigma = \prod_i \sigma_i$ , we have for any  $j \in \{1, \dots, k\}$  that

$$\begin{aligned} v_j(t+1) &= v_j(t) - \left( \prod_i v_i(t) - \sigma y \right) \prod_{i \neq j} v_i(t) = \sigma_j w_j(t) - \sigma \left( \prod_i w_i(t) - y \right) \prod_{i \neq j} \sigma_i w_i(t) \\ &= \sigma_j \left( w_j(t) - \left( \prod_i w_i(t) - y \right) \prod_{i \neq j} w_i(t) \right) = \sigma_j w_j(t+1). \end{aligned}$$

As a result,

$$F_\sigma(\mathbf{v}(t+1)) = \frac{1}{2} \left( \prod_i v_i(t+1) - \sigma y \right)^2 = \frac{1}{2} \left( \sigma \prod_i w_i(t+1) - \sigma y \right)^2 = F(\mathbf{w}(t+1)).$$

This establishes the inductive step for  $t+1$ , hence proving the lemma.  $\blacksquare$

The lemma implies that for studying the dynamics of gradient descent starting from any initial point  $(w_1(1), \dots, w_k(1))$ , we can arbitrarily change the signs of its coordinates, as long as the sign of  $y$  is changed accordingly. In particular, we will assume without loss of generality that all  $w_1(1), \dots, w_k(1)$  are positive (again, as long as the sign of  $y$  is fixed accordingly). The proof then proceeds as follows:

- The simplest case is when after the sign transformations,  $y > 0$ . By our assumptions, this implies that both  $y$  and  $\prod_i w_i(1)$  switched from being negative (and satisfying  $\prod_i w_i(1) > y$ ) to positive, hence we now have  $y > \prod_i w_i(1) > 0$ . In that case, Lemma 13 below implies that  $\exp(\tilde{\mathcal{O}}(k)) \log(1/\epsilon)$  iterations suffice.
- The case  $y < 0$  (which by our assumptions, implies  $y < 0 < \prod_i w_i(1)$ ) is more involved: First, we show that after  $t = \exp(\tilde{\mathcal{O}}(k))$  iterations, one (and only one) of the coordinates of  $\mathbf{w}(t)$  becomes non-positive (Lemma 15). Then, we show that after at most one additional iteration, that non-positive coordinate becomes negative and bounded away from 0 (Lemma 16), the other coordinates remaining strictly positive. By Lemma 11, we can then argue that at that time point, the dynamics become the same as the scenario where  $y > 0$ , and all coordinates of the iterate are strictly positive. Again applying Lemma 13, we get that  $\exp(\tilde{\mathcal{O}}(k)) \log(1/\epsilon)$  additional iterations suffice for convergence.

#### A.2.1. THE CASE $y > \prod_i w_i(1) > 0$

We will need the following auxiliary lemma:

**Lemma 12** For any  $a > 0$ ,  $b \geq 0$ ,  $\log(a+b) \leq \log(a) + \frac{b}{a}$ .

**Proof** Since  $\log(1+z) \leq z$  for all  $z \geq 0$ , we have  $\log(a+b) = \log(a(1+b/a)) = \log(a) + \log(1+b/a) \leq \log(a) + b/a$ .  $\blacksquare$

**Lemma 13** Fix some  $\gamma \geq \delta > 0$ . Suppose that  $y > 0$ , and gradient descent on  $F$  is initialized at some  $\mathbf{w}(1)$  such that  $\prod_i w_i(1) \in [0, y)$ ,  $w_{j^*}(1) \geq \delta$  for some  $j^* \in \arg \min_i w_i(1)$ , and  $w_j(1) \geq \gamma$  for all  $j \neq j^*$ . Assuming step size  $\eta \leq \delta^2/2ky^2$ , we have that  $F(\mathbf{w}(t)) \leq \epsilon$  for any  $t \geq \frac{\log(y^2/2\epsilon)}{k\delta^2\gamma^2(k-2)\eta}$ .

**Proof** Let

$$\mathcal{W} := \left\{ \mathbf{w} \in \mathbb{R}^k : \prod_i w_i(1) \in [0, y), \min_i w_i(1) \geq \delta, \forall j \neq j^* w_j(1) \geq \gamma \right\}$$

denote the set of points in  $\mathbb{R}^k$  which satisfy the initialization conditions of the lemma.

First, we show that if the step size  $\eta$  is small enough, then gradient descent will remain in  $\mathcal{W}$  forever. For that, it is enough to show that for any  $\mathbf{w} \in \mathcal{W}$ , the update  $\mathbf{w}' := \mathbf{w} - \eta \nabla F(\mathbf{w})$  produced by gradient descent is in  $\mathcal{W}$  as well. By definition of  $\mathcal{W}$ , it is easily verified that  $w'_i \geq w_i > 0$  for all  $i$ , so the only non-trivial condition to verify is that  $\prod_j w'_j < y$ . To show this, we note that by Lemma 12,

$$\begin{aligned} \log \left( \prod_j w'_j \right) &= \sum_j \log(w'_j) = \sum_j \log \left( w_j + \eta(y - \prod_i w_i) \prod_{i \neq j} w_i \right) \\ &\leq \sum_j \log(w_j) + \eta(y - \prod_i w_i) \sum_j \frac{\prod_{i \neq j} w_i}{w_j} \\ &= \log \left( \prod_j w_j \right) + \eta(y - \prod_i w_i) \left( \prod_i w_i \right) \sum_j \frac{1}{w_j^2} \\ &< \log \left( \prod_j w_j \right) + \eta(y - \prod_i w_i) y \sum_j \frac{1}{\delta^2} \\ &= \log \left( \prod_j w_j \right) + \eta \frac{yk}{\delta^2} \left( y - \prod_i w_i \right). \end{aligned}$$

Thus, to ensure that  $\prod_j w'_j < y$  (or equivalently,  $\log(\prod_j w'_j) < \log(y)$ ), it is enough to ensure that

$$\log \left( \prod_j w_j \right) + \eta \frac{yk}{\delta^2} \left( y - \prod_i w_i \right) \leq \log(y).$$

Rearranging the above, we require that

$$\eta \frac{yk}{\delta^2} \leq \frac{\log(y) - \log(\prod_j w_j)}{y - \prod_j w_j}.$$

By the mean value theorem and the fact that  $\prod_j w_j < y$ , the right hand side can be lower bounded by  $\min_{z \in (0, y]} \log'(z) = 1/y$ , so it is enough to require

$$\eta \frac{yk}{\delta^2} \leq \frac{1}{y} \Rightarrow \eta \leq \frac{\delta^2}{ky^2},$$

which indeed holds by assumption.

Having established that gradient descent will remain in  $\mathcal{W}$  forever, we now establish that the objective  $F$  has a  $\frac{2ky^2}{\delta^2}$ -Lipschitz gradient on  $\mathcal{W}$ : Indeed, the Hessian of  $F$  at any  $\mathbf{w} \in \mathcal{W}$  can be easily verified to equal

$$(\nabla^2 F(\mathbf{w}))_{r,s} = \begin{cases} (\prod_i w_i - y) \frac{\prod_i w_i}{w_r w_s} + \frac{\prod_i w_i^2}{w_r w_s} & r \neq s \\ \frac{\prod_i w_i^2}{w_r^2} & r = s \end{cases}.$$

Since  $\mathbf{w} \in \mathcal{W}$ , it follows that magnitude of each entry in the  $k \times k$  Hessian is at most  $y \cdot \frac{y}{\delta^2} + \frac{y^2}{\delta^2} = 2y^2/\delta^2$ , and therefore its spectral norm (which is at most the Frobenius norm) can be upper bounded by  $2ky^2/\delta^2$ .

The final ingredient we need is that  $F$  satisfies

$$\|\nabla F(\mathbf{w})\|^2 \geq 2k\delta^2\gamma^{2(k-2)}F(\mathbf{w})$$

for any  $\mathbf{w} \in \mathcal{W}$  (this type of inequality is known as the Polyak-Łojasiewicz condition, which ensures linear convergence rates for gradient descent on possibly non-convex functions – see [Polyak \(1963\)](#); [Karimi et al. \(2016\)](#)). This follows from  $\|\nabla F(\mathbf{w})\|^2$ , by definition, being equal to

$$\left(\prod_i w_i - y\right)^2 \sum_j \left(\prod_{i \neq j} w_i\right)^2 = 2F(\mathbf{w}) \sum_j \left(\prod_{i \neq j} w_i\right)^2 \geq 2F(\mathbf{w})k \left(\delta\gamma^{k-2}\right)^2.$$

Collecting these ingredients, we can now perform a standard analysis using the Polyak-Łojasiewicz condition: If we do a gradient step to get from  $\mathbf{w} \in \mathcal{W}$  to  $\mathbf{w}' \in \mathcal{W}$  (i.e.  $\mathbf{w}' := \mathbf{w} - \eta\nabla F(\mathbf{w})$ ), and assuming  $\eta \leq \delta^2/2ky^2$ , then

$$\begin{aligned} F(\mathbf{w}') &\leq F(\mathbf{w}) + \nabla F(\mathbf{w})^\top (\mathbf{w}' - \mathbf{w}) + \frac{ky^2}{\delta^2} \|\mathbf{w}' - \mathbf{w}\|^2 \\ &= F(\mathbf{w}) - \eta \|\nabla F(\mathbf{w})\|^2 + \frac{ky^2}{\delta^2} \|\eta \cdot \nabla F(\mathbf{w})\|^2 \\ &= F(\mathbf{w}) - \eta \left(1 - \frac{ky^2}{\delta^2}\eta\right) \|\nabla F(\mathbf{w})\|^2 \\ &\leq F(\mathbf{w}) - \eta \cdot \frac{1}{2} \cdot 2k\delta^2\gamma^{2(k-2)}F(\mathbf{w}) \\ &= \left(1 - k\delta^2\gamma^{2(k-2)}\eta\right) F(\mathbf{w}) \leq \exp(-k\delta^2\gamma^{2(k-2)}\eta)F(\mathbf{w}). \end{aligned}$$

Applying this inequality  $t$  times, we get that

$$F(\mathbf{w}(t)) \leq \exp\left(-k\delta^2\gamma^{2(k-2)}\eta t\right) F(\mathbf{w}(1)) \leq \frac{y^2}{2} \exp\left(-k\delta^2\gamma^{2(k-2)}\eta t\right).$$

Equating the bound above to the target accuracy  $\epsilon$  and solving for  $t$ , the result follows.  $\blacksquare$

#### A.2.2. THE CASE $y < 0 < \prod_i w_i(1)$

We first state the following auxiliary lemma, which establishes that the gaps between coordinates are monotonically increasing under suitable assumptions.

**Lemma 14** *Fix some coordinate indices  $j, j'$  and iteration  $t$ , and suppose that  $w_j(t) \leq w_{j'}(t)$ ,  $\min_i w_i(t) \geq 0$ , and  $y < 0$ . Then  $w_{j'}(t) - w_j(t) \leq w_{j'}(t+1) - w_j(t+1)$ .*

**Proof** Dropping the  $(t)$  index to simplify notation, we have by definition that  $w_{j'}(t+1) - w_j(t+1)$  equals

$$\begin{aligned} & \left( w_{j'} - \eta \left( \prod_i w_i - y \right) \prod_{i \neq j'} w_i \right) - \left( w_j - \eta \left( \prod_i w_i - y \right) \prod_{i \neq j} w_i \right) \\ &= (w_{j'} - w_j) - \eta \left( \prod_i w_i - y \right) \left( \prod_{i \neq j'} w_i - \prod_{i \neq j} w_i \right) \\ &= (w_{j'} - w_j) - \eta \left( \prod_i w_i - y \right) (w_j - w_{j'}) \prod_{i \notin \{j, j'\}} w_i = (w_{j'} - w_j) \left( 1 + \eta \left( \prod_i w_i - y \right) \prod_{i \notin \{j, j'\}} w_i \right). \end{aligned}$$

Since  $y < 0$  and  $w_i \geq 0$  for all  $i$ , the above is at least  $w_{j'} - w_j = w_{j'}(t) - w_j(t)$  as required.  $\blacksquare$

**Lemma 15** *Suppose  $y < 0$ , and that  $\mathbf{w}(1)$  has positive entries which satisfy the theorem assumptions. If  $\eta \leq k^{-C}$  for some sufficiently large constant  $C$ , then the following hold for some iteration  $t_0 \leq \exp(\tilde{\mathcal{O}}(k))/\eta$ :*

- *There exists a unique  $j^* = \arg \min_i w_i(t_0)$ , and  $-\mathcal{O}(1) \cdot \eta \leq w_{j^*}(t_0) \leq 0$ .*
- *$\min_{j \neq j^*} w_j(t_0) \geq k^{-\mathcal{O}(1)}$ ,  $\max_{j \neq j^*} w_j(t_0) \leq \mathcal{O}(1)$ , and  $\max_j \prod_{i \notin \{j, j^*\}} w_i(t_0) \leq \mathcal{O}(1)$ .*

It is important to note that the constants hidden in the  $\mathcal{O}(\cdot)$  notation do not depend on  $\eta$  (although they may depend on  $C$ ).

**Proof** By Lemma 14 and the theorem assumptions, the following holds for all iterations  $t = 1, 2, \dots, T$  where  $\min_{i, t < T} w_i(t) \geq 0$ : There exists a unique  $j^* = \arg \min_i w_i(1)$ ,  $w_{j^*}(t)$  remains the unique smallest value among  $w_1(t), \dots, w_k(t)$ , and its distance from any other coordinate (which was initially  $k^{-\mathcal{O}(1)}$ ) is monotonically increasing in  $t$ . In particular, for any  $t < T$ ,  $\min_{j \neq j^*} w_j(t) \geq k^{-\mathcal{O}(1)}$ . As a result, recalling that  $y < 0$ , we have for all  $t < T$  that

$$\begin{aligned} w_{j^*}(t+1) &= w_{j^*}(t) - \eta \left( \prod_i w_i(t) - y \right) \prod_{i \neq j^*} w_i(t) \leq w_{j^*}(t) - \eta \left( \prod_i w_i(t) - y \right) \left( k^{-\mathcal{O}(1)} \right)^{k-1} \\ &\leq w_{j^*}(t) + \eta y \exp(-\tilde{\mathcal{O}}(k)) \leq w_{j^*}(t) - \eta \cdot \exp(-\tilde{\mathcal{O}}(k)). \end{aligned}$$

We also assume that initially  $w_{j^*}(1) \leq \mathcal{O}(1)$ . Therefore, after at most  $t_0 = \exp(\tilde{\mathcal{O}}(k))/\eta$  iterations, we will have  $w_{j^*}(t_0) \leq 0$  for the first time.

It remains to show that  $w_{j^*}(t_0) \geq -\mathcal{O}(1) \cdot \eta$ , as well as the second bullet in the lemma. To that end, we note that up till iteration  $t_0$ , for any  $j$ , both  $w_j(t)$  and  $\prod_{i \neq j} w_i(t)$  are monotonically decreasing in  $t$ , and moreover,  $t_0 > 1$  (since  $w_{j^*}(t_0) \leq 0$  and we assume  $w_{j^*}(1) > 0$ ). Thus, by Assumption 4,

$$\begin{aligned} w_{j^*}(t_0) &= w_{j^*}(t_0 - 1) - \eta \left( \prod_i w_i(t_0 - 1) - y \right) \prod_{i \neq j^*} w_i(t_0 - 1) \\ &> 0 - \eta \left( \prod_i w_i(1) - y \right) \prod_{i \neq j^*} w_i(1) \geq -\eta \cdot \mathcal{O}(1). \end{aligned}$$

Using this inequality, we have for any  $j \neq j^*$

$$w_j(t_0) = w_{j^*}(t_0) + (w_j(t_0) - w_{j^*}(t_0)) \geq -\eta \cdot \mathcal{O}(1) + k^{-\mathcal{O}(1)},$$

which is at least  $k^{-\mathcal{O}(1)}$  if  $\eta \leq k^{-C}$  for some sufficiently large constant  $C$ . Finally, since  $w_j(t)$  for any  $j \neq j^*$  is positive and monotonically decreasing up to iteration  $t_0$ , we have  $w_j(t_0) \leq w_j(1) \leq \mathcal{O}(1)$  and  $\prod_{i \notin \{j, j^*\}} w_i(t_0) \leq \prod_{i \notin \{j, j^*\}} w_i(1) \leq \mathcal{O}(1)$  by Assumption 4.  $\blacksquare$

**Lemma 16** *Under the conditions of Lemma 15,*

- *There exists a unique  $j^* = \arg \min_i w_i(t_0 + 1)$ , and  $w_{j^*}(t_0 + 1) \leq -\eta \cdot \exp(-\tilde{\mathcal{O}}(k))$ .*
- *$\min_{j \neq j^*} w_j(t_0 + 1) \geq k^{-\mathcal{O}(1)}$  and  $\prod_{i \neq j^*} w_i(t_0 + 1) \leq \mathcal{O}(1)$ .*

**Proof** By Lemma 15, we have  $\prod_i w_i(t_0) \leq 0$ , as well as  $\prod_i w_i(t_0) = w_{j^*}(t_0) \cdot w_j(t_0) \cdot \prod_{i \notin \{j, j^*\}} w_i(t_0) \geq -\mathcal{O}(1) \cdot \eta$  (where  $j$  is arbitrary). This implies that for sufficiently small  $\eta$ ,  $\frac{y}{2} \leq \prod_i w_i(t_0) \leq 0$ . By definition of the gradient descent update, it follows that  $w_{j^*}(t_0 + 1) \leq w_{j^*}(t_0)$  and for all  $j \neq j^*$ ,  $w_j(t_0 + 1) \geq w_j(t_0)$ , which implies that  $j^*$  remains the unique coordinate with smallest value as we move from iteration  $t_0$  to iteration  $t_0 + 1$ , as well as  $\min_{j \neq j^*} w_j(t_0 + 1) \geq \min_{j \neq j^*} w_j(t_0) \geq k^{-\mathcal{O}(1)}$ .

We now turn to prove  $w_{j^*}(t_0 + 1) \leq -\eta \cdot \exp(-\tilde{\mathcal{O}}(k))$ . Using the fact that  $\frac{y}{2} \leq \prod_i w_i(t_0) \leq 0$  as noted earlier,

$$w_{j^*}(t_0 + 1) = w_{j^*}(t_0) - \eta \left( \prod_i w_i(t_0) - y \right) \prod_{i \neq j^*} w_i(t_0) \leq 0 + \eta \cdot \frac{y}{2} \cdot (k^{-\mathcal{O}(1)})^{k-1} \leq -\eta \cdot \exp(-\tilde{\mathcal{O}}(k)).$$

Finally, to prove  $\prod_{i \neq j^*} w_i(t_0 + 1) \leq \mathcal{O}(1)$ , we have by definition that for any  $j \neq j^*$ ,

$$w_j(t_0 + 1) = w_j(t_0) - \eta \left( \prod_i w_i(t_0) - y \right) w_{j^*}(t_0) \cdot \prod_{i \notin \{j, j^*\}} w_i(t_0).$$

Using the fact that  $\frac{y}{2} \leq \prod_i w_i(t_0) \leq 0$  as shown earlier, and noting that by Lemma 15,  $|w_{j^*}(t_0)| \leq \mathcal{O}(1) \cdot \eta$  and  $\left| \prod_{i \notin \{j, j^*\}} w_i(t_0) \right| \leq \mathcal{O}(1)$ , it follows from the displayed equation above that  $w_j(t_0 + 1) \leq w_j(t_0) + \mathcal{O}(1) \cdot \eta^2$ . Therefore,

$$\prod_{i \neq j^*} w_i(t_0 + 1) \leq \prod_{i \neq j^*} (w_i(t_0) + \mathcal{O}(1) \cdot \eta^2) \leq \prod_{i \neq j^*} \left( w_i(t_0) \left( 1 + \frac{\mathcal{O}(1) \eta^2}{w_i(t_0)} \right) \right).$$

Since  $\min_{j \neq j^*} w_j(t_0) \geq k^{-\mathcal{O}(1)}$  (where the  $\mathcal{O}(1)$  does not depend on  $\eta$ ), then by picking  $\eta \leq k^{-C}$  for a sufficiently large  $C$ , the above is at most  $\left( \prod_{i \neq j^*} w_i(t_0) \right) \left( 1 + \frac{\mathcal{O}(1)}{k} \right)^{k-1} \leq \left( w_j(t_0) \prod_{i \notin \{j, j^*\}} w_i(t_0) \right) \cdot \mathcal{O}(1) \leq \mathcal{O}(1)$ , where we used Lemma 15 and where  $j$  is arbitrary.  $\blacksquare$



## A.2.3. PUTTING EVERYTHING TOGETHER

As discussed at the beginning of the proof, we can assume w.l.o.g. that  $w_1(1), \dots, w_k(1)$  are all positive (and in fact,  $\min_i w_i(1) \geq k^{-\mathcal{O}(1)}$  by our assumptions), and only consider the cases  $y > \prod_i w_i(1) > 0$  and  $y < 0 < \prod_i w_i(1)$ .

- If  $y > \prod_i w_i(1) > 0$ , we can apply Lemma 13 with  $\gamma = \delta = k^{-\mathcal{O}(1)}$  and any  $\eta = k^{-c}$  for some large enough constant  $c$ , to get a convergence to an  $\epsilon$ -optimal solution in  $\exp(\tilde{\mathcal{O}}(k)) \cdot \max\{1, \log(1/\epsilon)\}$  iterations.
- If  $y < 0 < \prod_i w_i(1)$ , and assuming  $\eta = k^{-c}$  for some large enough constant  $c > 0$ , then Lemma 15 and Lemma 16 together tell us that after at most  $\exp(\tilde{\mathcal{O}}(k))$  iterations, we get to an iteration  $t = t_0 + 1$  where  $w_{j^*}(t) \leq -\exp(-\tilde{\mathcal{O}}(k))$  for some  $j^*$ ,  $w_j(t) \geq k^{-\mathcal{O}(1)}$  for all  $j \neq j^*$ , and  $0 > \prod_i w_i(t) \geq -\mathcal{O}(1) \cdot \eta \geq -\mathcal{O}(1) \cdot k^{-\Omega(1)} > y$  for large enough  $k$ . Therefore, by Lemma 11, the dynamics of gradient descent from this time point is identical to case where we switch the signs of  $y$  and  $w_{j^*}$ , so that  $y > 0$ ,  $w_{j^*}(t) \geq \exp(-\tilde{\mathcal{O}}(k))$ ,  $w_j(t) \geq k^{-\mathcal{O}(1)}$  for all  $j \neq j^*$ , and  $y > \prod_i w_i(t) > 0$  for large enough  $k$ . Now applying Lemma 13 with  $\delta = \exp(-\tilde{\mathcal{O}}(k))$ ,  $\gamma = k^{-\mathcal{O}(1)}$ , and any step size  $\eta = k^{-c}$  for some large enough  $c$ , we get that  $\exp(\tilde{\mathcal{O}}(k)) \cdot \max\{1, \log(1/\epsilon)\}$  additional iterations suffice for convergence. Overall,  $\exp(\tilde{\mathcal{O}}(k)) \cdot \max\{1, \log(1/\epsilon)\}$  iterations are sufficient.