

A Proof of Lemma 3

Proof of Lemma 3. We first prove the upper bound of A_t . The essential proof is actually due to Cheung et al. [2019a] in analyzing sliding window based approach. For self-containedness, we restate here in the notations of our proposed restarted strategy.

$$\begin{aligned} & \left\| V_{t-1}^{-1} \left(\sum_{s=t_0}^{t-1} X_s X_s^T (\theta_s - \theta_t) \right) \right\|_2 \\ &= \left\| V_{t-1}^{-1} \left(\sum_{s=t_0}^{t-1} X_s X_s^T \left(\sum_{p=s}^{t-1} (\theta_p - \theta_{p+1}) \right) \right) \right\|_2 \\ &= \left\| V_{t-1}^{-1} \left(\sum_{p=t_0}^{t-1} \left(\sum_{s=t_0}^p X_s X_s^T (\theta_p - \theta_{p+1}) \right) \right) \right\|_2 \quad (21) \end{aligned}$$

$$\leq \sum_{p=t_0}^{t-1} \left\| V_{t-1}^{-1} \left(\sum_{s=t_0}^p X_s X_s^T (\theta_p - \theta_{p+1}) \right) \right\|_2 \quad (22)$$

$$\leq \sum_{p=t_0}^{t-1} \lambda_{\max} \left(V_{t-1}^{-1} \left(\sum_{s=t_0}^p X_s X_s^T \right) \right) \|\theta_p - \theta_{p+1}\|_2 \quad (23)$$

$$\leq \|\theta_p - \theta_{p+1}\|_2, \quad (24)$$

where (21) holds by rearranging over the index pair of (s, p) , (22) holds due to the triangle inequality, (23) and (24) can be obtained by the same argument in Appendix B of Cheung et al. [2019b]. We thus prove the upper bound of A_t .

We proceed to prove the upper bound of B_t . From the self-normalized concentration inequality [Abbasi-Yadkori et al., 2011, Theorem 1], restated in Theorem 5 of Appendix C, we know that

$$\begin{aligned} & \left\| \sum_{s=t_0}^{t-1} \eta_s X_s \right\|_{V_{t-1}^{-1}} \\ & \stackrel{(32)}{\leq} \sqrt{2R^2 \log \left(\frac{\det(V_{t-1})^{1/2} \det(\lambda I)^{-1/2}}{\delta} \right)} \\ & \leq R \sqrt{2 \log \frac{1}{\delta} + d \log \left(1 + \frac{(t-t_0)L^2}{d} \right)}, \end{aligned}$$

where the last inequality is obtained from the analysis of the determinant, as shown in the proof of Lemma 4.

Meanwhile, since $V_{t-1} \succeq \lambda I$, we know that

$$\|\lambda \theta_t\|_{V_{t-1}^{-1}}^2 \leq 1/\lambda_{\min}(V_{t-1}) \|\lambda \theta_t\|_2^2 \leq \frac{1}{\lambda} \|\lambda \theta_t\|_2^2 \leq \lambda S^2.$$

Therefore, the upper bound of B_t can be immediately obtained by combining the above inequalities. \square

B Bandit-over-Bandits Mechanism and Proof of Theorem 4

The RestartUCB algorithm requires prior information of the path-length P_T , which is generally unknown. Such a limitation can be avoided by utilizing the Bandits-over-bandits (BOB) mechanism, proposed by Cheung et al. [2019a] in designing parameter-free algorithm for non-stationary linear bandits based on sliding window least square estimator.

In the following, we first describe how to apply the BOB mechanism to eliminate the requirement of the unknown path-length in RestartUCB. Then, we present the proof of Theorem 4.

B.1 RestartUCB with BOB Mechanism

We name the RestartUCB algorithm with Bandit-over-Bandits mechanism as ‘‘RestartUCB-BOB’’, whose main idea is illustrated in Figure 4.

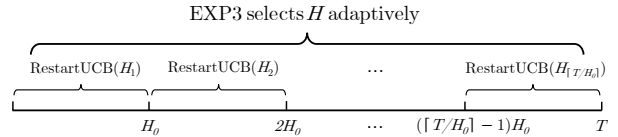


Figure 4: Illustration of Bandit-over-Bandits mechanism with application to RestartUCB algorithm.

From a high-level view, although the exact value of the optimal epoch size (or equivalently, the path-length P_T) is not clear, we can make some random guesses of its possible value, since the P_T is always bounded. Then, we can use a certain meta-algorithm to adaptively track the best epoch size, based on the returned reward returned. Specifically, The RestartUCB-BOB algorithm first sets an update period H_0 , and then runs the RestartUCB with a particular epoch size in each period, and the epoch size will be adaptively adjusted by employing EXP3 [Auer et al., 2002] as the meta-algorithm. We refer the reader to Section 7.3 of Cheung et al. [2019b] for more descriptions of design motivations and algorithmic details.

In the configuration of RestartUCB-BOB, we set $H_0 = \lceil d\sqrt{T} \rceil$ and the pool of epoch sizes J as

$$J = \{H_i = \lfloor (d/(2S))^{2/3} \cdot 2^{i-1} \rfloor \mid i = 1, 2, \dots, N\},$$

where $N = \lceil \ln(d^{1/3} T^{1/2} (2S)^{2/3}) \rceil + 1$.

Denoted by H_{\min} (H_{\max}) the minimal (maximal) epoch size in the pool J , we know that

$$H_{\min} = \lfloor (d/(2S))^{2/3} \rfloor, H_{\max} = \lfloor d\sqrt{T} \rfloor \leq H_0. \quad (25)$$

B.2 Proof of Theorem 4

Proof of Theorem 4. We begin with the following decomposition of the dynamic regret.

$$\begin{aligned}
 & \sum_{t=1}^T \langle X_t^*, \theta_t \rangle - \langle X_t, \theta_t \rangle \\
 = & \underbrace{\sum_{t=1}^T \langle X_t^*, \theta_t \rangle - \sum_{i=1}^{\lceil T/H_0 \rceil} \sum_{t=(i-1)H_0+1}^{iH_0} \langle X_t(H^\dagger), \theta_t \rangle}_{\text{term (i)}} \\
 & + \underbrace{\sum_{i=1}^{\lceil T/H_0 \rceil} \sum_{t=(i-1)H_0+1}^{iH_0} \langle X_t(H^\dagger), \theta_t \rangle - X_t(H_i), \theta_t \rangle}_{\text{term (ii)}},
 \end{aligned}$$

where H^\dagger is the best epoch size to approximate the optimal epoch size H^* in the pool J , and $H^* = \lfloor (dT/(1+P_T))^{2/3} \rfloor$. Hence, it suffices to bound terms (i) and (ii). In the following, we consider two cases, either $(1+P_T) \geq d^{-1/2}T^{1/4}$ or $(1+P_T) < d^{-1/2}T^{1/4}$.

Case 1. when $(1+P_T) \geq d^{-1/2}T^{1/4}$.

In this case, it is easy to verify that $H^* \leq H_{\max}$ and we thus conclude that H^* lies in the the range of $[H_{\min}, H_{\max}]$. Furthermore, from the configuration of the pool J , we confirm that there exists an epoch size $H^\dagger \in J$ such that $H^\dagger \leq H^* \leq 2H^\dagger$. So term (ii) can be upper bounded by

$$\text{term (ii)} \leq \sum_{i=1}^{\lceil T/H_0 \rceil} \tilde{O}\left(H^\dagger P_i + \frac{dH_0}{\sqrt{H^\dagger}}\right) \quad (26)$$

$$\begin{aligned}
 & = \tilde{O}\left(H^\dagger P_T + \frac{dT}{\sqrt{H^\dagger}}\right) \quad (27) \\
 & \leq \tilde{O}\left(H^* P_T + \frac{dT}{\sqrt{2H^*}}\right) \\
 & = \tilde{O}(d^{2/3}P_T^{1/3}T^{2/3}),
 \end{aligned}$$

where (26) is due to Theorem 2 and P_i denotes the path-length in the i -th update period. (27) follows by summing over all update periods, and the last inequality holds since the optimal epoch size H^* is provably in the range of $[H_{\min}, H_{\max}]$ and satisfies $H^\dagger \leq H^* \leq 2H^\dagger$.

Next, we bound the term (i),

$$\begin{aligned}
 \text{term (i)} & \leq \tilde{O}(\sqrt{H_0NT}) \\
 & \leq \tilde{O}(d^{1/2}T^{3/4}) \quad (28) \\
 & \leq \tilde{O}(d^{2/3}T^{2/3}(1+P_T)^{1/3}),
 \end{aligned}$$

where the first inequality follows by the same argument as in the sliding window based approach [Cheung et al., 2019b, Lemma 13], building upon the of EXP3. In

addition, the last inequality holds due to the fact that $(1+P_T) \geq d^{-1/2}T^{1/4}$ implies,

$$d^{1/2}T^{3/4} = d^{2/3}T^{2/3}d^{-1/3}T^{1/6} \leq d^{2/3}T^{2/3}(1+P_T)^{1/3}.$$

Hence, by combining the upper bounds of term (i) and term (ii), we know that the dynamic regret of RestartUCB-BOB is bounded by $\tilde{O}(d^{2/3}T^{2/3}(1+P_T)^{1/3})$ under the condition of $(1+P_T) \geq d^{-1/2}T^{1/4}$.

Case 2. when $(1+P_T) < d^{-1/2}T^{1/4}$.

In this case, we cannot guarantee that the optimal epoch size H^* lies in the range of $[H_{\min}, H_{\max}]$, so we set H^\dagger as H_0 ,

$$\begin{aligned}
 \text{term (ii)} & \leq \tilde{O}\left(H^\dagger P_T + \frac{dT}{\sqrt{H^\dagger}}\right) \\
 & \leq \tilde{O}\left(H_0 P_T + \frac{dT}{\sqrt{H_0}}\right) \\
 & = \tilde{O}\left(d\sqrt{T}P_T + d^{1/2}T^{3/4}\right) \\
 & \leq \tilde{O}\left(d^{1/2}T^{3/4}\right)
 \end{aligned}$$

where the last inequality holds by exploiting the condition of $(1+P_T) \leq d^{-1/2}T^{1/4}$. The result in conjunction with the upper bound of term (i) in (28) gives the $\tilde{O}(d^{1/2}T^{3/4})$ dynamic regret under this condition.

Finally, note that the dynamic regret of above two cases can be rewritten in the following unified form,

$$\text{term (i)} + \text{term (ii)} \leq \tilde{O}\left(d^{2/3}T^{2/3}(\max\{P_T, d^{-1/2}T^{1/4}\})^{1/3}\right).$$

Hence, we complete the proof of Theorem 4. \square

C Technical Lemmas

In this section, we provide several technical lemmas that frequently used in the proofs.

Theorem 5 (Self-Normalized Bound for Vector-Valued Martingales [Abbasi-Yadkori et al., 2011, Theorem 1]). *Let $\{F_t\}_{t=0}^\infty$ be a filtration. Let $\{\eta_t\}_{t=0}^\infty$ be a real-valued stochastic process such that η_t is F_t -measurable and conditionally R -sub-Gaussian for some $R > 0$, namely,*

$$\forall \lambda \in \mathbb{R}, \quad \mathbb{E}[\exp(\lambda\eta_t)|F_{t-1}] \leq \exp\left(\frac{\lambda^2 R^2}{2}\right). \quad (29)$$

Let $\{X_t\}_{t=1}^\infty$ be an \mathbb{R}^d -valued stochastic process such that X_t is F_{t-1} -measurable. Assume that V is a $d \times d$ positive definite matrix. For any $t \geq 0$, define

$$\bar{V}_t = V + \sum_{\tau=1}^t X_\tau X_\tau^\top, \quad S_t = \sum_{\tau=1}^t \eta_\tau X_\tau. \quad (30)$$

Then, for any $\delta > 0$, with probability at least $1 - \delta$, for all $t \geq 0$,

$$\|S_t\|_{\bar{V}_t}^2 \leq 2R^2 \log \left(\frac{\det(\bar{V}_t)^{1/2} \det(V)^{-1/2}}{\delta} \right). \quad (31)$$

Lemma 4 (Elliptical Potential Lemma). *Suppose $U_0 = \lambda I$, $U_t = U_{t-1} + X_t X_t^\top$, and $\|X_t\|_2 \leq L$, then*

$$\sum_{t=1}^T \|U_{t-1}^{-\frac{1}{2}} X_t\|_2 \leq \sqrt{2dT \log \left(1 + \frac{L^2 T}{\lambda d} \right)}. \quad (32)$$

Proof. First, we have the following decomposition,

$$U_t = U_{t-1} + X_t X_t^\top = U_{t-1}^{\frac{1}{2}} (I + U_{t-1}^{-\frac{1}{2}} X_t X_t^\top U_{t-1}^{-\frac{1}{2}}) U_{t-1}^{\frac{1}{2}}.$$

Taking the determinant on both sides, we get

$$\det(U_t) = \det(U_{t-1}) \det(I + U_{t-1}^{-\frac{1}{2}} X_t X_t^\top U_{t-1}^{-\frac{1}{2}}),$$

which in conjunction with Lemma 5 yields

$$\begin{aligned} \det(U_t) &= \det(U_{t-1}) (1 + \|U_{t-1}^{-\frac{1}{2}} X_t\|_2^2) \\ &\geq \det(U_{t-1}) \exp(\|U_{t-1}^{-\frac{1}{2}} X_t\|_2^2 / 2). \end{aligned}$$

Note that in the first inequality, we utilize the fact that $1 + x \geq \exp(x/2)$ holds for any $x \in [0, 1]$. By taking advantage of the telescope structure, we have

$$\sum_{t=1}^T \|U_{t-1}^{-\frac{1}{2}} X_t\|_2^2 \leq 2 \log \frac{\det(U_T)}{\det(U_0)} \leq 2d \log \left(1 + \frac{L^2 T}{\lambda d} \right),$$

where the last inequality follows from the fact that $\text{Tr}(U_T) \leq \text{Tr}(U_0) + L^2 T = \lambda d + L^2 T$, and thus $\det(U_T) \leq (\lambda + L^2 T/d)^d$.

Therefore, Cauchy-Schwartz inequality gives,

$$\begin{aligned} \sum_{t=1}^T \|U_{t-1}^{-\frac{1}{2}} X_t\|_2 &\leq \sqrt{T \sum_{t=1}^T \|U_{t-1}^{-\frac{1}{2}} X_t\|_2^2} \\ &\leq \sqrt{2dT \log \left(1 + \frac{L^2 T}{\lambda d} \right)}. \end{aligned}$$

□

Lemma 5.

$$\det(I + \mathbf{v} \mathbf{v}^\top) = 1 + \|\mathbf{v}\|_2^2. \quad (33)$$

Proof. Notice that

- (i) $(I + \mathbf{v} \mathbf{v}^\top) \mathbf{v} = (1 + \|\mathbf{v}\|_2^2) \mathbf{v}$, therefore, \mathbf{v} is its eigenvector with $(1 + \|\mathbf{v}\|_2^2)$ as the eigenvalue;
- (ii) $(I + \mathbf{v} \mathbf{v}^\top) \mathbf{v}^\perp = \mathbf{v}^\perp$, therefore, $\mathbf{v}^\perp \perp \mathbf{v}$ is its eigenvector with 1 as the eigenvalue.

Consequently, $\det(I + \mathbf{v} \mathbf{v}^\top) = 1 + \|\mathbf{v}\|_2^2$. □