

# Weighted Model Counting with Conditional Weights for Bayesian Networks (Supplementary Material)

Paulius Dilkas<sup>1</sup>

Vaishak Belle<sup>1</sup>

<sup>1</sup>University of Edinburgh, Edinburgh, UK

## 1 PROOFS

**Theorem 1.** *The function  $\mu_\nu$  is a measure.*

*Proof.* Note that  $\mu_\nu(\perp) = 0$  since there are no atoms below  $\perp$ . Let  $a, b \in 2^{2^U}$  be such that  $a \wedge b = \perp$ . By elementary properties of Boolean algebras, all atoms below  $a \vee b$  are either below  $a$  or below  $b$ . Moreover, none of them can be below both  $a$  and  $b$  because then they would have to be below  $a \wedge b = \perp$ . Thus

$$\begin{aligned} \mu_\nu(a \vee b) &= \sum_{\{u\} \leq a \vee b} \nu(u) = \sum_{\{u\} \leq a} \nu(u) + \sum_{\{u\} \leq b} \nu(u) \\ &= \mu_\nu(a) + \mu_\nu(b) \end{aligned}$$

as required.  $\square$

**Theorem 3.** *For any set  $U$  and measure  $\mu: 2^{2^U} \rightarrow \mathbb{R}_{\geq 0}$ , there exists a set  $V \supseteq U$ , a factorable measure  $\mu': 2^{2^V} \rightarrow \mathbb{R}_{\geq 0}$ , and a formula  $f \in 2^{2^V}$  such that  $\mu(x) = \mu'(x \wedge f)$  for all formulas  $x \in 2^{2^U}$ .*

*Proof.* Let  $V = U \cup \{f_m \mid m \in 2^U\}$ , and  $f = \bigwedge_{m \in 2^U} \{m\} \leftrightarrow f_m$ . We define weight function  $\nu: 2^V \rightarrow \mathbb{R}_{\geq 0}$  as  $\nu = \prod_{v \in V} \nu_v$ , where  $\nu_v(\{v\}) = \mu(\{m\})$  if  $v = f_m$  for some  $m \in 2^U$  and  $\nu_v(x) = 1$  for all other  $v \in V$  and  $x \in 2^{\{v\}}$ . Let  $\mu': 2^{2^V} \rightarrow \mathbb{R}_{\geq 0}$  be the measure induced by  $\nu$ . It is enough to show that  $\mu$  and  $x \mapsto \mu'(x \wedge f)$  agree on the atoms in  $2^{2^U}$ . For any  $\{a\} \in 2^{2^U}$ ,

$$\begin{aligned} \mu'(\{a\} \wedge f) &= \sum_{\{x\} \leq \{a\} \wedge f} \nu(x) = \nu(a \cup \{f_a\}) \\ &= \nu_{f_a}(\{f_a\}) = \mu(\{a\}) \end{aligned}$$

as required.  $\square$

**Lemma 1.** *Let  $X \in \mathcal{V}$  be a random variable with parents  $\text{pa}(X) = \{Y_1, \dots, Y_n\}$ . Then  $\text{CPT}_X: 2^{\mathcal{E}^*(X)} \rightarrow$*

$\mathbb{R}_{\geq 0}$  is such that for any  $x \in \text{im } X$  and  $(y_1, \dots, y_n) \in \prod_{i=1}^n \text{im } Y_i$ ,

$$\text{CPT}_X(T) = \Pr(X = x \mid Y_1 = y_1, \dots, Y_n = y_n),$$

where  $T = \{\lambda_{X=x}\} \cup \{\lambda_{Y_i=y_i} \mid i = 1, \dots, n\}$ .

*Proof.* If  $X$  is binary, then  $\text{CPT}_X$  is a sum of  $2 \prod_{i=1}^n |\text{im } Y_i|$  terms, one for each possible assignment of values to variables  $X, Y_1, \dots, Y_n$ . Exactly one of these terms is nonzero when applied to  $T$ , and it is equal to  $\Pr(X = x \mid Y_1 = y_1, \dots, Y_n = y_n)$  by definition.

If  $X$  is not binary, then  $(\sum_{i=1}^m [\lambda_{X=x_i}])(T) = 1$ , and  $(\prod_{i=1}^m \prod_{j=i+1}^m (\overline{[\lambda_{X=x_i}]} + \overline{[\lambda_{X=x_j}]}))(T) = 1$ , so  $\text{CPT}_X(T) = \Pr(X = x \mid Y_1 = y_1, \dots, Y_n = y_n)$  by a similar argument as before.  $\square$

**Lemma 2.** *Let  $\mathcal{V} = \{X_1, \dots, X_n\}$ . Then*

$$\phi(T) = \begin{cases} \Pr(x_1, \dots, x_n) & \text{if } T = \{\lambda_{X_i=x_i}\}_{i=1}^n \text{ for} \\ & \text{some } (x_i)_{i=1}^n \in \prod_{i=1}^n \text{im } X_i \\ 0 & \text{otherwise,} \end{cases}$$

for all  $T \in 2^U$ .

*Proof.* If  $T = \{\lambda_{X=v_X} \mid X \in \mathcal{V}\}$  for some  $(v_X)_{X \in \mathcal{V}} \in \prod_{X \in \mathcal{V}} \text{im } X$ , then

$$\begin{aligned} \phi(T) &= \prod_{X \in \mathcal{V}} \Pr \left( X = v_X \mid \bigwedge_{Y \in \text{pa}(X)} Y = v_Y \right) \\ &= \Pr \left( \bigwedge_{X \in \mathcal{V}} X = v_X \right) \end{aligned}$$

by Lemma 1 and the definition of a Bayesian network. Otherwise there must be some non-binary random variable  $X \in \mathcal{V}$  such that  $|\mathcal{E}(X) \cap T| \neq 1$ . If  $\mathcal{E}(X) \cap T = \emptyset$ , then  $(\sum_{i=1}^m [\lambda_{X=x_i}])(T) = 0$ , and so  $\text{CPT}_X(T) = 0$ , and

$\phi(T) = 0$ . If  $|\mathcal{E}(X) \cap T| > 1$ , then we must have two different values  $x_1, x_2 \in \text{im } X$  such that  $\{\lambda_{X=x_1}, \lambda_{X=x_2}\} \subseteq T$  which means that  $([\lambda_{X=x_1}] + [\lambda_{X=x_2}])(T) = 0$ , and so, again,  $\text{CPT}_X(T) = 0$ , and  $\phi(T) = 0$ .  $\square$

**Theorem 4.** For any  $X \in \mathcal{V}$  and  $x \in \text{im } X$ ,

$$(\exists_U(\phi \cdot [\lambda_{X=x}])(\emptyset) = \Pr(X = x).$$

*Proof.* Let  $\mathcal{V} = \{X, Y_1, \dots, Y_n\}$ . Then

$$\begin{aligned} (\exists_U(\phi \cdot [\lambda_{X=x}])(\emptyset) &= \sum_{T \in 2^U} (\phi \cdot [\lambda_{X=x}])(T) \\ &= \sum_{\lambda_{X=x} \in T \in 2^U} \phi(T) \\ &= \sum_{\lambda_{X=x} \in T \in 2^U} \left( \prod_{Y \in \mathcal{V}} \text{CPT}_Y \right) (T) \\ &= \sum_{(y_i)_{i=1}^n \in \prod_{i=1}^n \text{im } Y_i} \Pr(x, y_1, \dots, y_n) \\ &= \Pr(X = x) \end{aligned}$$

by:

- the proof of Theorem 1 by Dudek et al. [2020];
- if  $\lambda_{X=x} \notin T \in 2^U$ , then  $(\phi \cdot [\lambda_{X=x}])(T) = \phi(T) \cdot [\lambda_{X=x}](T \cap \{\lambda_{X=x}\}) = \phi(T) \cdot 0 = 0$ ;
- Lemma 2;
- marginalisation of a probability distribution.

$\square$

## References

Jeffrey M. Dudek, Vu Phan, and Moshe Y. Vardi. ADDMC: weighted model counting with algebraic decision diagrams. In *The Thirty-Fourth AAAI Conference on Artificial Intelligence, AAAI 2020, The Thirty-Second Innovative Applications of Artificial Intelligence Conference, IAAI 2020, The Tenth AAAI Symposium on Educational Advances in Artificial Intelligence, EAAI 2020, New York, NY, USA, February 7-12, 2020*, pages 1468–1476. AAAI Press, 2020. ISBN 978-1-57735-823-7. URL <https://aaai.org/ojs/index.php/AAAI/article/view/5505>.